# BRAVO Biased Locking for Reader-Writer Locks 

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## Reader-Writer Locks

- Allow shared access for read-only use of a resource
- Ubiquitous in modern systems



## Reader-Writer Locks

- Allow shared access for read-only use of a resource
- Ubiquitous in modern systems

- Have to keep track of the presence of active readers


## The "shared counter" approach

## 0

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## 0



## The "shared counter" approach



## The "shared counter" approach



## The "shared counter" approach



## The "distributed" approach



## The "distributed" approach



The Scalable "Reader Indicator" Dilemma


## The BRAVO approach

 Compact \& Scalable
## The BRAVO approach



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## The BRAVO approach



## Evaluation

- Easy to integrate with existing locks
- Compact
- Accelerates reads
- Handles writes gracefully


## Evaluation: Easy to integrate

- Brandenburg-Anderson (BA) reader-writer lock
- POSIX Pthread reader-writer lock
- Linux kernel rwsem


## Evaluation: Compact

| Locks | Memory footprint |
| :--- | :--- |
| BA | 40 |
| BA + BRAVO | $40+12+32 \mathrm{~KB}$ (for a table) |
| Per-CPU | 9216 (on a system with 72 CPUs) |
| Cohort-RW | 896 (dual-socket) |


| Intel Xeon E5-2699 v3 CPU |
| :---: |
| 2 sockets |
| 72 logical CPUs in total |

## Evaluation: Accelerates reads <br> RWBench with 1 out of every 10000 are writes



## Evaluation: Handles writes gracefully <br> RWBench with 9 out of every 10 are writes



## Linux Kernel rwsem

- Counter + waiting queue protected by a spin lock
- Reader atomically increments the counter and checks its value


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- Stress-test with will-it-scale: page_fault and mmap


## Evaluation with will-it-scale

| Intel Xeon E7-8895 v3 CPU |
| :---: |
| 4 sockets |
| 144 logical CPUs in total |


page_fault

mmap

## Conclusion

- Builds into any existing lock
- Reads are accelerated
- Avoids write overhead
- Very compact

- Overall, takes the "reader indicator" dilemma away


## Thank you! Questions?

## Future Work

- Dynamic table sizing
- Probing multiple table locations
- Adaptive policies for enabling bias
- Revocation scan via SIMD instructions and non-temporal loads

