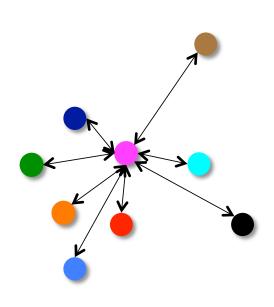
Unsynchronized Techniques for Approximate Parallel Computing

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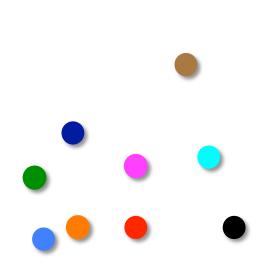
Barnes Hut N-Body Simulation



N Interacting Bodies (Stars, Molecules)

- At each step
 - Compute force
 - Acting on each body
 - From all other bodies

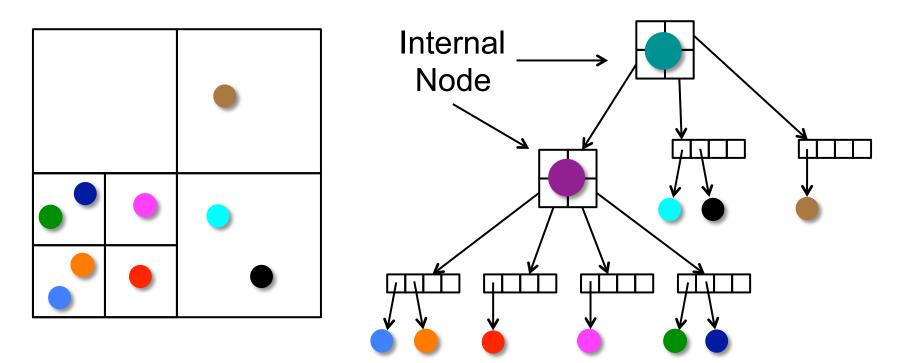
Barnes Hut N-Body Simulation



N Interacting Bodies (Stars, Molecules)

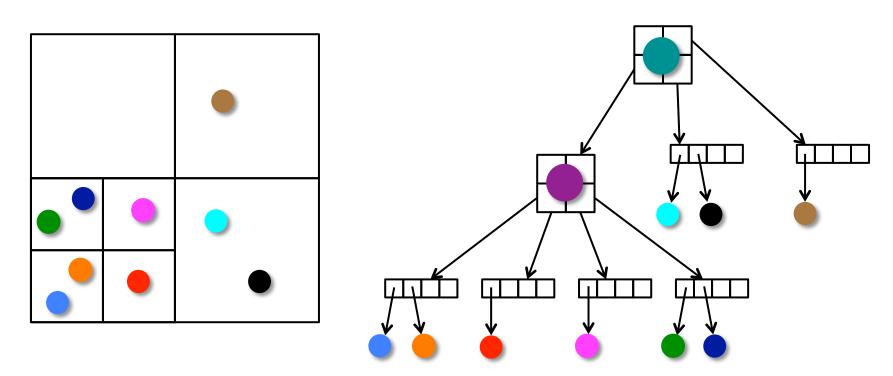
- At each step
 - Compute force
 - Acting on each body
 - From all other bodies
 - Use forces to move bodies
- Repeat

Space Subdivision Tree



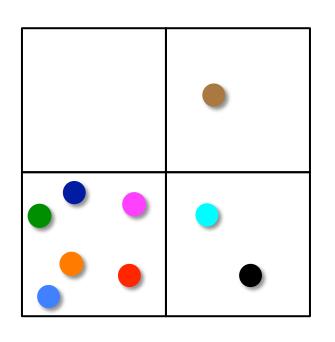
- Internal node contains center of mass for subtree
- Center of mass approximation for distant bodies
- Changes N² algorithm into N log N algorithm

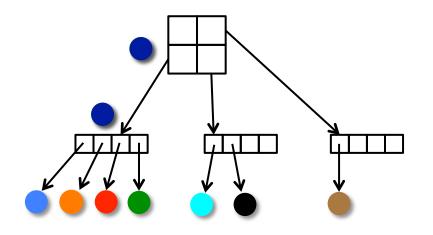
Barnes-Hut Algorithm



- 1) Build space subdivision tree
- 2) Use tree to compute forces acting on each body
- 3) Move bodies
- 4) Repeat

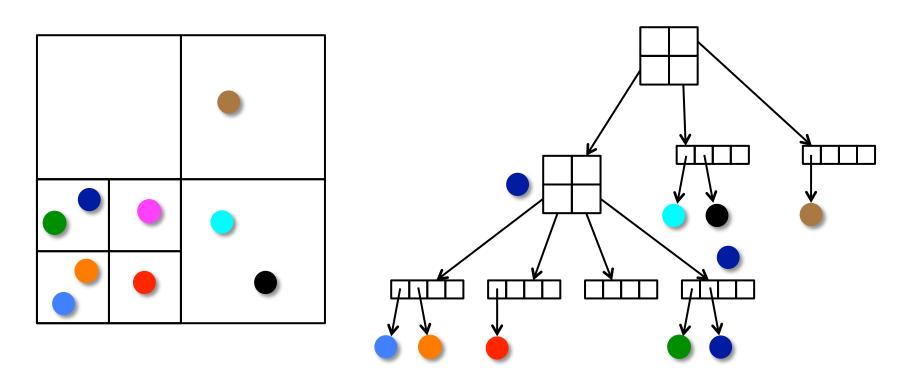
Tree Construction Algorithm





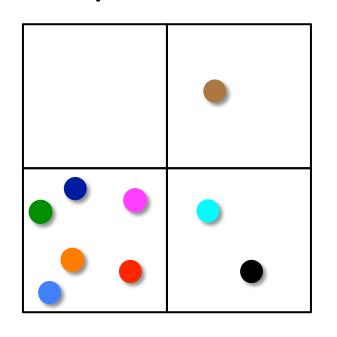
- 1) Drop bodies into the tree from the top
- 2) Bodies percolate down the tree
- 3) Insert each body in correct position
- 4) Add internal nodes as necessary

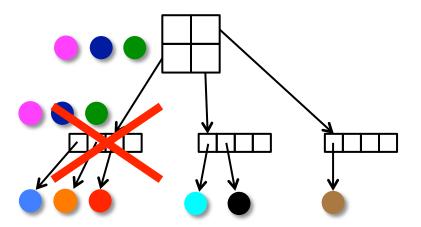
Tree Construction



- 1) Drop bodies into the tree from the top
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Parallel Tree Construction Algorithm, Version 1 – No Synchronization, Just Insert Bodies in Parallel



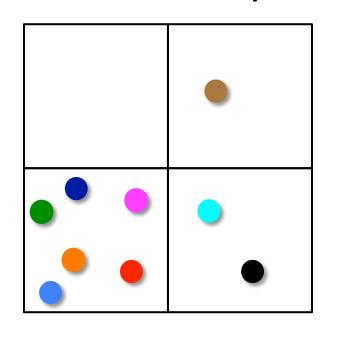


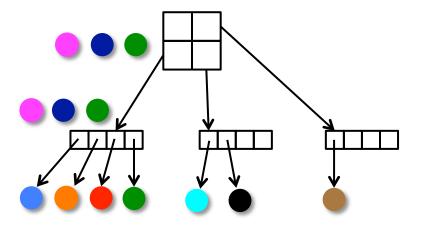
Data Races Corruption Crash!

- 1) Drop bodies into the tree from the top
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Standard Solution

- Add synchronization
 - Complicate program
 - Add synchronization overhead
 - Add space overhead
 - Maybe contention, even deadlock too
- We aren't going to do this



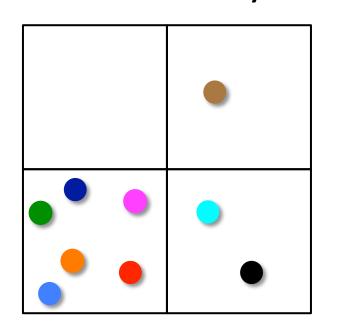


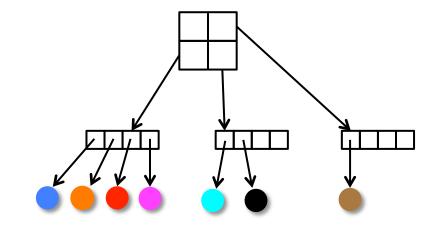
Potential Outcome:

Drop

•

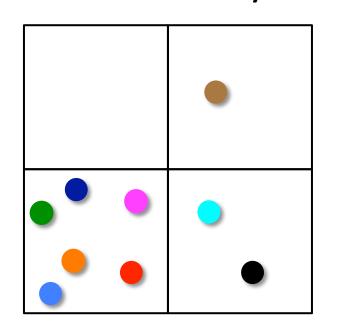
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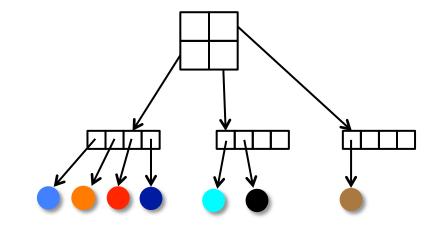




Dropped: • •

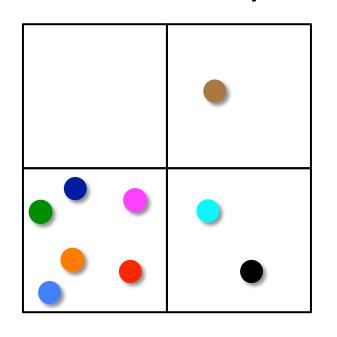
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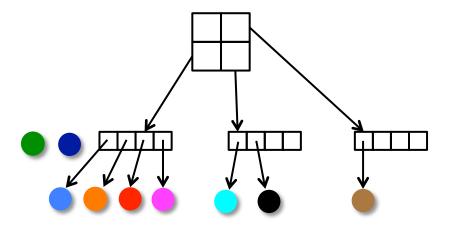




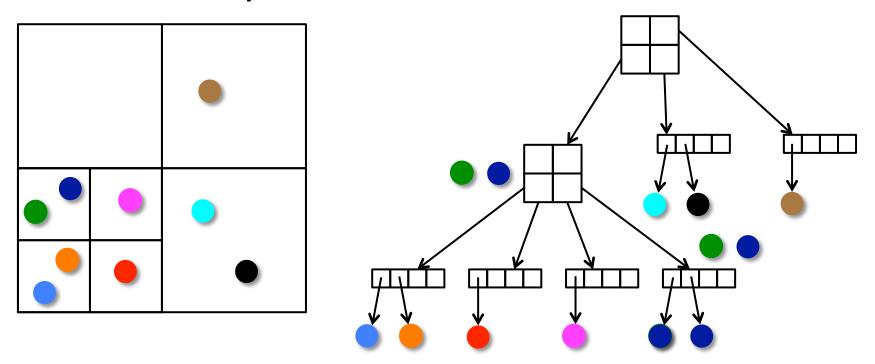
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Effect of Unsynchronized Construction

- Always produces a tree that force computation calculation can use
- But tree may not contain all bodies
 - Forces computed as if dropped bodies don't exist
 - Get an approximate force computation
 - But force computation is already approximate because of center of mass approximation
- Preserves integrity
- May affect accuracy

Parallel Tree Construction Options

Unsynchronized

- Data races
- Always produces a usable tree, but may drop bodies
 - Preserves integrity of computation
 - May affect accuracy

Tree Locking (standard approach)

- No data races, no dropped bodies
- Lots of synchronization, lots of contention at top of tree

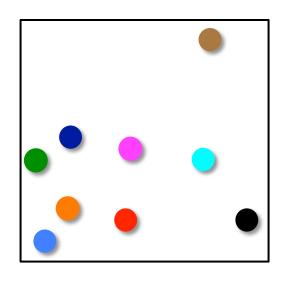
Update Locking (synchronize updates but not reads)

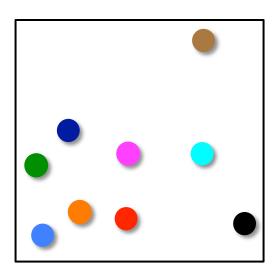
- Data races, no dropped bodies
- Some synchronization, little contention
- Complex, scary algorithm

Accuracy Evaluation

- Need a comparison point
- Full N² computation too expensive
- Compare with hyperaccurate version
 - -Uses a center of mass approximation
 - But goes deeper into the tree before using the center of mass approximation
 - So more accurate

Accuracy Metric





- Start with two corresponding configurations
 - Reference (hyperaccurate)
 - Comparison (tree locking, update locking, unsynchronized)
- Compute sum of distances between corresponding bodies
- Divide sum by distance between corners of bounding box

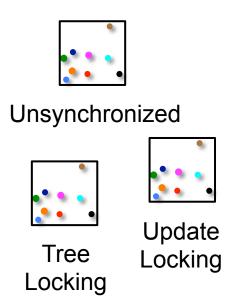
Accuracy Result

- Accuracy metric between
 - Hyperaccurate version and
 - All other versions
 - All other numbers of processors
- Is 1.02% (to three significant digits)

Visually



Hyperaccurate



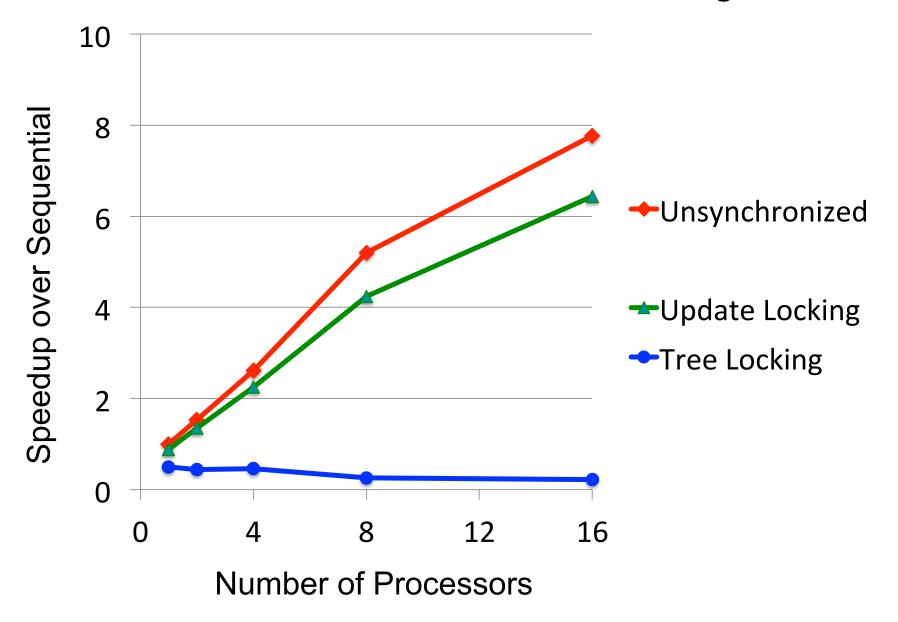
Accuracy Result in Context

- Difference between hyperaccurate and synchronized version is two orders of magnitude larger than difference between synchronized and unsynchronized versions
- Changing center of mass approximation has an accuracy effect
 two orders of magnitude more than removing synchronization

Performance Evaluation

- Implemented different versions
 - Tree locking
 - Update locking
 - Unsynchronized
 - Evaluated performance of different versions
- Measured speedup over sequential tree construction algorithm

Performance of Parallel Tree Construction Algorithms



Key Technical Concepts and Results

- Safe approximate unsynchronized data structures
 - Look just like standard sequential data structures
 - Give programmer sense of security and comfort
 - Building blocks in paper
- Accuracy evaluation methodology
 - To evaluate effect of unsynchronized data structures
 - Adjust accepted accuracy knobs
 - Compare with effect of removing synchronization
- Evidence shows programs are oversynchronized

Bigger Picture

- Field has traditionally aspired to perfection
- But perfection is increasingly unavailable
 - Huge systems, huge data sets
 - Something always broken
- Software is inherently resilient and malleable
- Enables more mature and productive approaches
- Goal is acceptability, not perfection
 - Integrity (program does not crash)
 - Accuracy (accurate enough, often enough)
 - End-to-end perspective
- Opens up new, counterintuitive possibilities