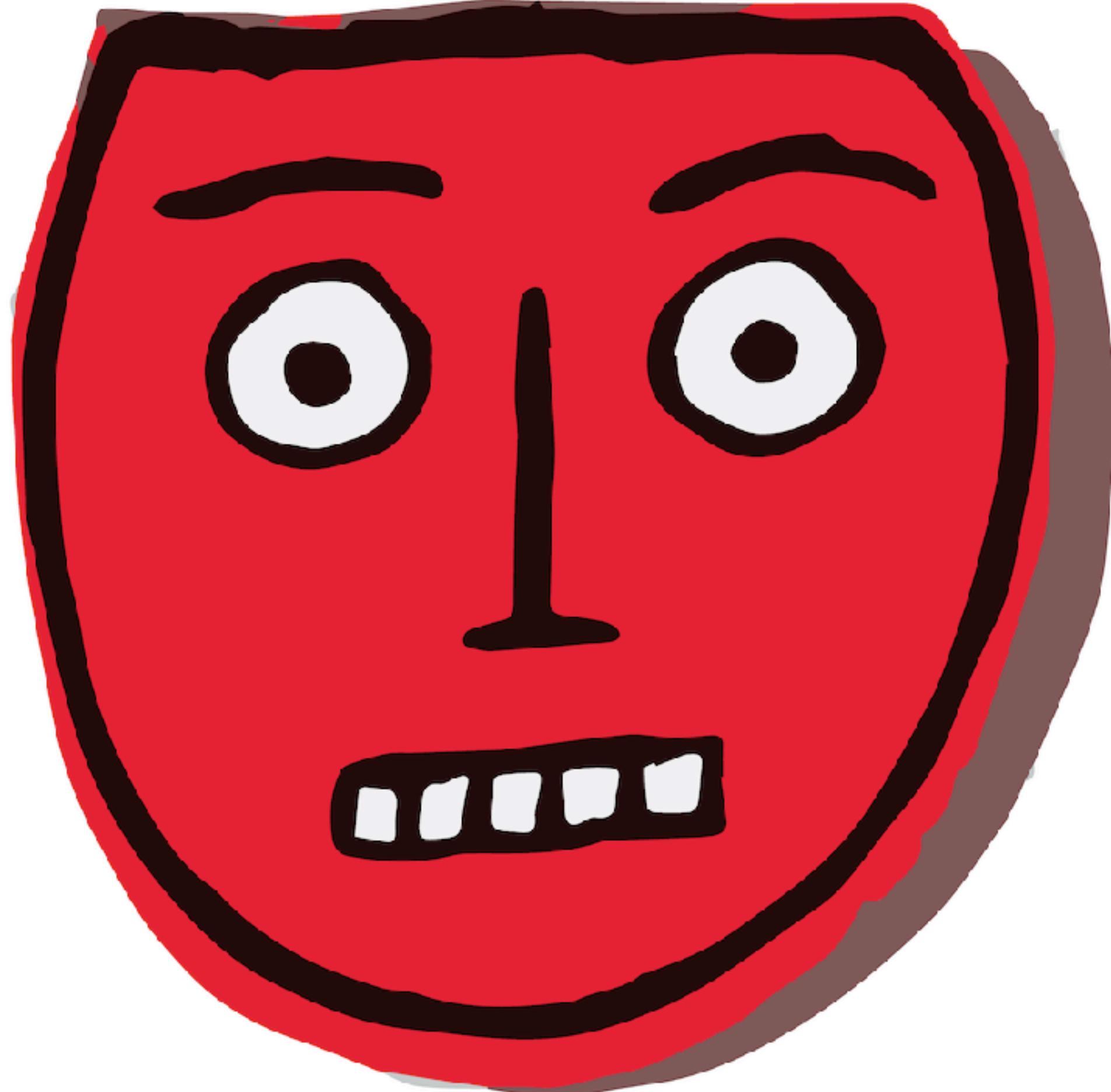


OK Log 

Distributed and
coördination-free
logging



@peterbourgon
github.com/peterbourgon
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fastly

Contextualizing

@peterbourgon



Outline

Gremlins of distsys

Logging systems

OK Log design



Outline



Gremlins of distsys

Logging systems

OK Log design



My incredible journey

- Software engineer — not computer scientist
- (Though I love a good paper)
- Distributed systems journeyman
- This shit is hard!!!



Why is it hard

- Failure modes are intractable
- Symptoms are subtle
- Diagnosis and triage requires specialized knowledge
- Computer science fundamentals are insufficient



You know this

```
>>> x = 1
```

```
>>> print x  
1
```



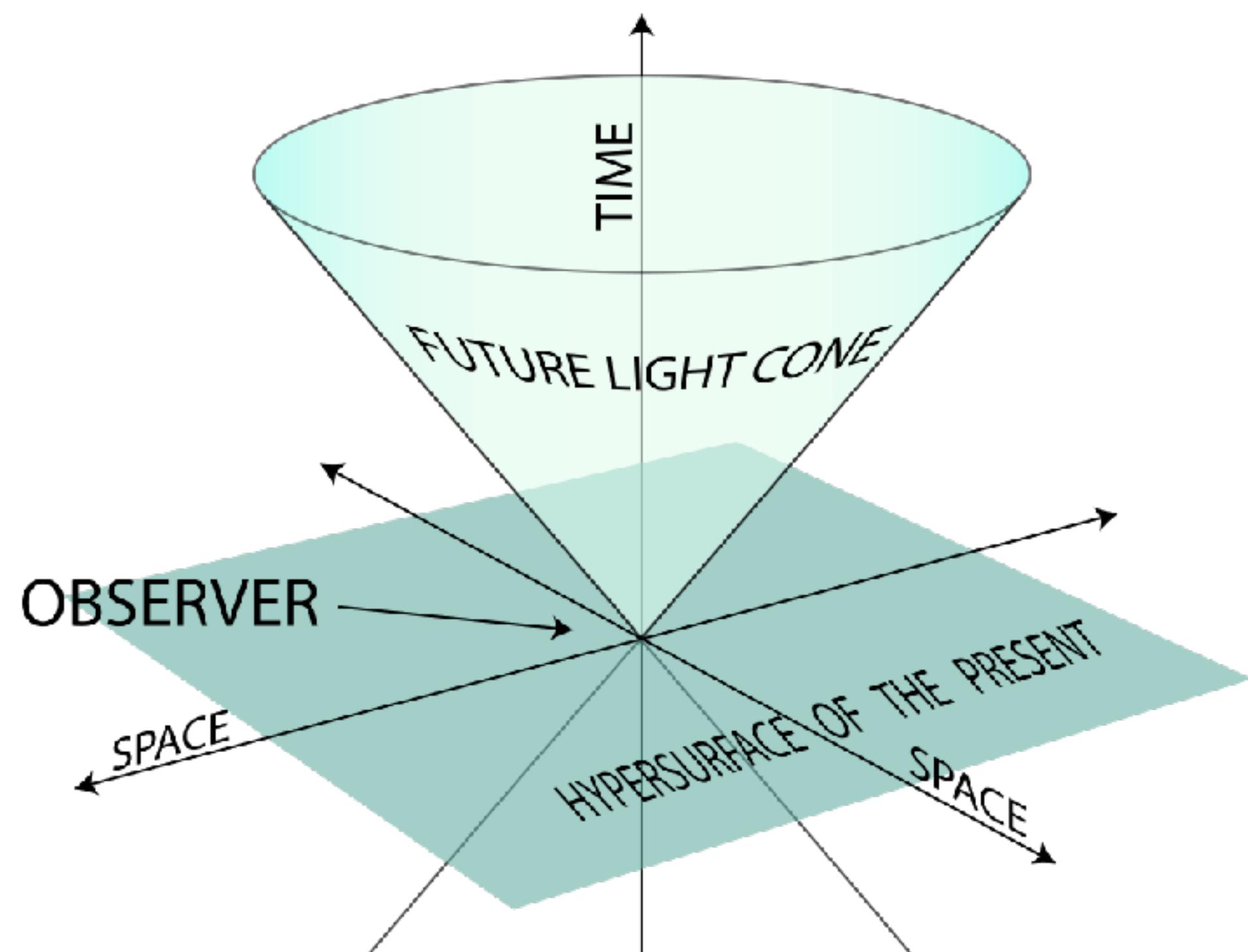
You know this

```
$ curl -XPOST http://db:9091/vars/x/1  
HTTP 502 (Bad Gateway)
```

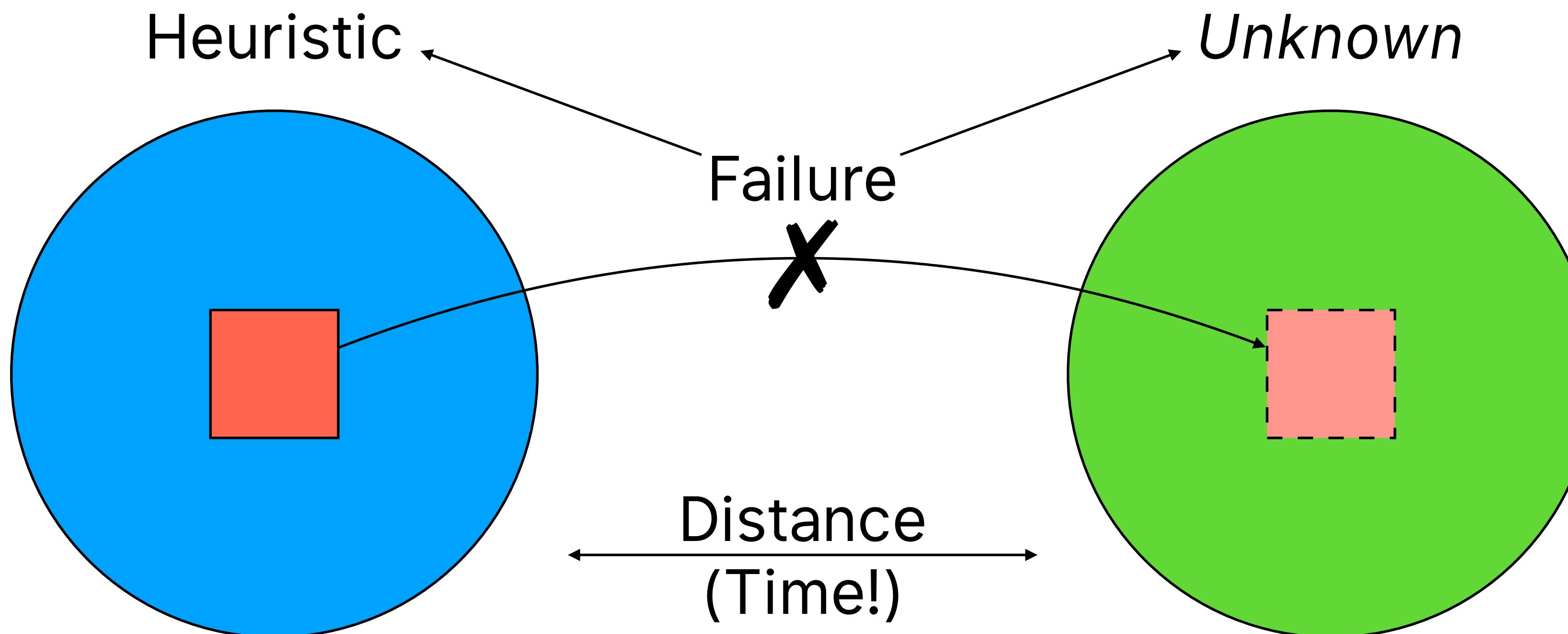
```
$ curl -XGET http://db:9091/vars/x  
HTTP 500 (Internal Server Error)
```



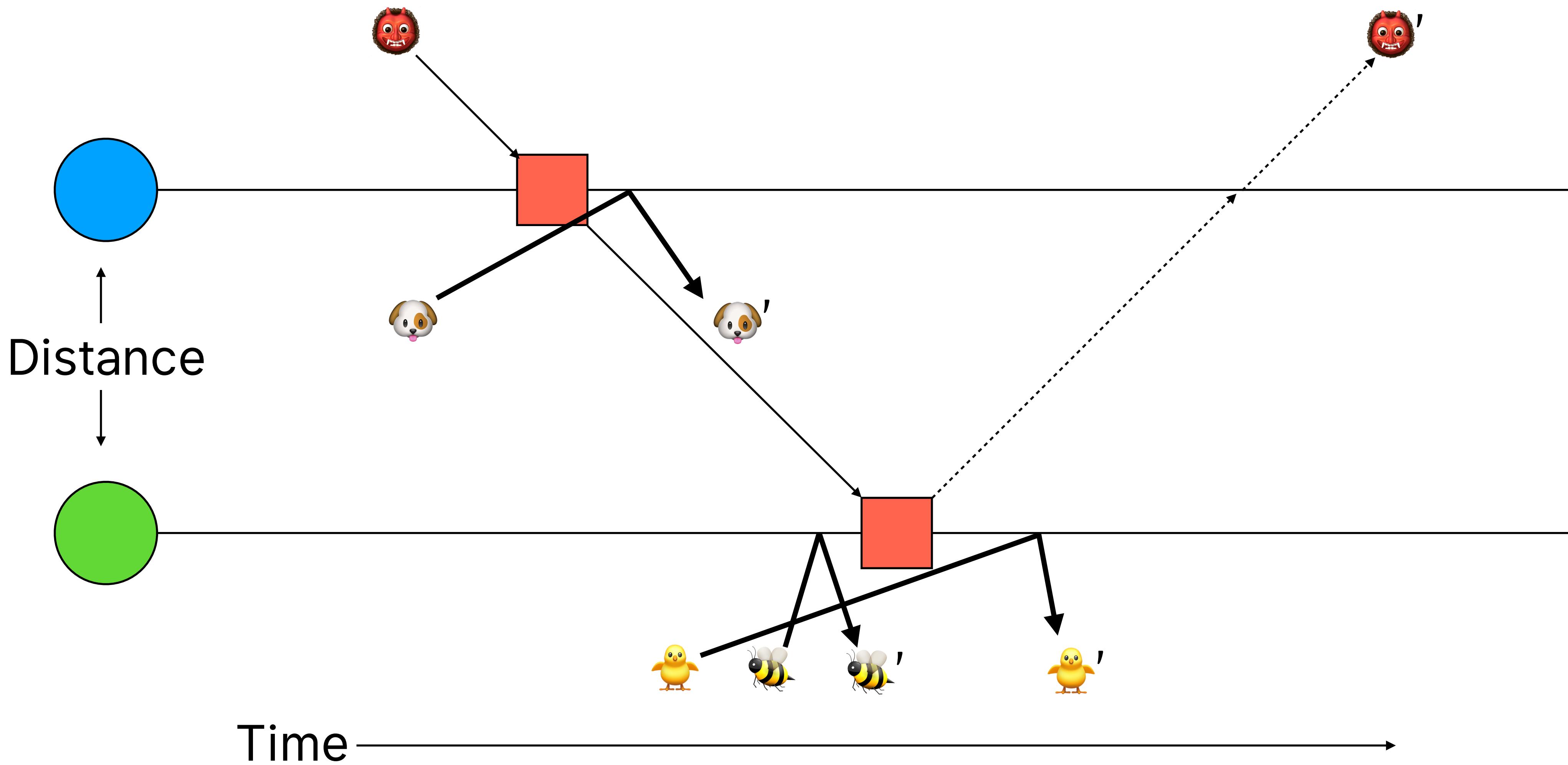
Fundamentally hard



Fundamentally hard



Fundamentally hard



Strong consistency models

Secure | https://aphyr.com/posts/313-strong-consistency-models

Light cones

However, this is not the full story: in almost every real-world system, processes are *distant* from each other. An uncached value in memory, for instance, is likely on a DIMM *thirty centimeters* away from the CPU. It takes light over a full nanosecond to travel that distance—and real memory accesses are much slower. A value on a computer in a different datacenter could be thousands of kilometers—hundreds of milliseconds—away. We just can't send information there any faster; physics, thus far, forbids it.

This means our operations are *no longer instantaneous*. Some of them might be so fast as to be negligible, but in full generality, operations *take time*. We *invoke* a write of a variable; the write travels to memory, or another computer, or the moon; the memory changes state; a confirmation travels back; and then we *know* the operation took place.

The delay in sending messages from one place to another implies *ambiguity* in the history of operations. If messages travel faster or slower, they could take place in unexpected orders. Here, the bottom process invokes a read when the value is `a`. While the read is in flight, the top process writes `b`—and by happenstance, its write arrives *before* the read. The bottom process finally completes its read and finds `b`, not `a`.

This history violates our concurrent register consistency model. The bottom process did *not* read the current value at the time it invoked the read. We might try to use the completion time, rather than the invocation time, as the “true time” of the operation, but this fails by symmetry as well; if the read arrives *before* the write, the process would receive `a` when the current value is `b`.

The diagram illustrates two processes and their light cones.
 Top process: A horizontal bar divided into two regions labeled 'a' and 'b'. Two dashed arrows point away from the 'b' region: one labeled 'w' pointing up-left and one labeled 'r'' pointing up-right.
 Bottom process: A similar horizontal bar divided into 'a' and 'b'. A solid arrow points down from the 'a' region to the 'b' region, labeled 'r'.
 Interaction: A solid arrow points from the 'b' region of the top process to the 'b' region of the bottom process, labeled 'r' = b.
 This visualizes how a read operation (r) in the bottom process can see the value 'b' written by a previous write (w) in the top process, due to the finite speed of light and message transmission.

The diagram shows two processes interacting with a shared variable.
 Top process: A horizontal bar divided into two regions. The left region contains 'w' and 'w'' above it, with an arrow pointing to the right. The right region contains 'b' above it.
 Bottom process: A similar horizontal bar divided into 'a' and 'b'. A solid arrow points down from the 'a' region to the 'b' region, labeled 'r'.
 Interaction: A solid arrow points from the 'b' region of the top process to the 'b' region of the bottom process, labeled 'r' = b.
 This visualizes a race condition where the bottom process's read 'r' might capture the value 'a' instead of 'b' if the write 'w' arrives after the read 'r' but before the write 'w'' completes.

It's phenomenological

- Immanuel Kant → Edmund Husserl → ...
- Actual thing (*Noumena*) vs. our experience of it (*Phenomena*)
- There is no "true" state — just a union of expressed intents
- Set aside the fiction of a global truth!
- We can only reason about each node's unique perception
- Successful distsys thinking starts with *phenomenological reduction*



Idioms

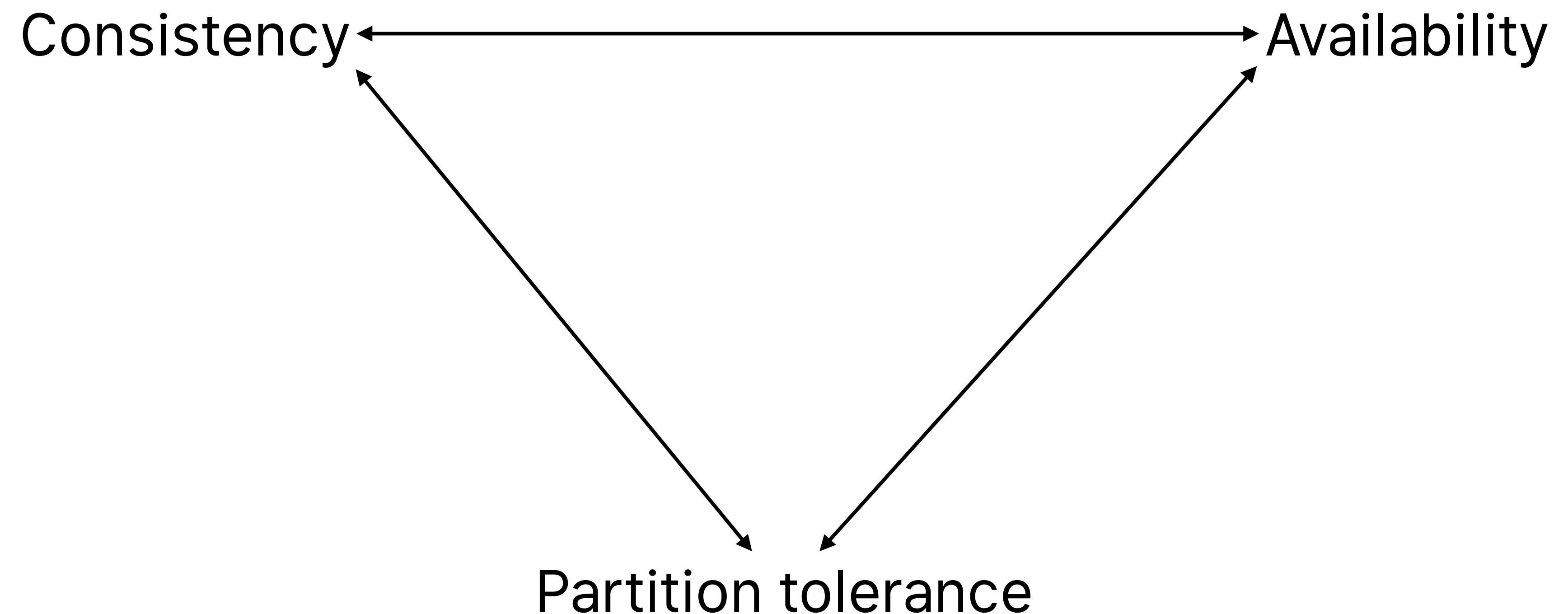
1980s: RPC

1990s: CORBA

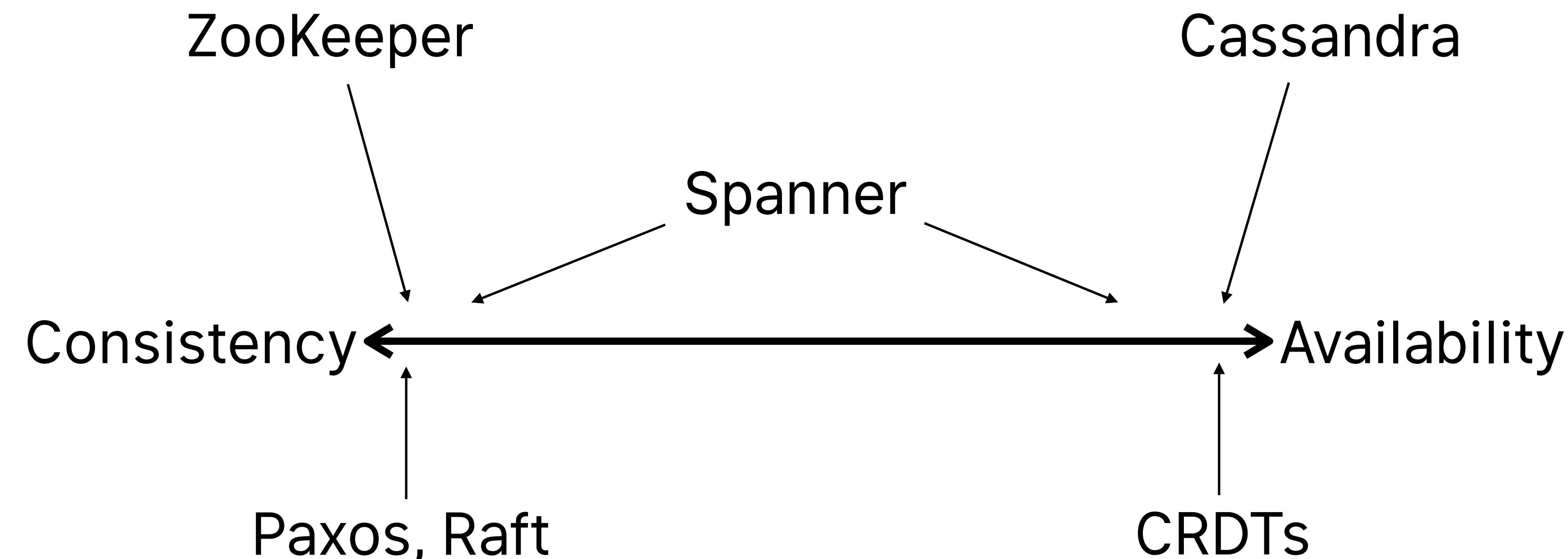
2000s: CAP



CAP



CAP



PACELC

- In the case of network **Partition**
- You have to choose between **Availability** and **Consistency**
- **Else**
- You can choose between **Latency** and **Consistency**
- → PA/EL, PA/EC, PC/EL, PC/EC



This is all very interesting

- Main point: situation is complex
- And becoming moreso



Danger: Hot Takes



Jonathan Edwards
@jonathoda

Follow ▾

Replying to @garybernhardt

We like it that way. Programmers wallow in complexity like pigs in shit.

4:28 AM - 20 Aug 2017

17 Retweets 29 Likes



2 17 29

@peterbourgon



Outline



Gremlins of distsys

Logging systems

OK Log design

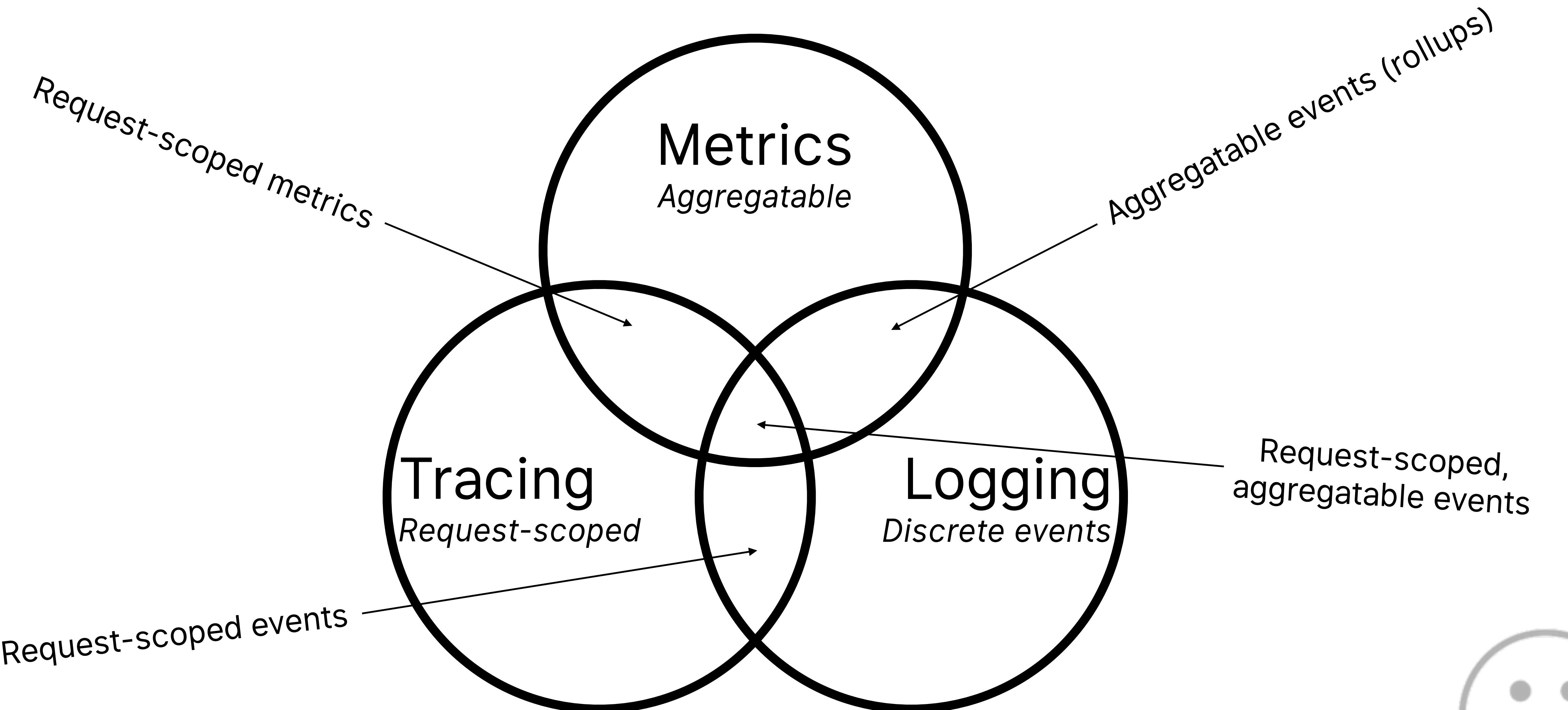


Outline

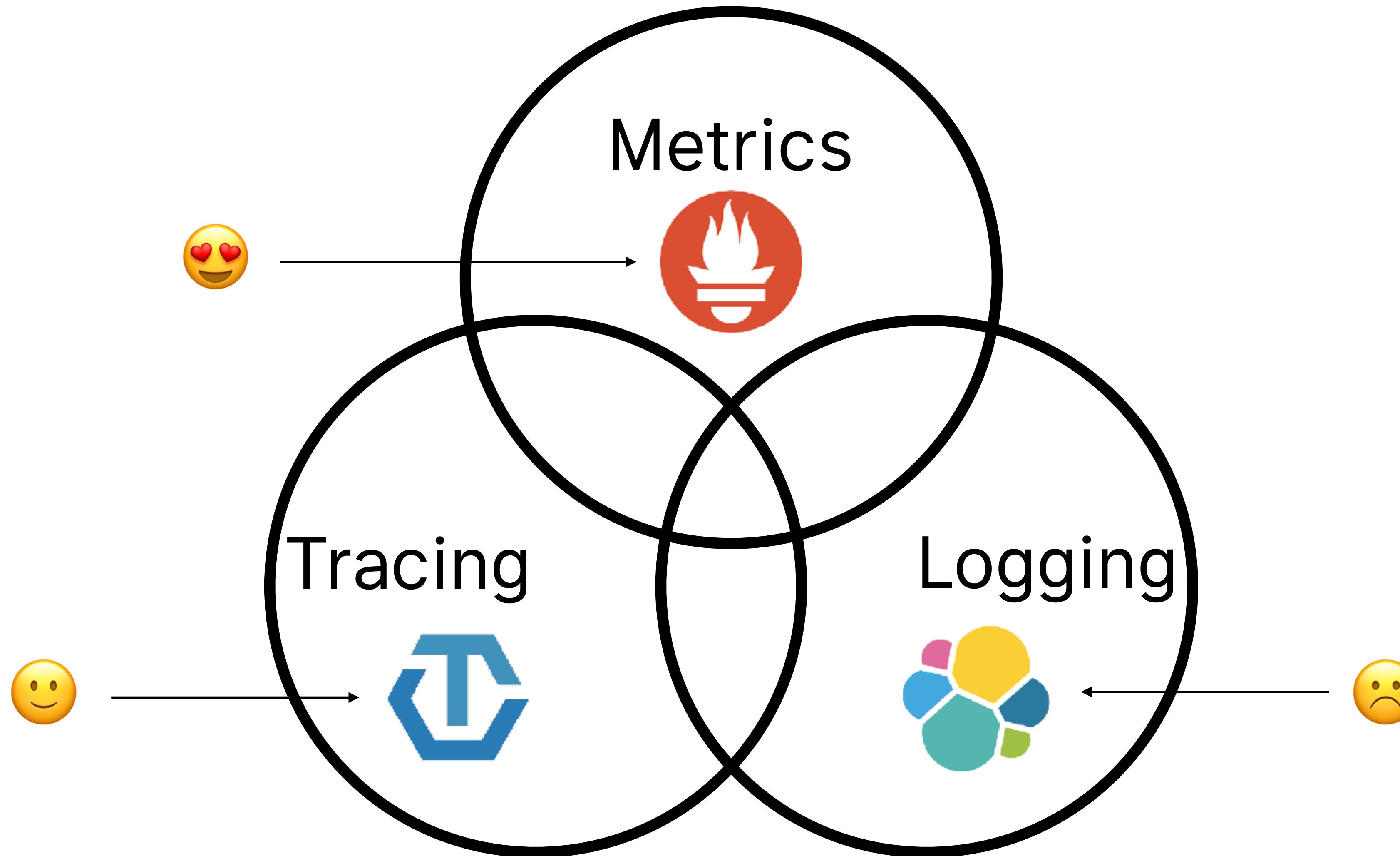
Gremlins of distsys

- 👉 Logging systems
- OK Log design

Observability



Observability





- Lucene not the best for log data
- Operationally complex, can be unstable
- Resource heavy
- Does too much — most of the time I just wanna, like, grep my logs

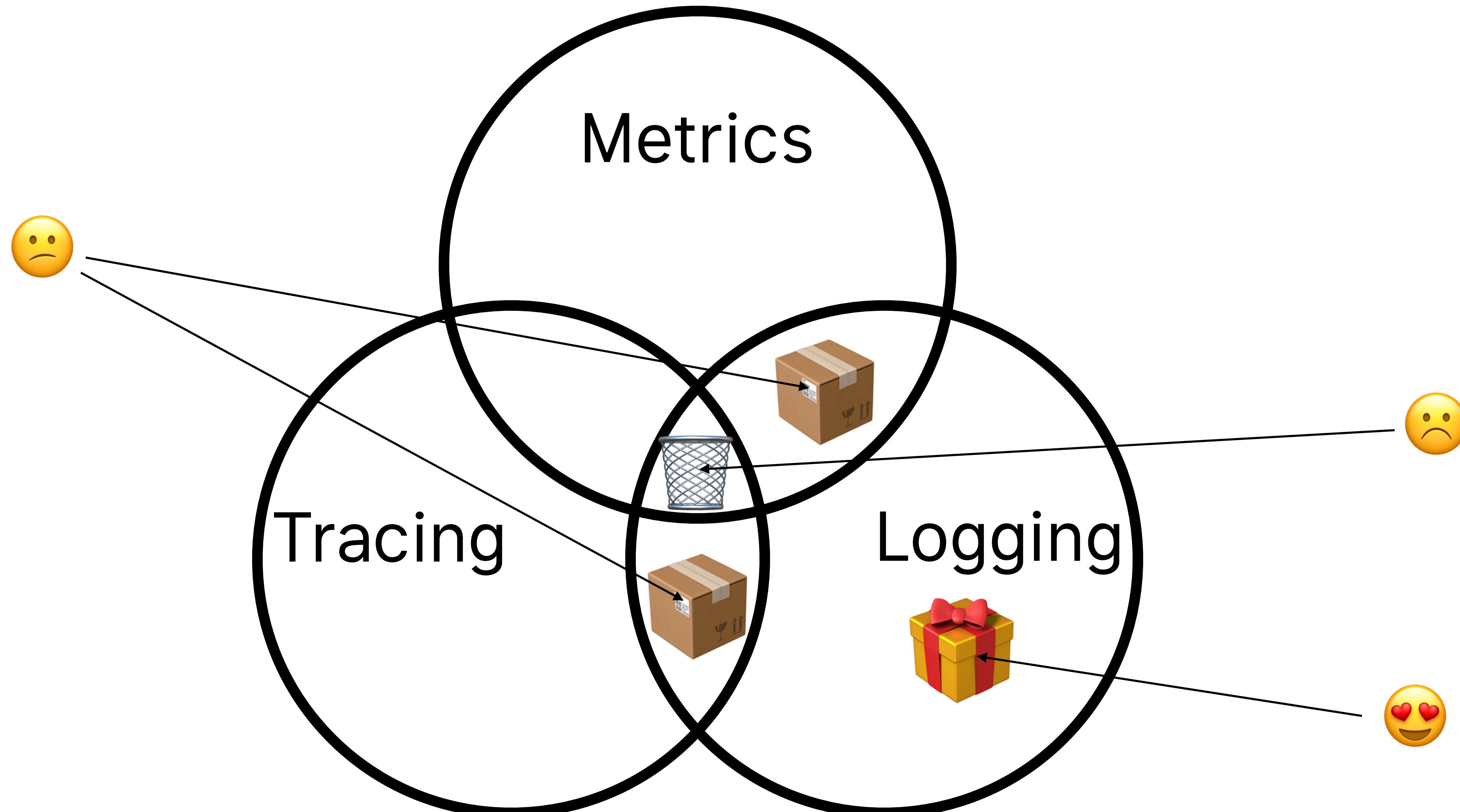


What are we doing here

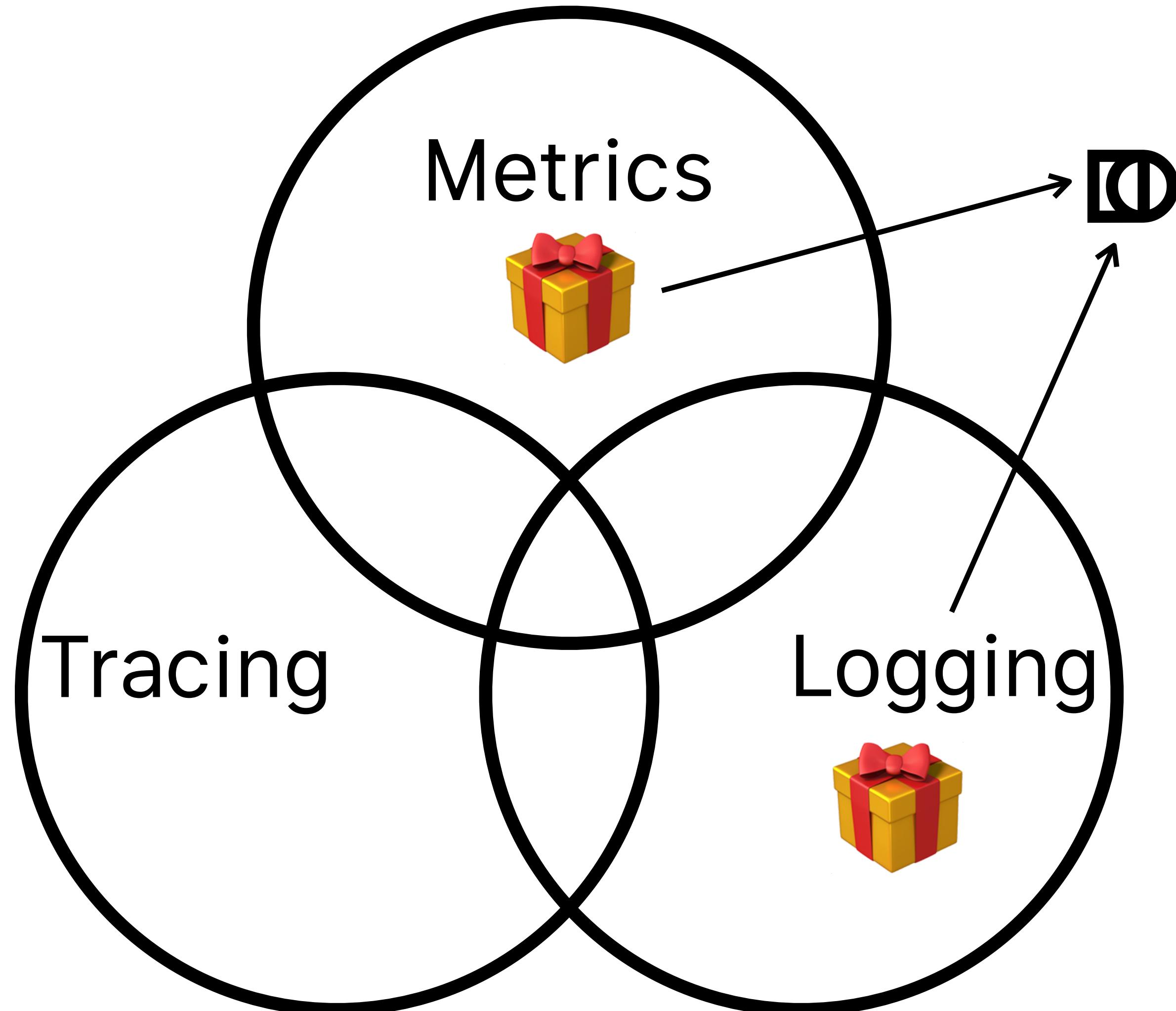
- Observability = mechanisms for understanding
- Given a *what*, logging helps us answer the *why* and *how*
 - *What* = high-level, via metrics and alerts
 - *Why, how* = low-level, via error messages and contextual info



Well-defined scope



Well-defined scope



Actual requirements

- Often: grep for some error in the recent past
- Often: grep for some customer/transaction ID
- Sometimes: grep for a class of things in the last days/weeks
- Rarely: roll-ups, statistical analysis — *this is the domain of metrics!*



Products

- Kafka? Whoa boy
- HBase? Hadoop lol
- Heka? Abandonware
- Ekanite? syslog seems wrong
- Fluentd/Logstash? Only part of the problem
- Splunk/Loggly/Honeycomb/...? Not on-prem, \$/€/£

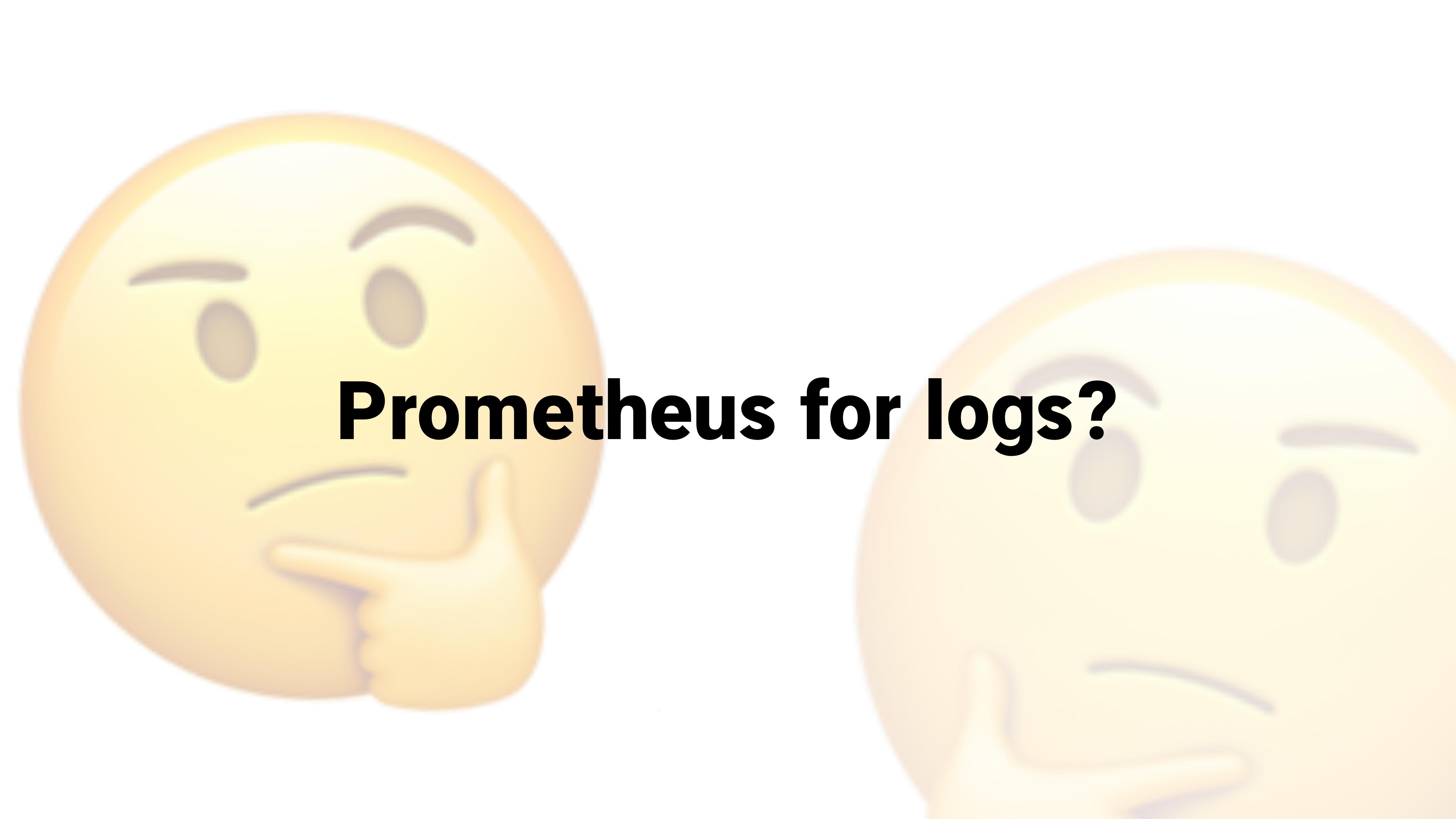


Why is Prometheus good?

IMO

- Run it yourself — open-source and deployable on-prem
- "Cloud-Native" — dynamic, containerized, microservicey workloads
- Easy to operate — local storage, no clustering, pull model
- Complete system — doesn't need separate web UI, TSDB, etc.
- Scales up & out — 90% of use cases satisfied without special effort





Prometheus for logs?

Outline

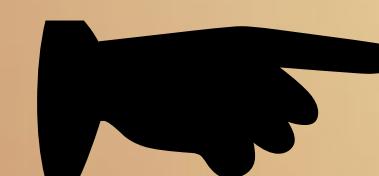
Gremlins of distsys

- 👉 Logging systems
- OK Log design

Outline

Gremlins of distsys

Logging systems



OK Log design



👉 OK Log design
High-level goals
Components
Behaviors

👉 OK Log design
→ *High-level goals*
Components
Behaviors

High-level goals

- Prometheus for logs
 - Run it yourself — open-source and deployable on-prem ✓
 - "Cloud-Native" — dynamic, microservicey workloads ✓
 - Easy to operate — local storage, no clustering, pull model ?
 - Complete system — doesn't need separate web UI, TSDB, etc. ✓
 - Scales up & out — 80% of use cases satisfied w/o special effort ✓



High-level goals

- Principal use case: application logs from typical microservices
 - High-volume, low-QOS
 - Secondary use case: event or audit logs
 - Low-volume, high-QOS
 - Universal log consumer, including from third-party software
 - Can't mandate structure; should treat each record as opaque



Constraining the problem

- Opaque records imply a pure *transport* system
- Leave application concerns to the edges
 - e.g. Contextual annotation can happen prior to ingest
 - e.g. Queryable indices, aggregates modeled as stream consumers
- Only available metadata = timestamp at ingest



Query interface

- Is time-bounded grep sufficient as a query interface?
- Yes
- Maybe
- Let's say yes



A note on scale

- Prometheus has it easy — metrics are aggregatable
- A single server can conceivably handle 10k+ target services
- Core Prometheus opts-out of clustering/distsys — *hell yeah*
- To operate at desired scale, we (unfortunately) need more capacity
- Logging systems are distributed systems, by necessity



"Easy to operate"

- Must not require specialized knowledge or arcana
- << 1x FTE to deploy and maintain
- Embeddable in product team infrastructure — sysadmin not required
- No byzantine or pathological failure modes!!!
 - Kill processes, [delete files,] start again = must work



"Easy to operate"

- As few knobs as possible — sensible defaults
- Automatic peer discovery and cluster management
- Adding nodes = adding capacity, ideally no rebalancing/reindexing/...
- Losing nodes = decrease capacity, ideally no other side effects
- ⚡ Minimize transactions and other coördination



Coördination-free

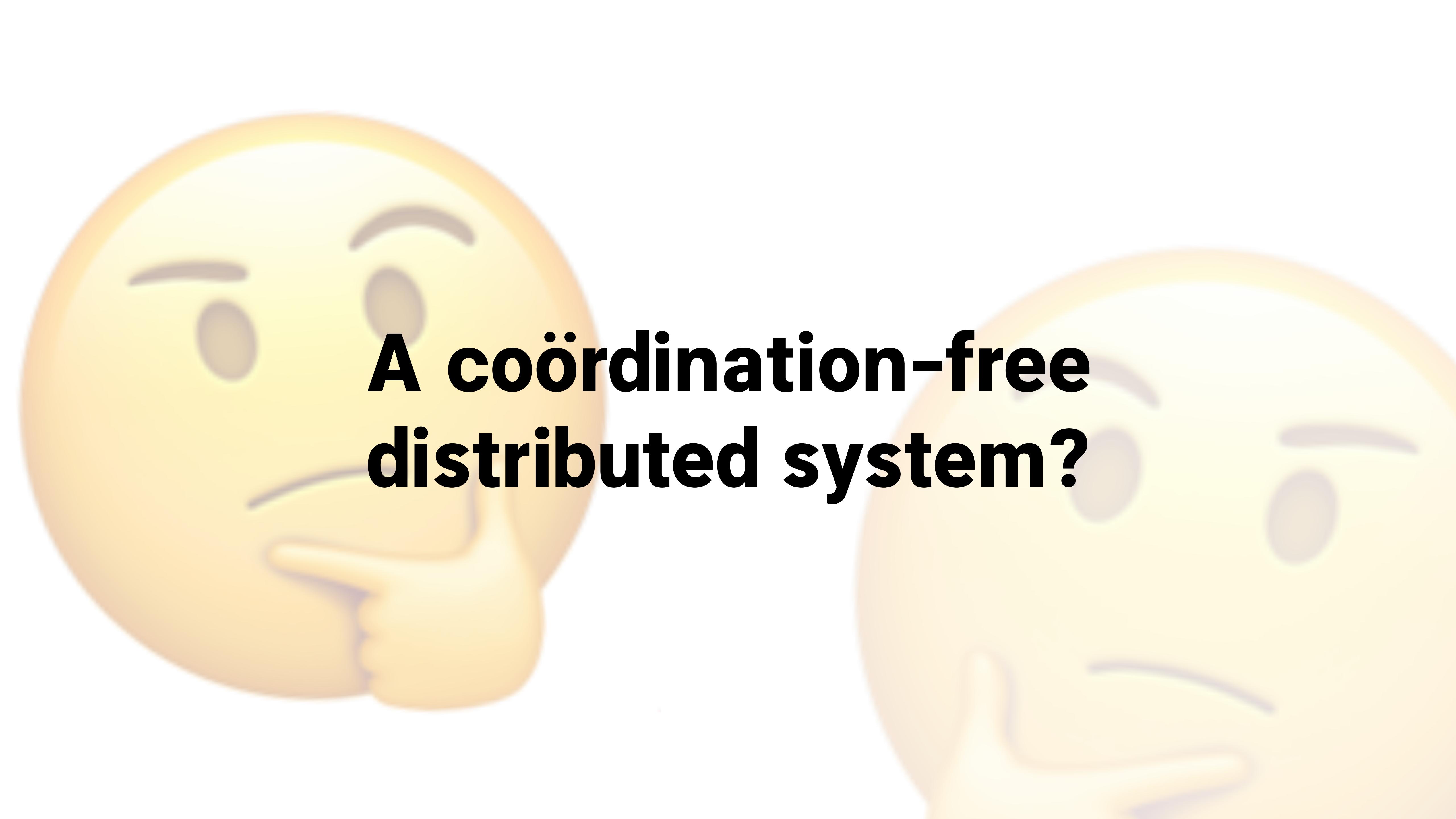
- Coördination: lifting from local truth (*Phenomena*) to global (*Noumena*)
- Expensive, error-prone, often literally impossible (CAP)
- "The best way to solve a problem is to opt-out of it" —*A smart*
- Set aside the fiction of a global truth!
- Coördination-free = phenomenological reduction



In practice

- No master election
- No knowledge of allocations
- No shard table
- No distributed index
- No vnodes, no ring



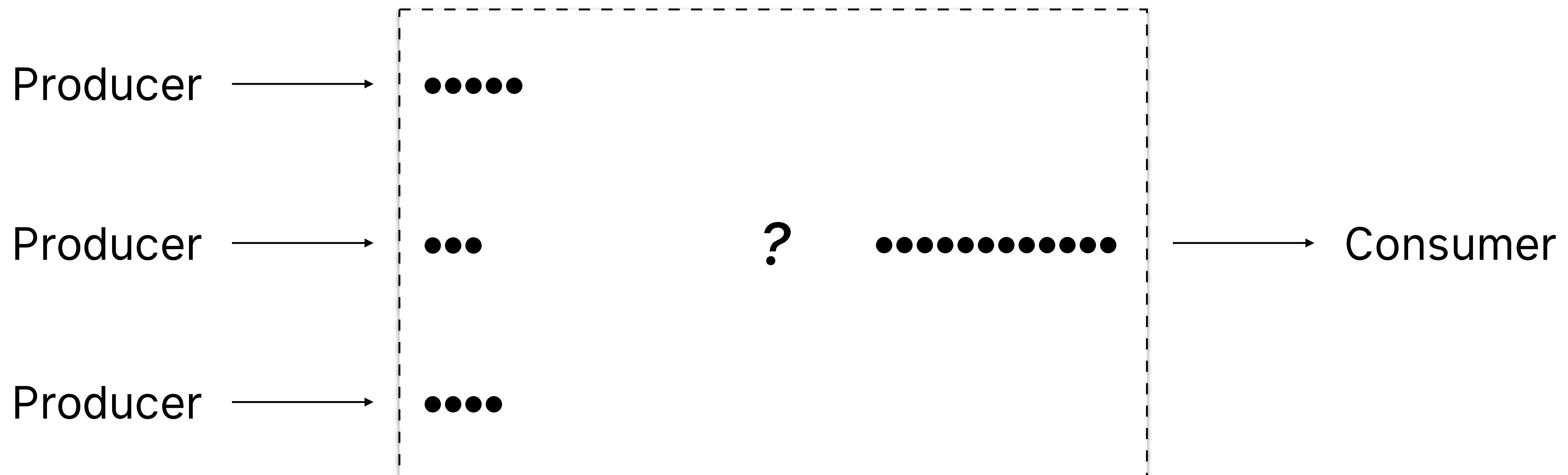
A large, semi-transparent watermark of two yellow cartoon dogs is centered behind the text. The dog on the left has its head tilted back, showing its tongue and teeth. The dog on the right has its head tilted forward, looking towards the first dog. They appear to be communicating or playing together.

A coördination-free distributed system?

👉 OK Log design
→ *High-level goals*
Components
Behaviors

👉 OK Log design
High-level goals
→ Components
Behaviors

The importance of writes



Ingestion

- Different performance requirements for ingestion vs. querying
 - Ingest must take precedence
- Makes sense to isolate those workloads to different machines
 - They could be colocated for small or trial installations
 - Ideally ingestion is limited by hardware



ULID

0	1	2	3
0 1 2 3 4 5 6 7 8 9 0	1 2 3 4 5 6 7 8 9 0	1 2 3 4 5 6 7 8 9 0 1	
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+			
32_bit_uint_time_high			
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+			
16_bit_uint_time_low		16_bit_uint_random	
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+			
32_bit_uint_random			
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+			
32_bit_uint_random			
+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+-----+			



ULID

Timestamp
10 chars, 48 bits
ms precision

01AN4Z07BY79KA1307SR9X4MV3

Entropy
16 chars, 80 bits
1.208 septillion per ms

Good enough

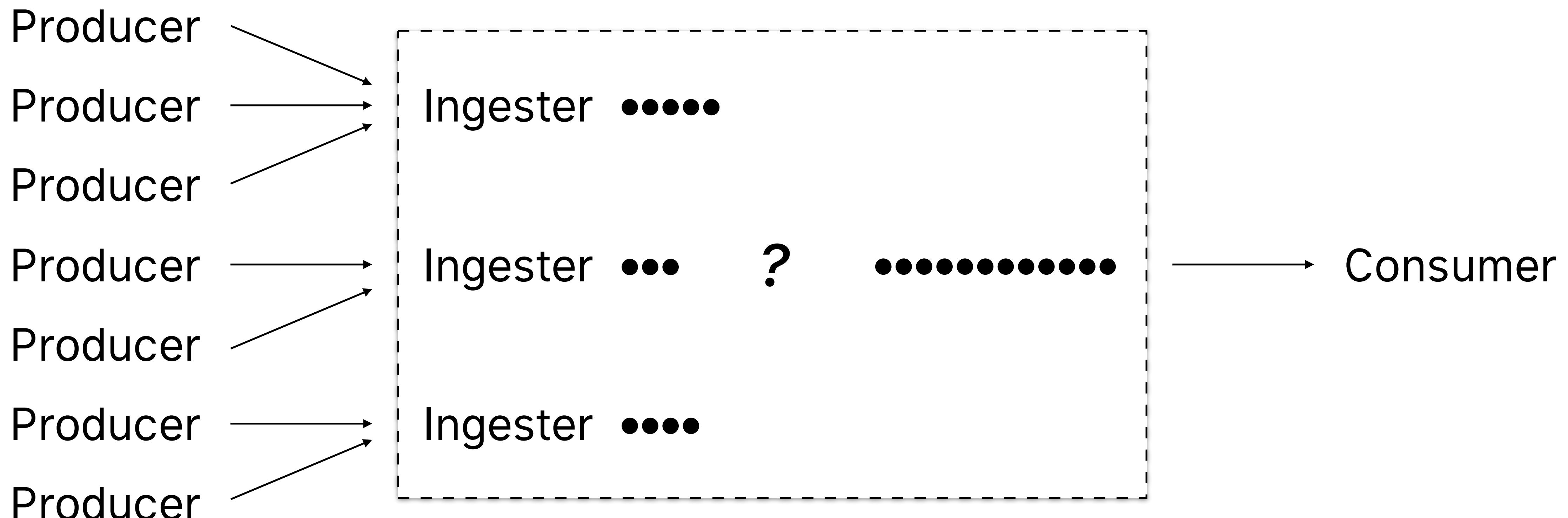


Durability modes

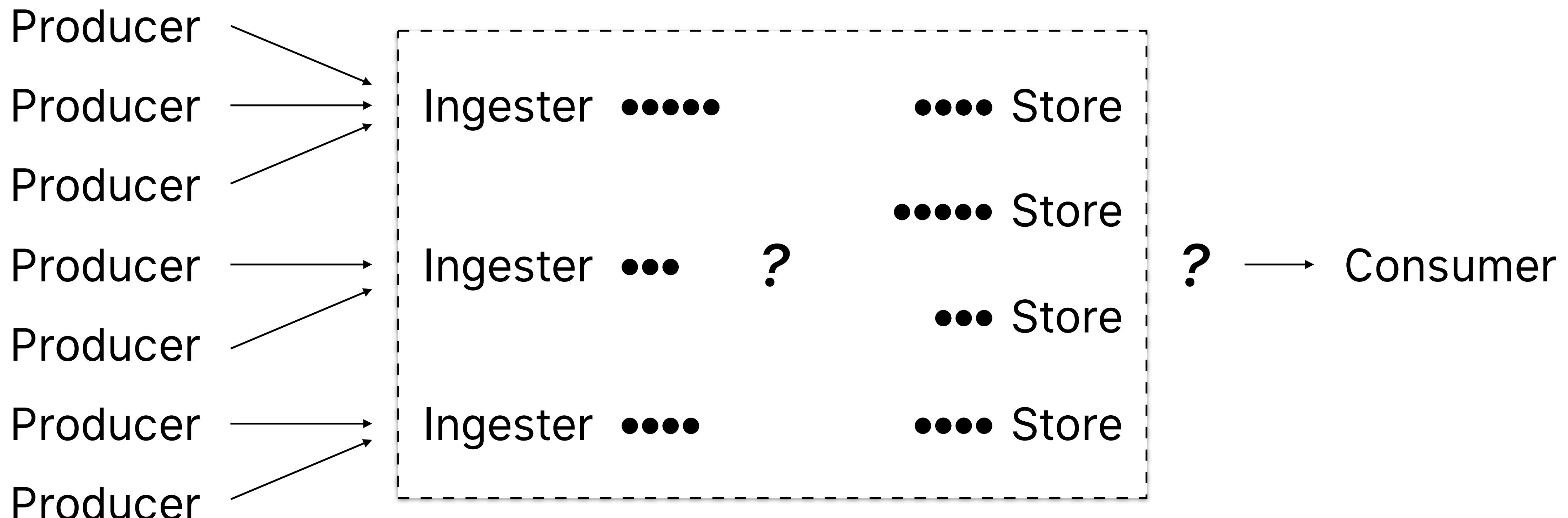
- *Fast* mode — typical for application logs
 - Write to file descriptor without explicit disk sync
- *Durable* mode — typical for event logs
 - Regular syncs to disk
- *Bulk* mode — typical for compliance/audit logs
 - Ingest only acks when entire segment has been replicated



Component model



Component model



👉 OK Log design

High-level goals

→ Components

Behaviors

👉 OK Log design

High-level goals

Components

→ Behaviors

Ingest

Fast writes



Ingest

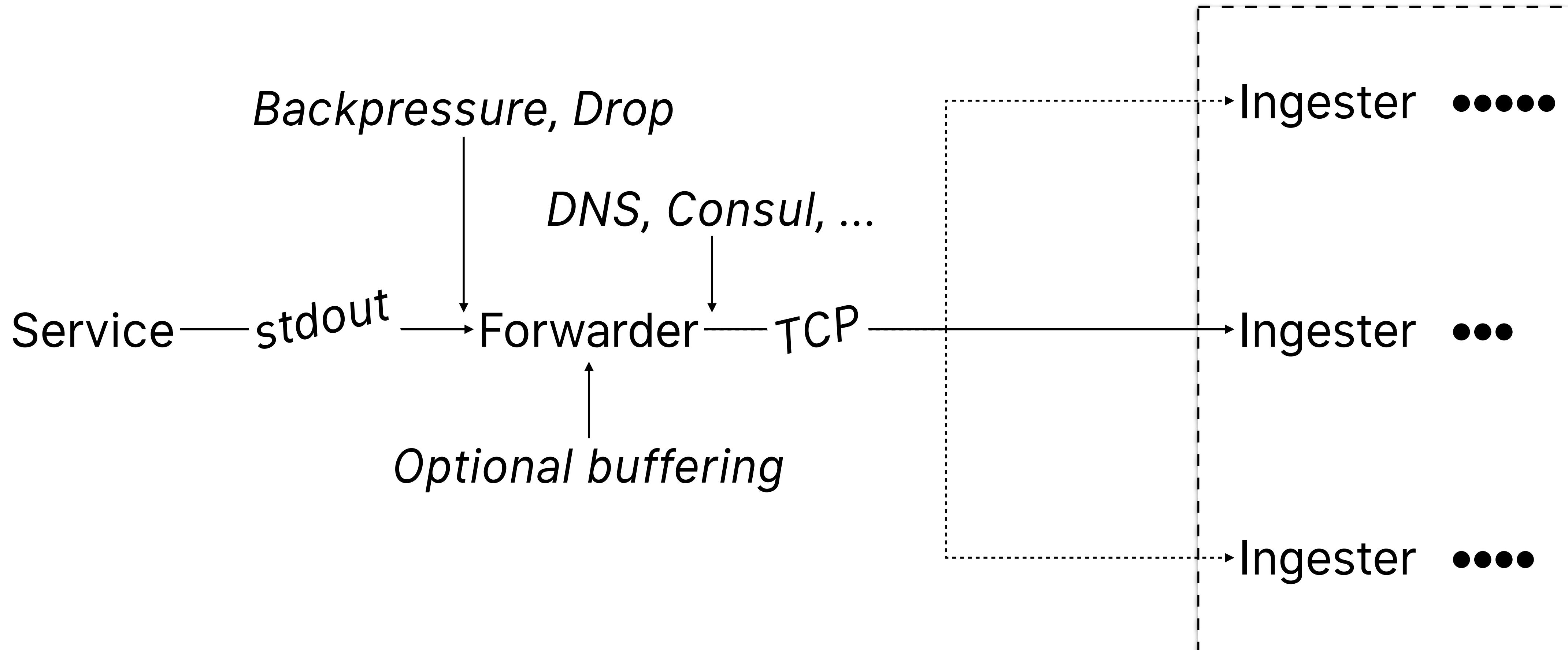
Producer → Ingestor •••

Ingestor ••••

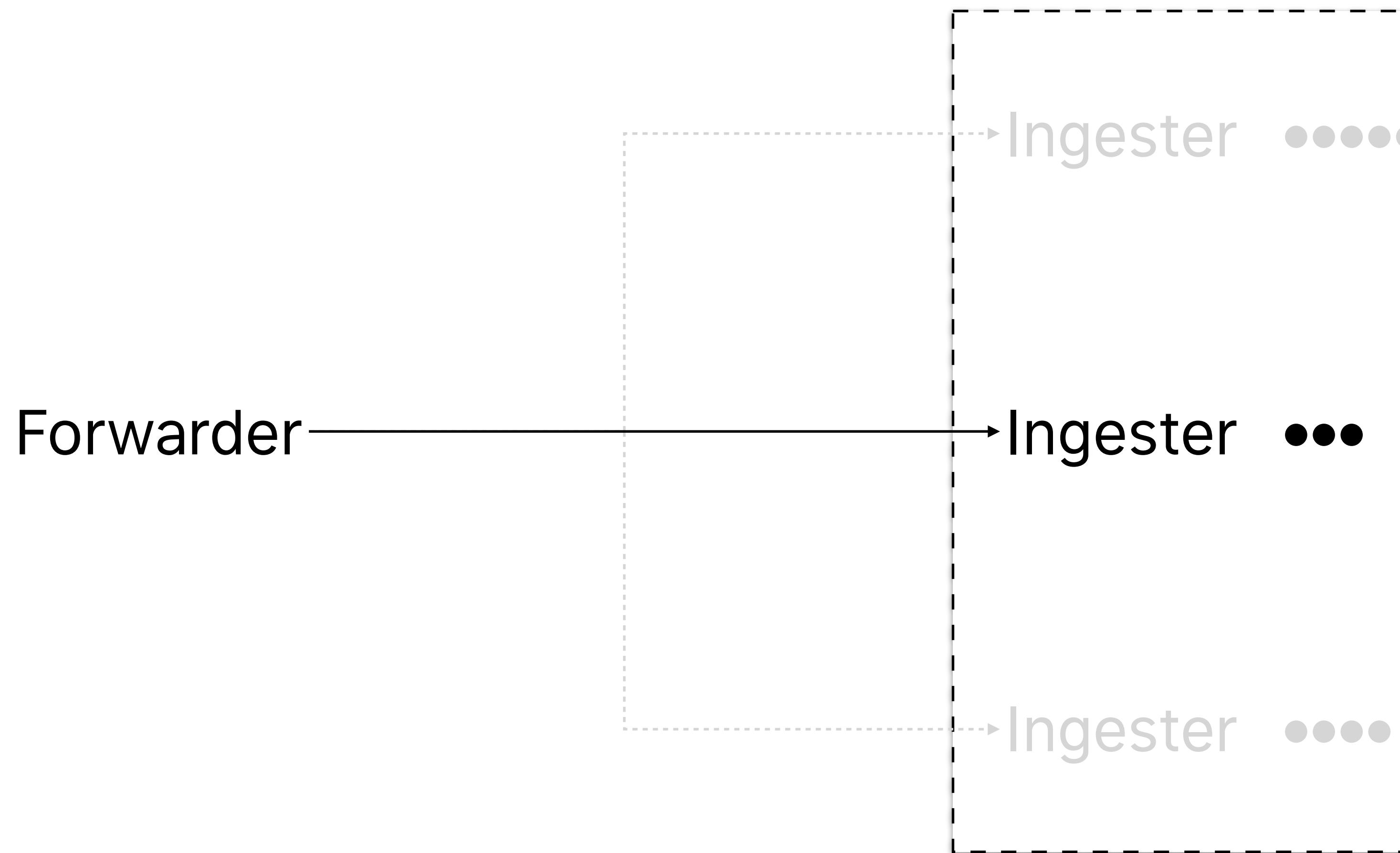
Ingestor •••

Ingestor ••••

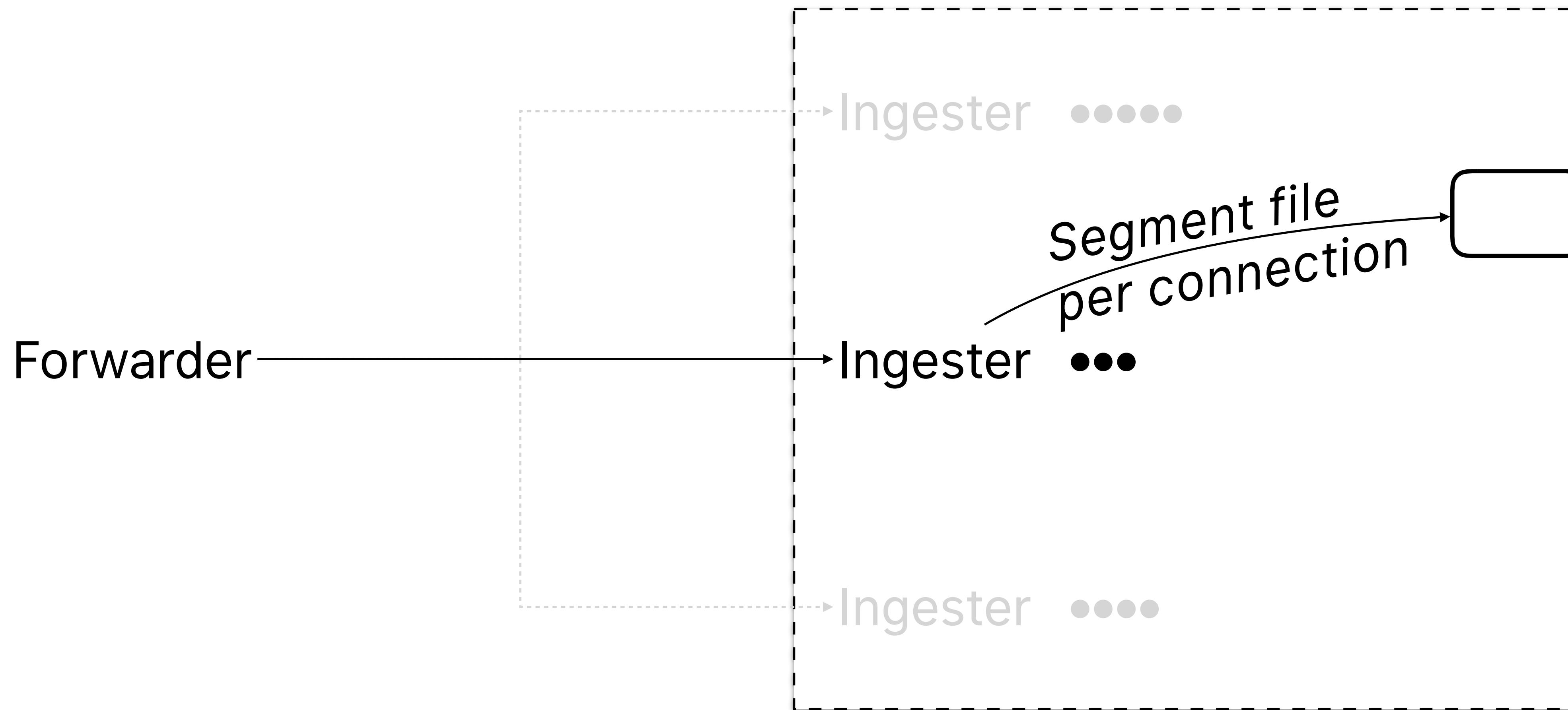
Ingest



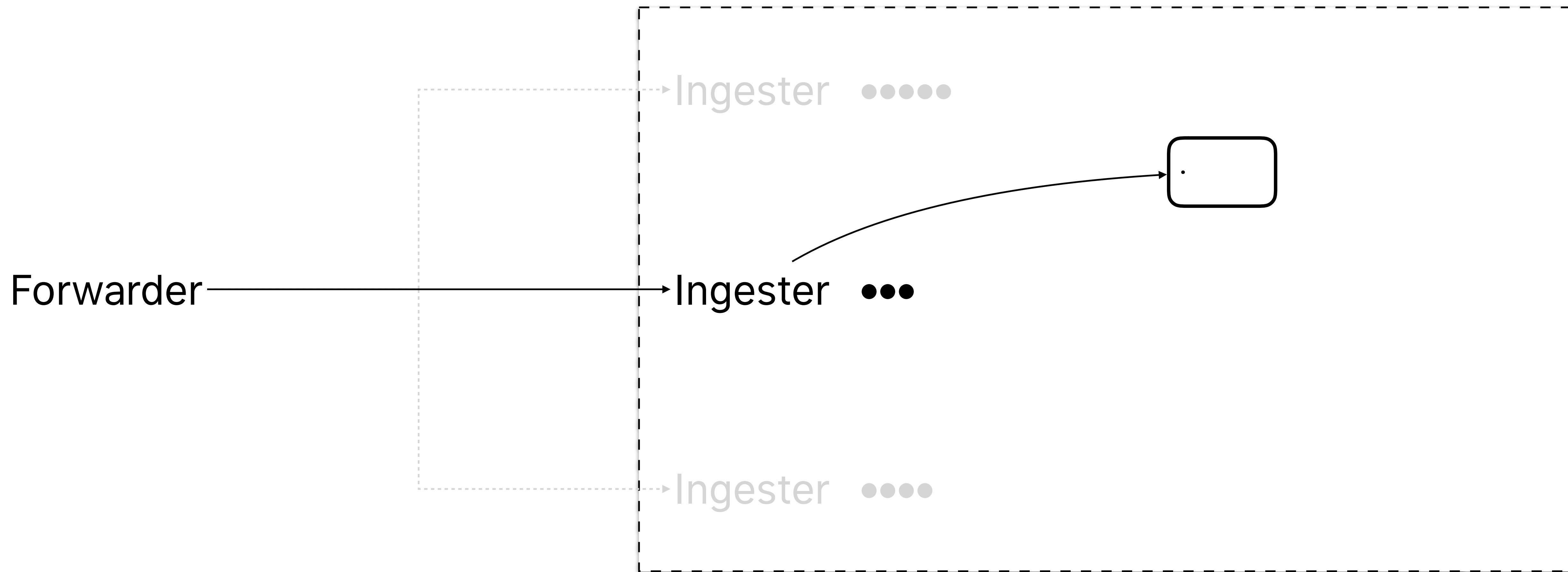
Ingest



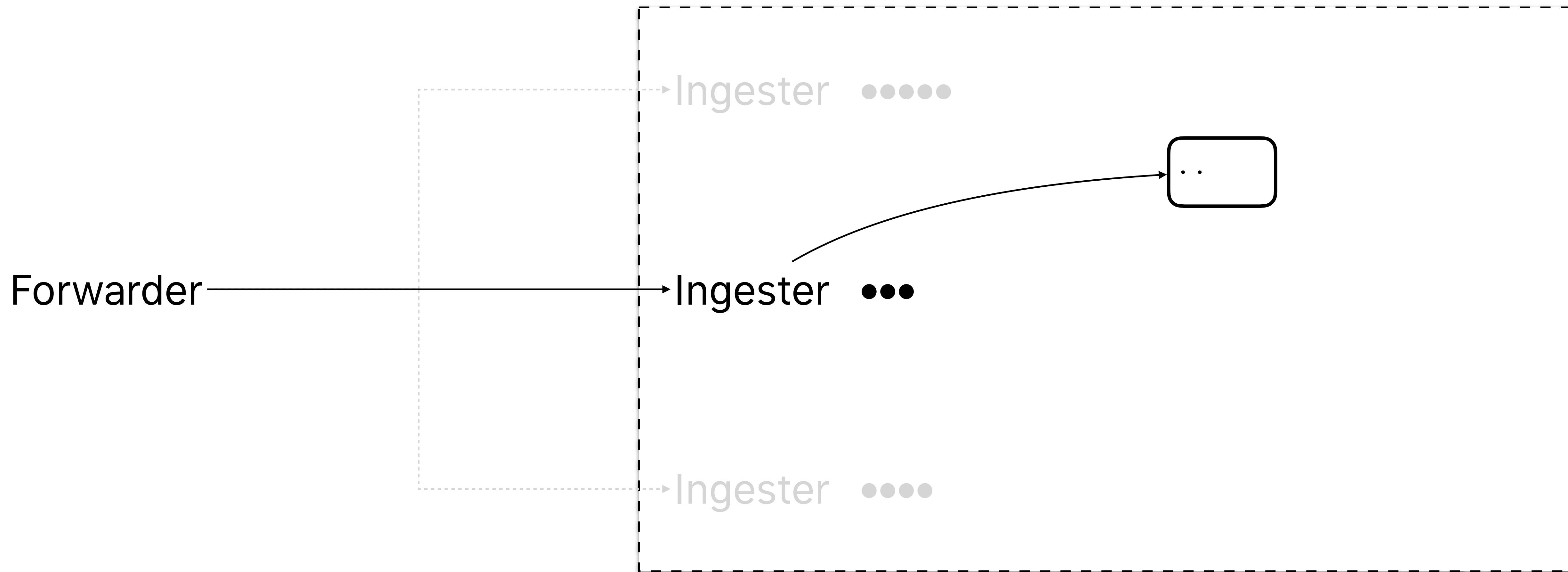
Ingest



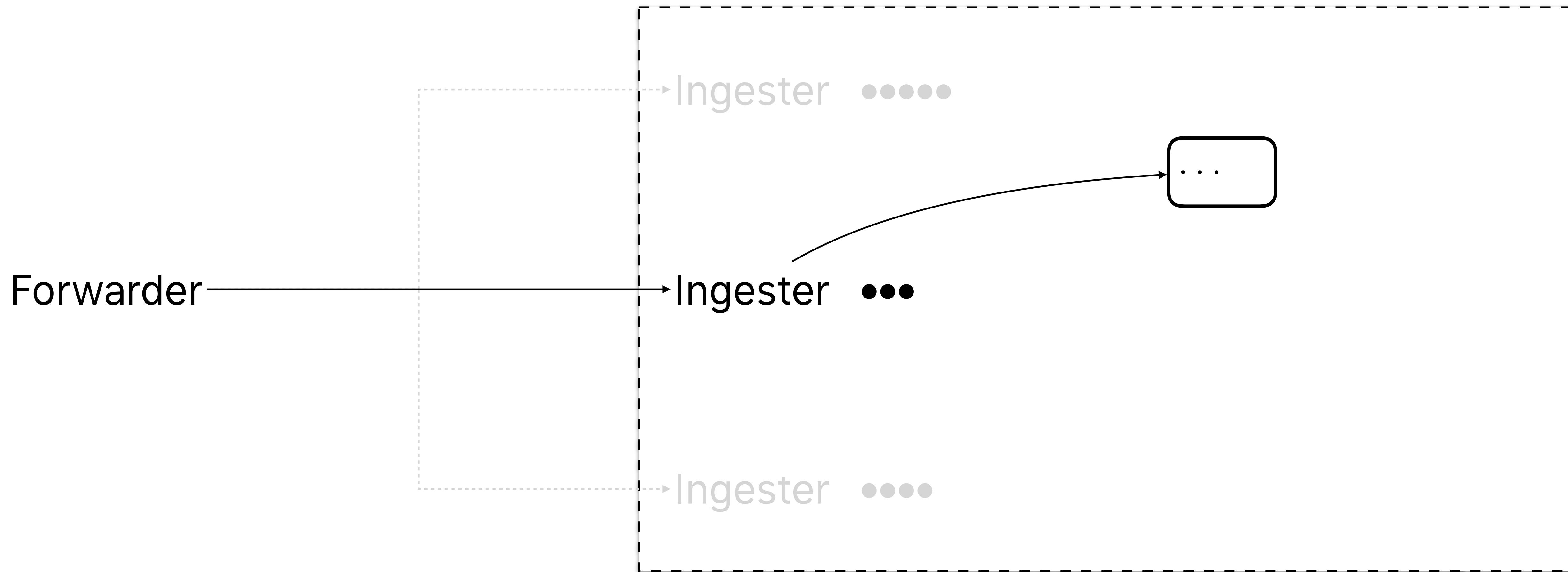
Ingest



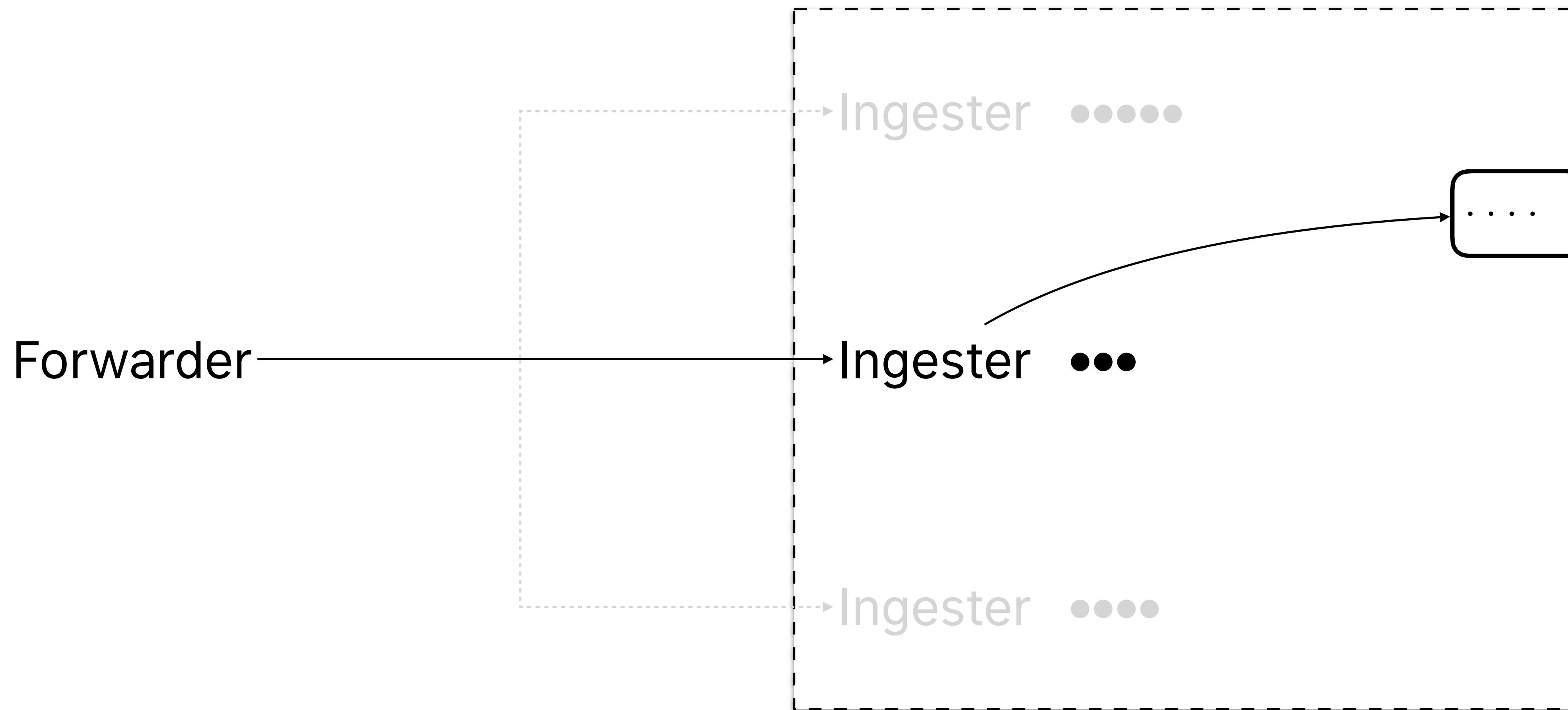
Ingest



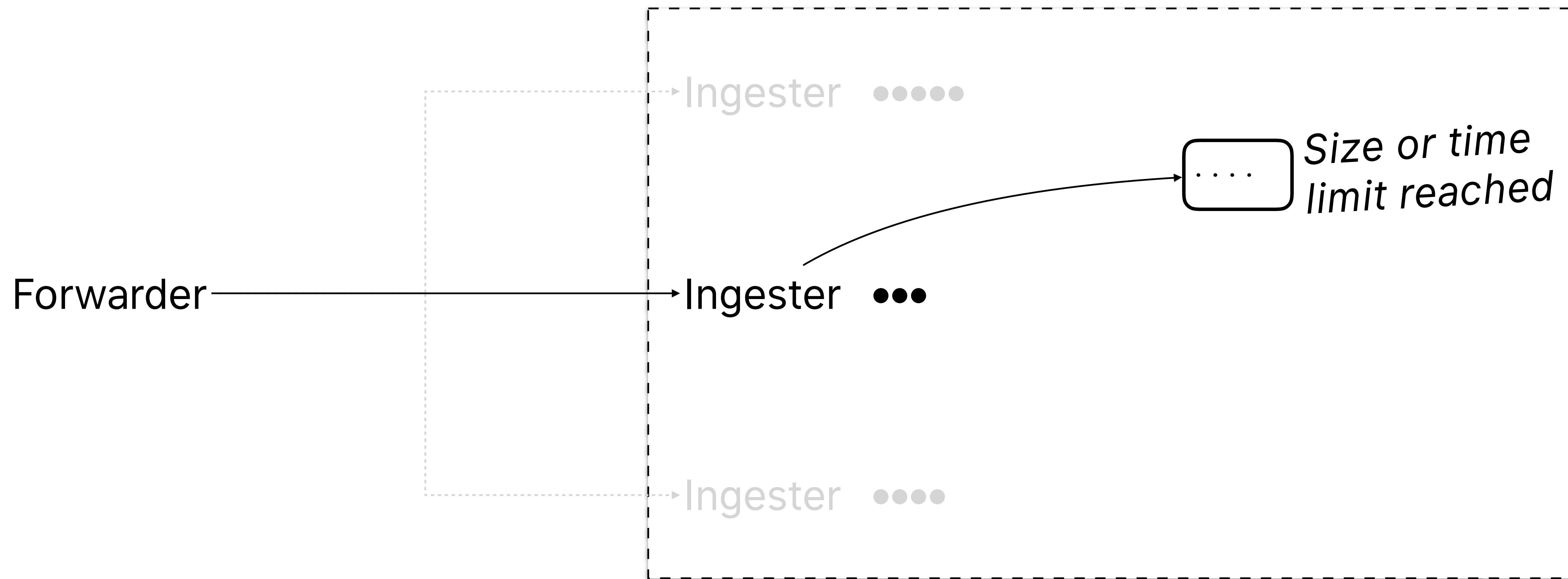
Ingest



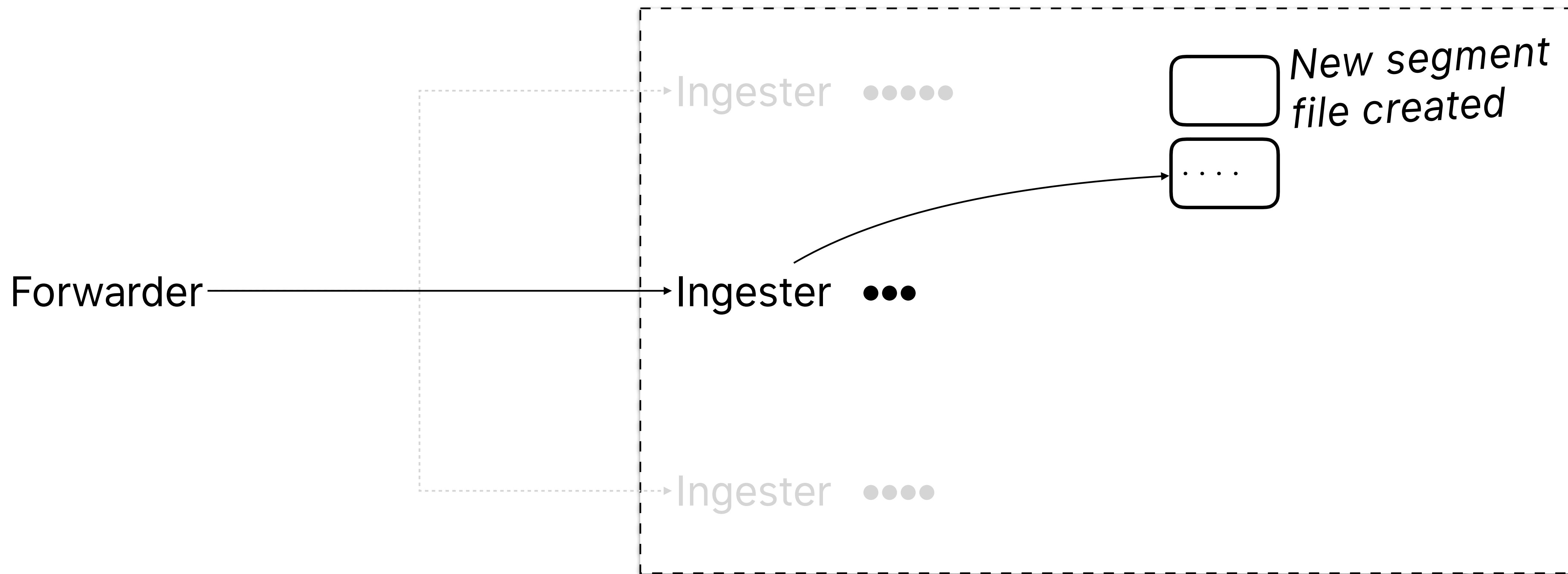
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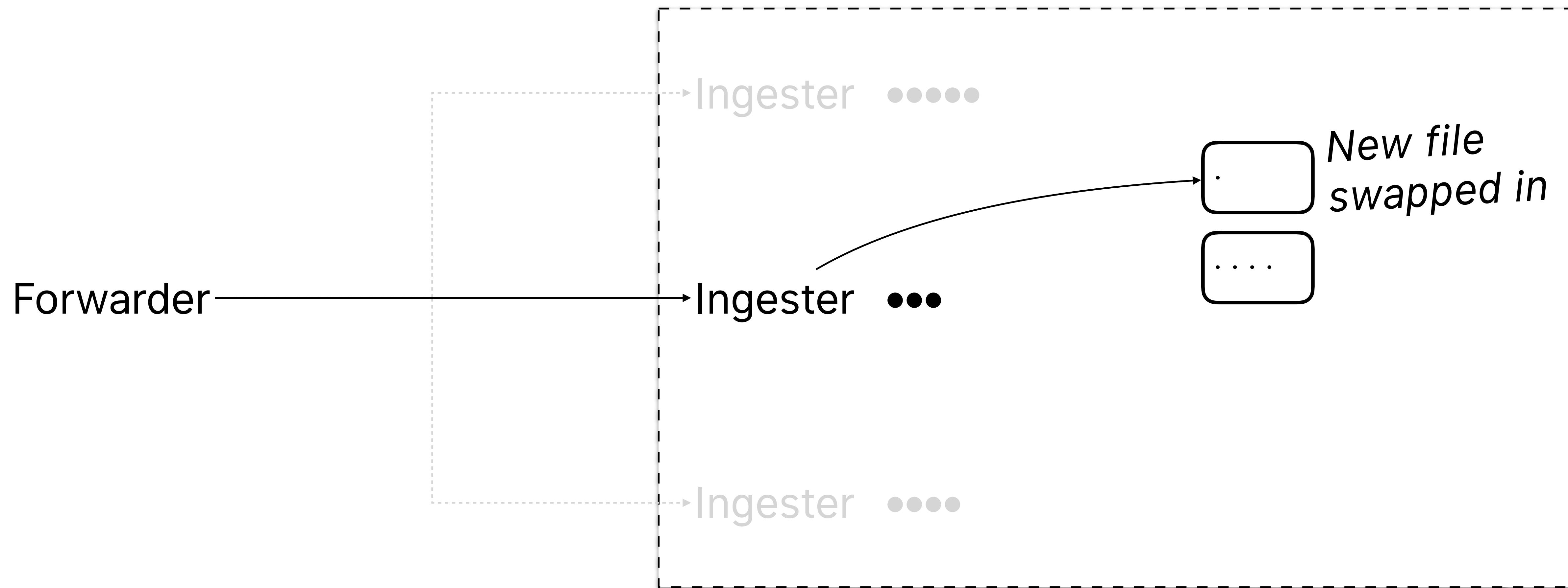
Ingest



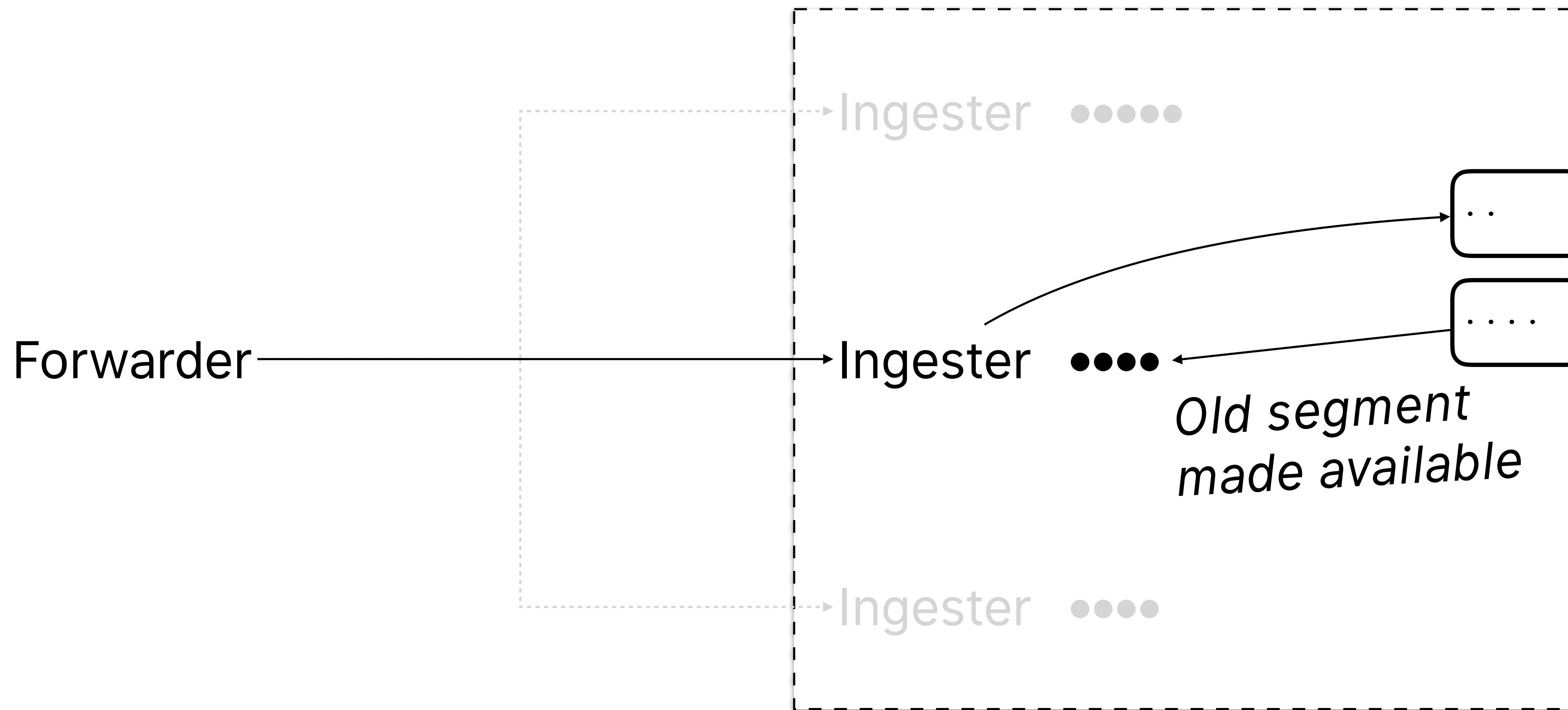
Ingest



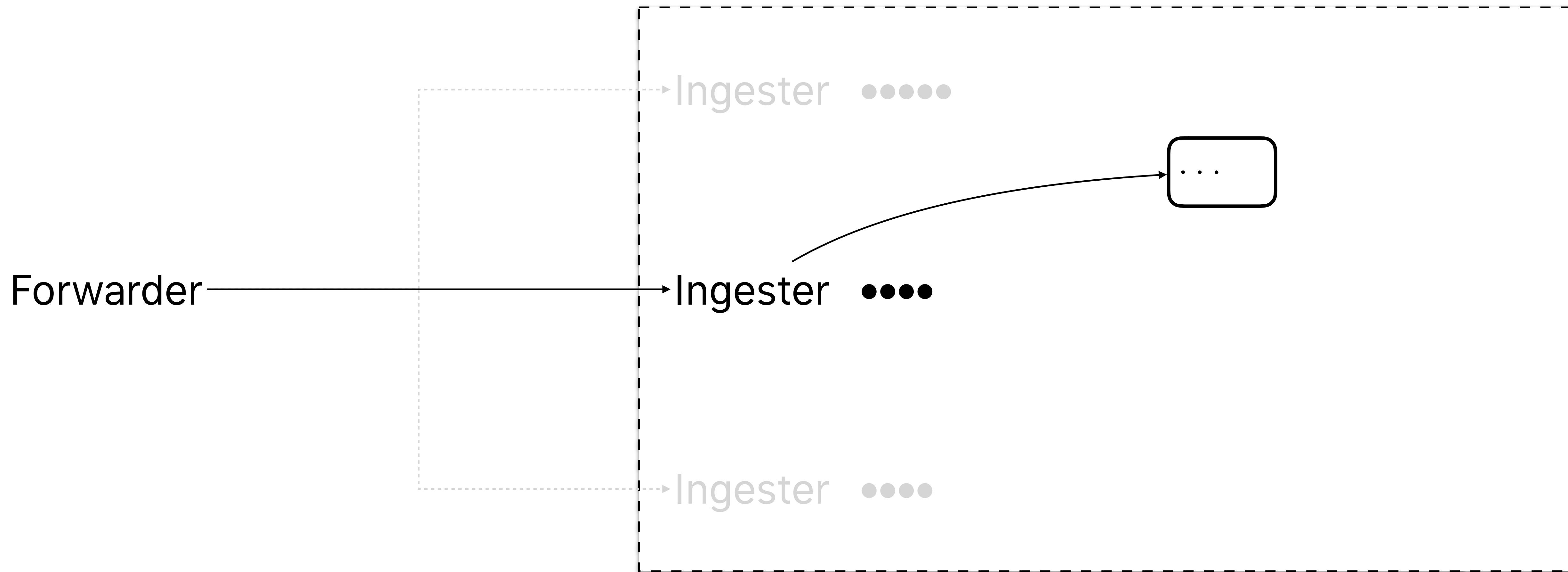
Ingest



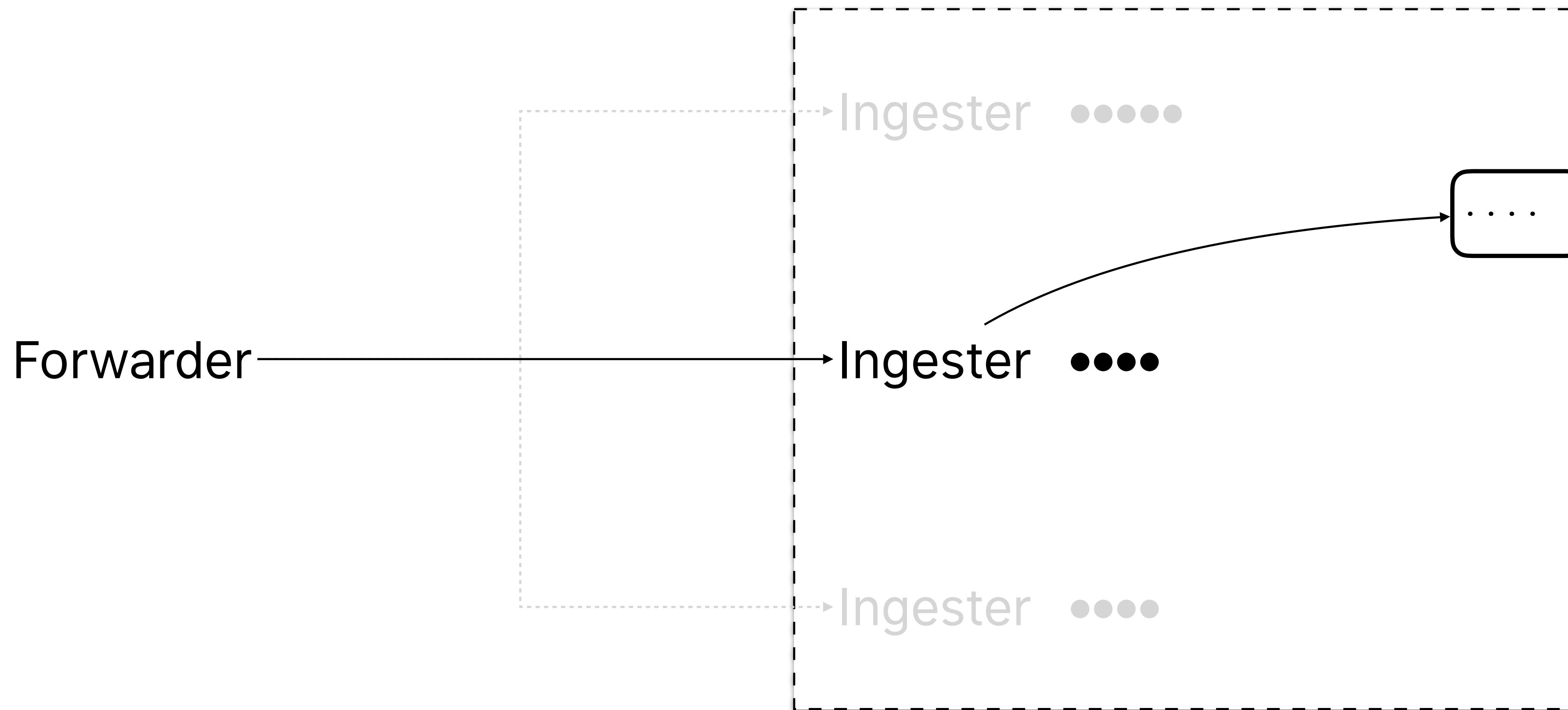
Ingest



Ingest



Ingest



Priorities

- Writes are the hot path, they should succeed
- Everything is optimized to get writes to disk
- Segment size limits affect liveness in high-throughput environments
- Segment time limits affect liveness in low-throughput environments
- Reasonable defaults of 16MB and 1s respectively

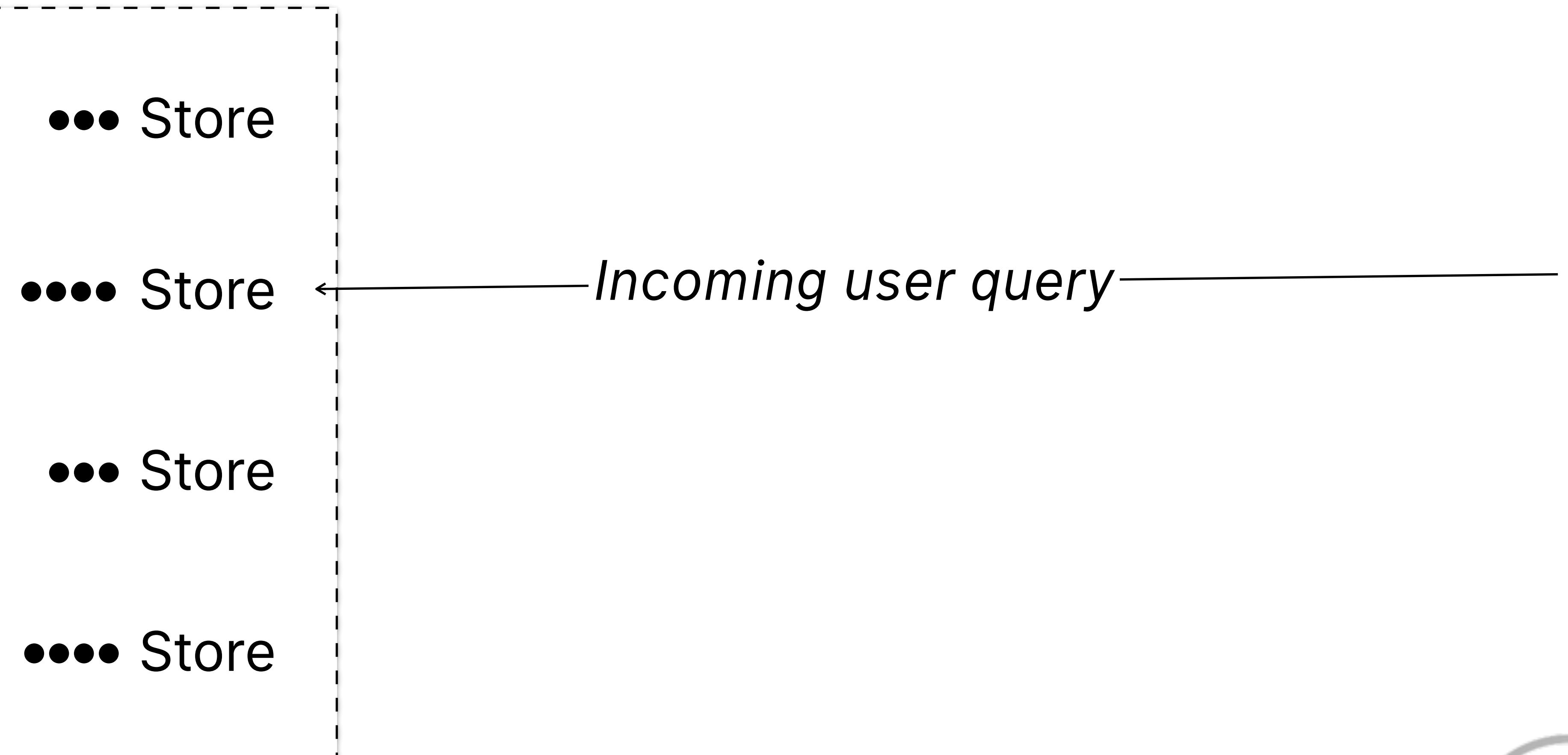


Query

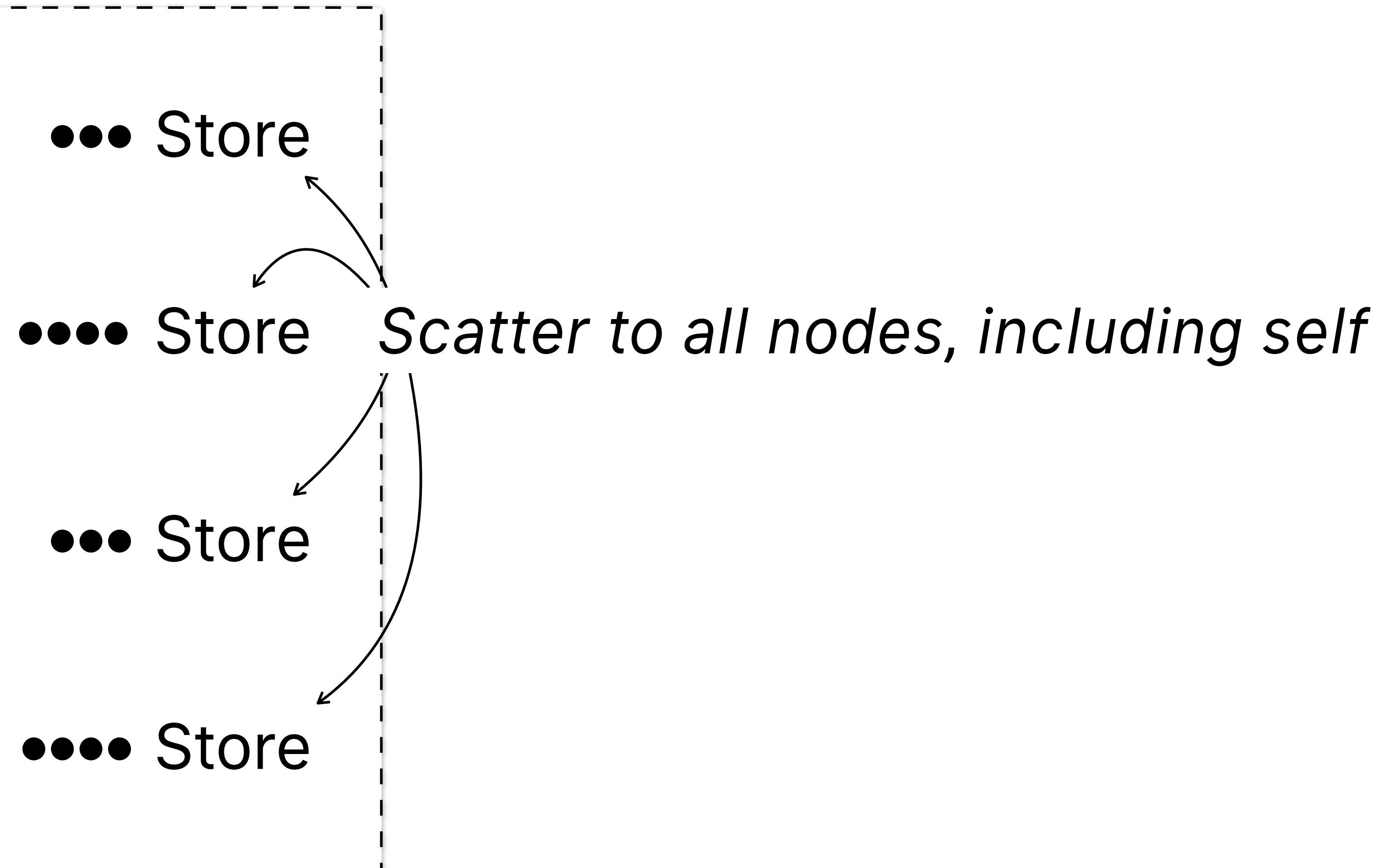
Classic scatter/gather



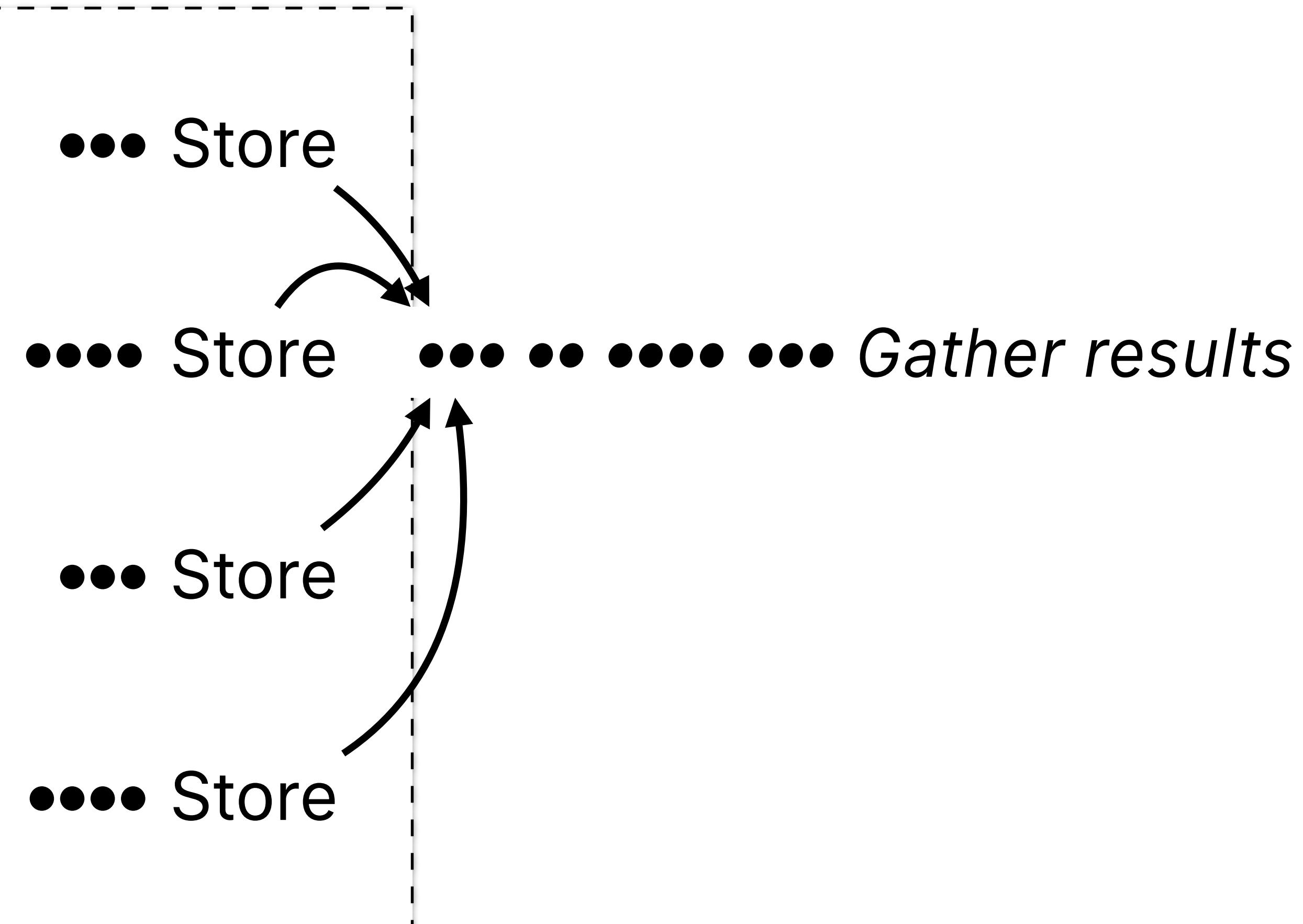
Query



Query



Query



Query

••• Store

•••• Store

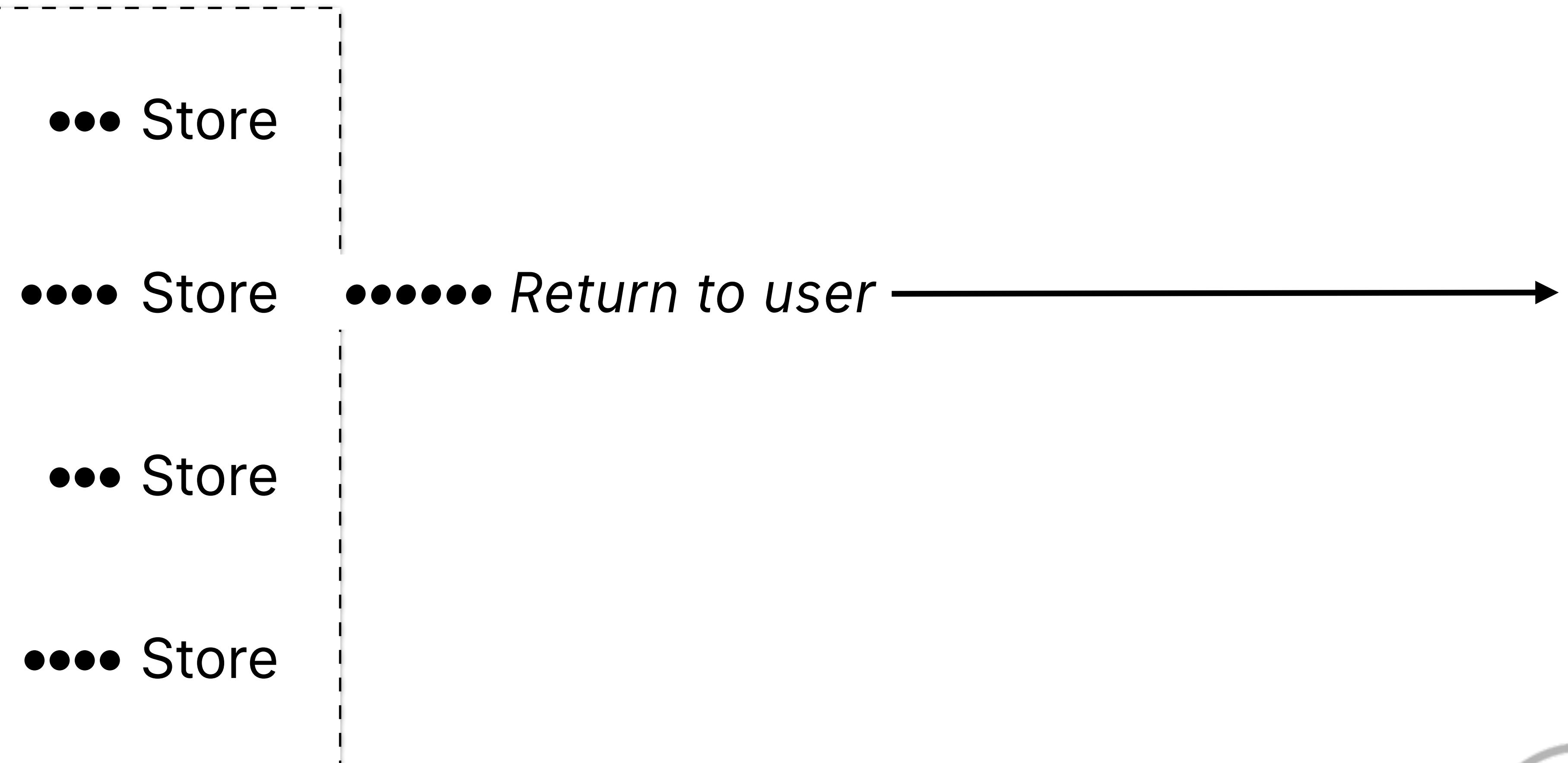
••• Store

•••• Store

••••• *Merge and deduplicate*



Query



Tradeoffs

- True: queries are only as fast as your slowest machine
- Need aggressive timeouts to maintain liveness
- Still, we can tolerate partial node failure without compromising results
- And more substantial node failure degrades gracefully
- Since we get duplicates, we can also do basic read-repair

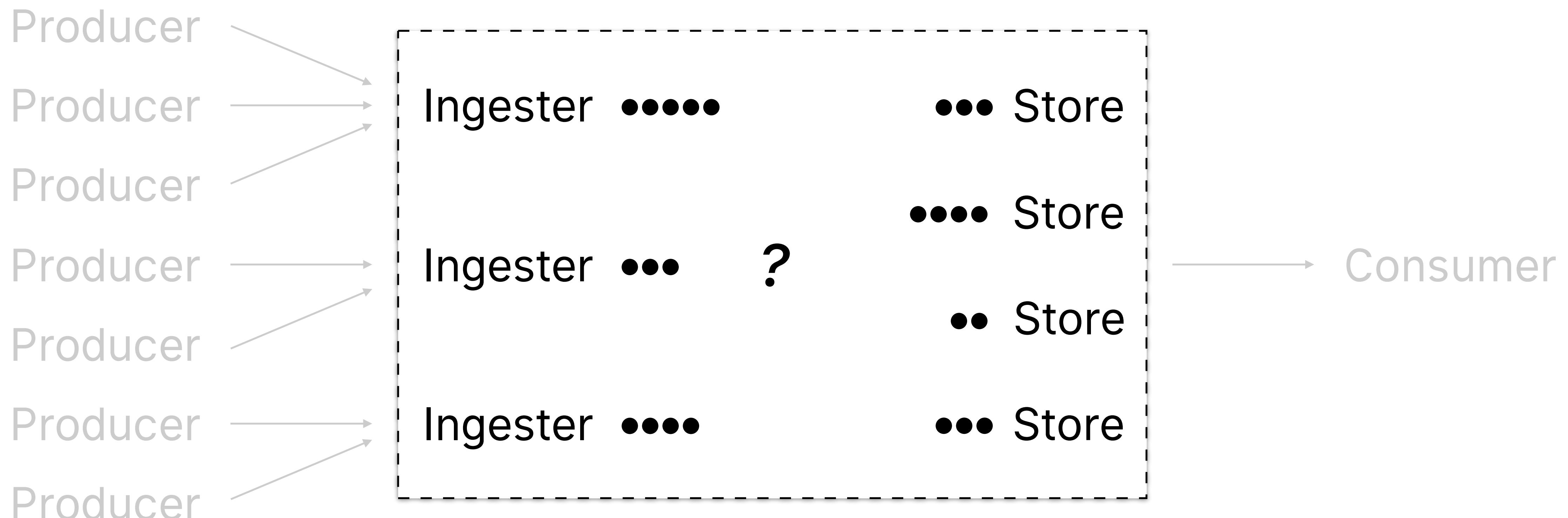


Replication

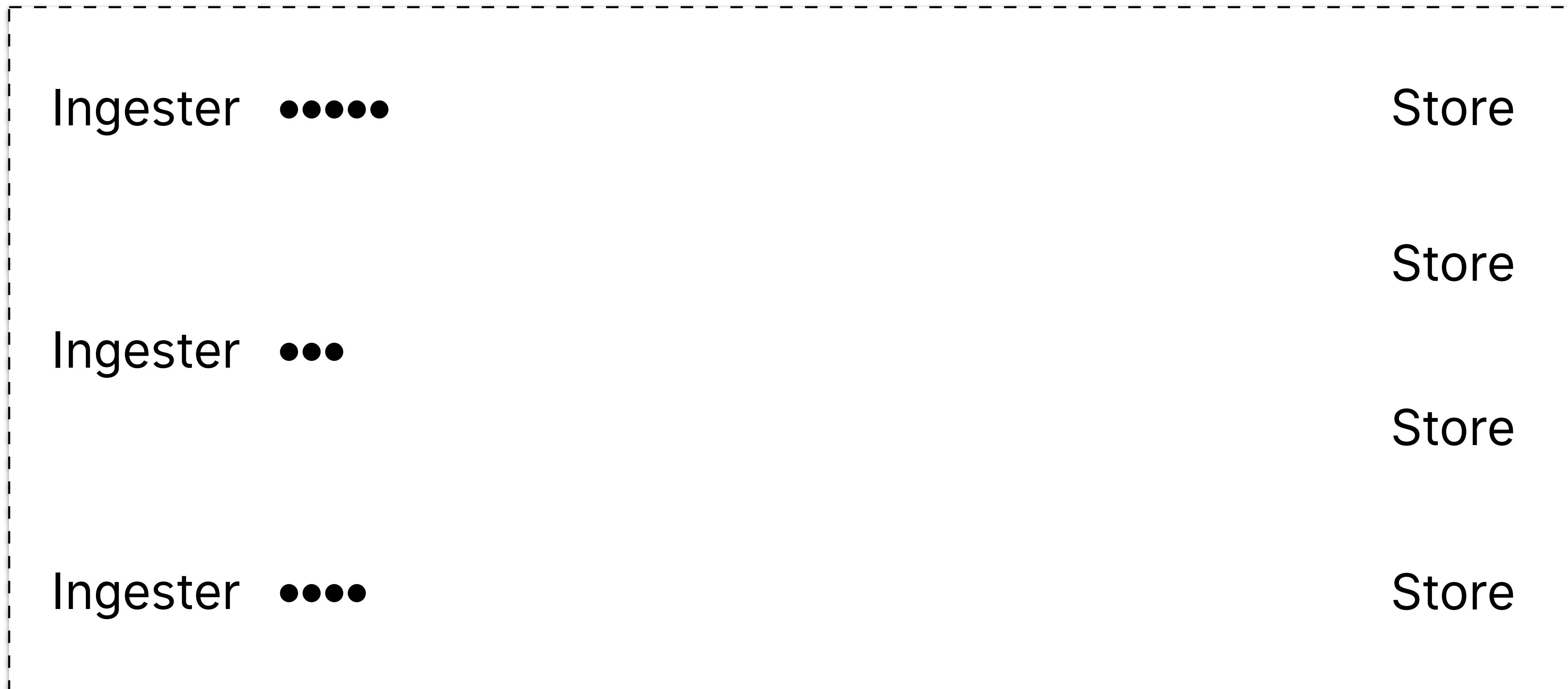
Inspired by Prometheus



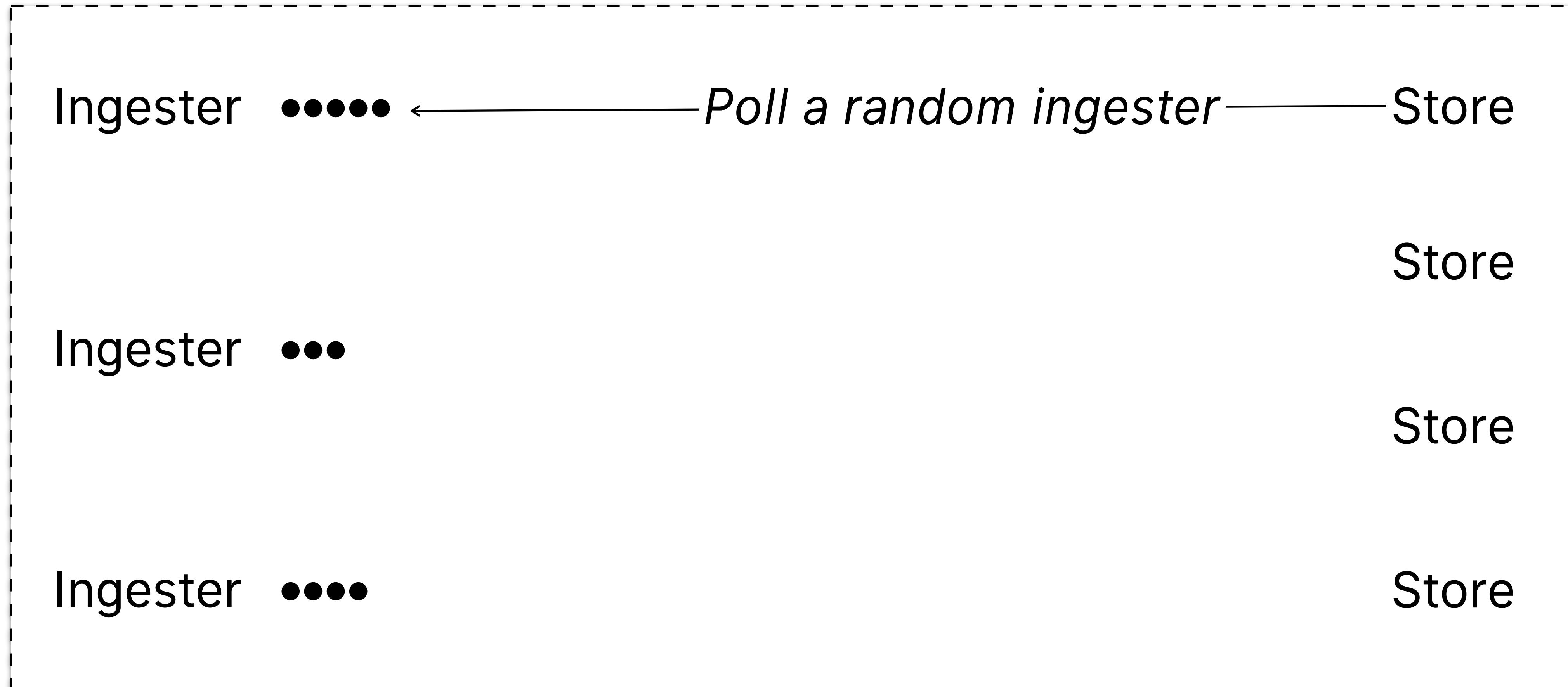
Replication



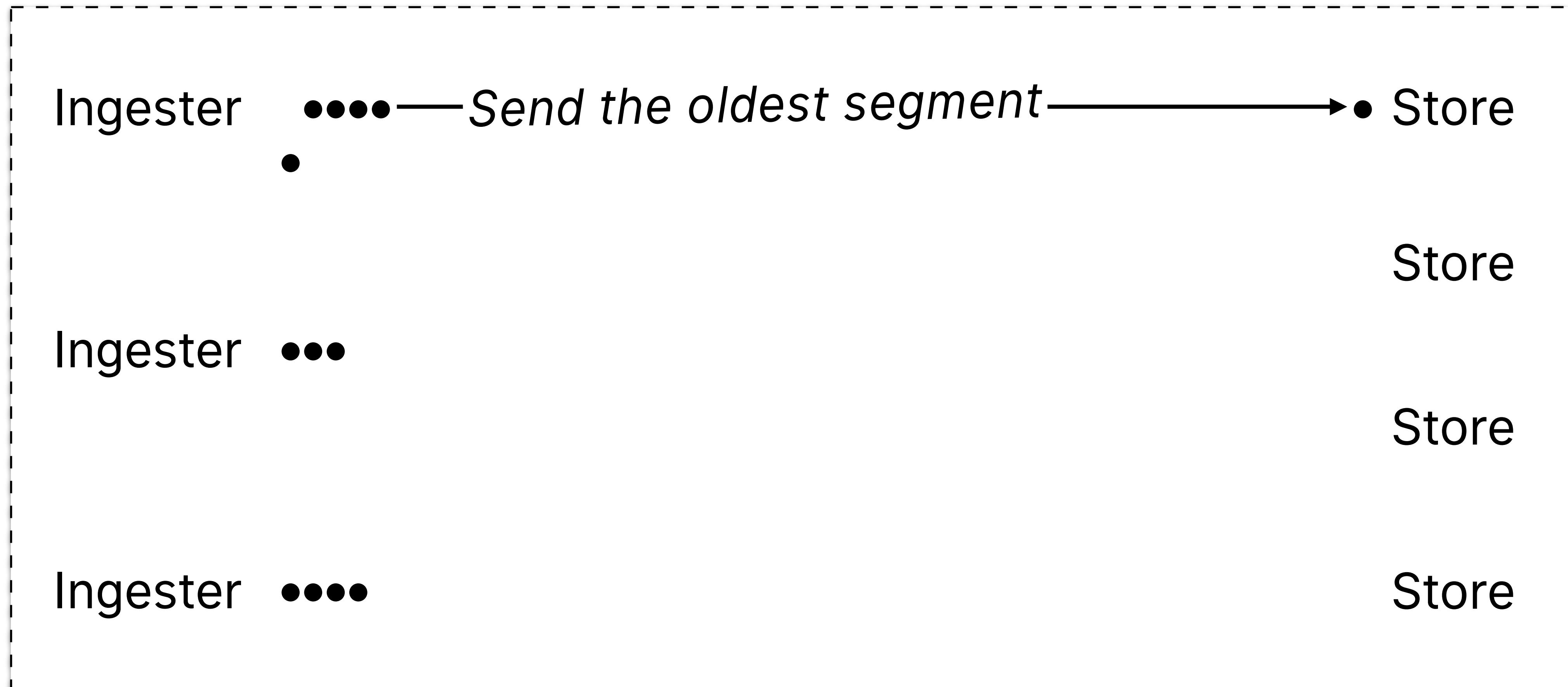
Replication



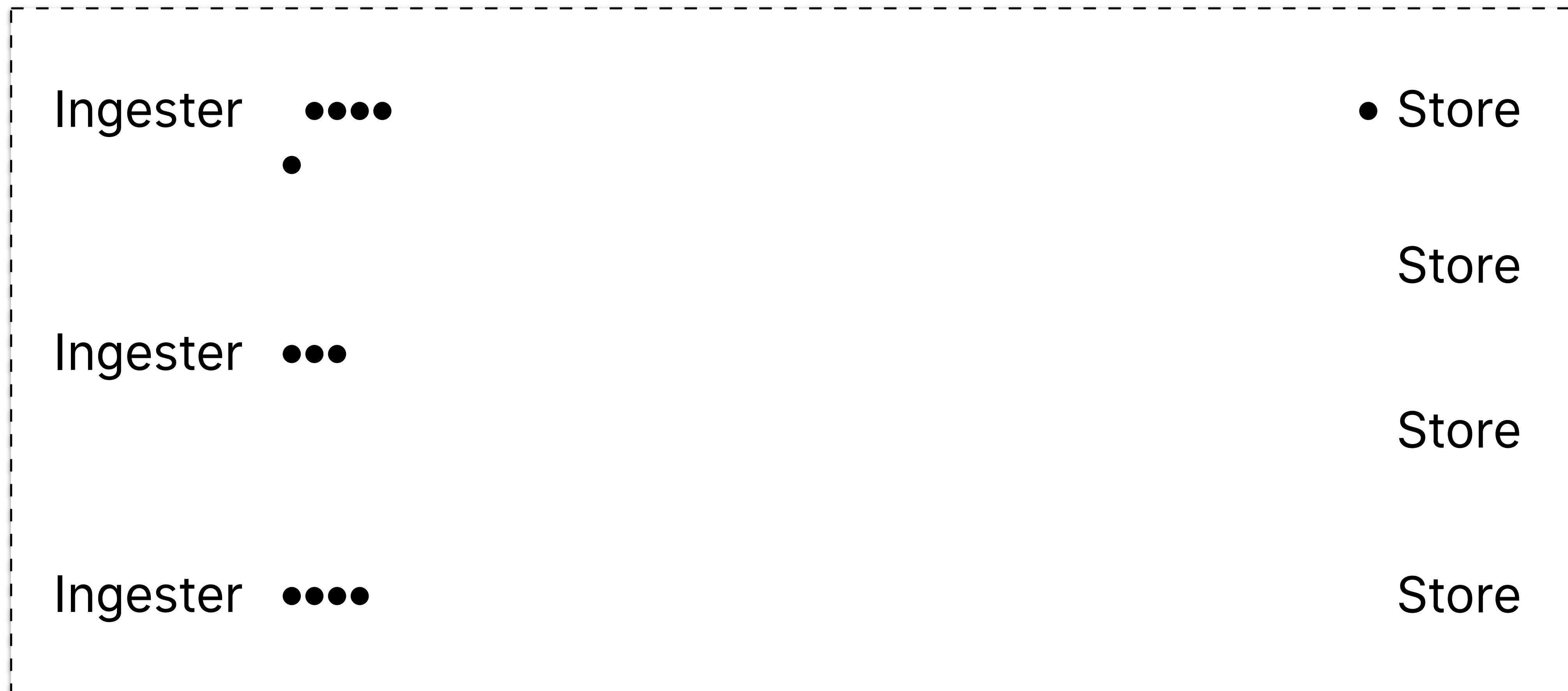
Replication



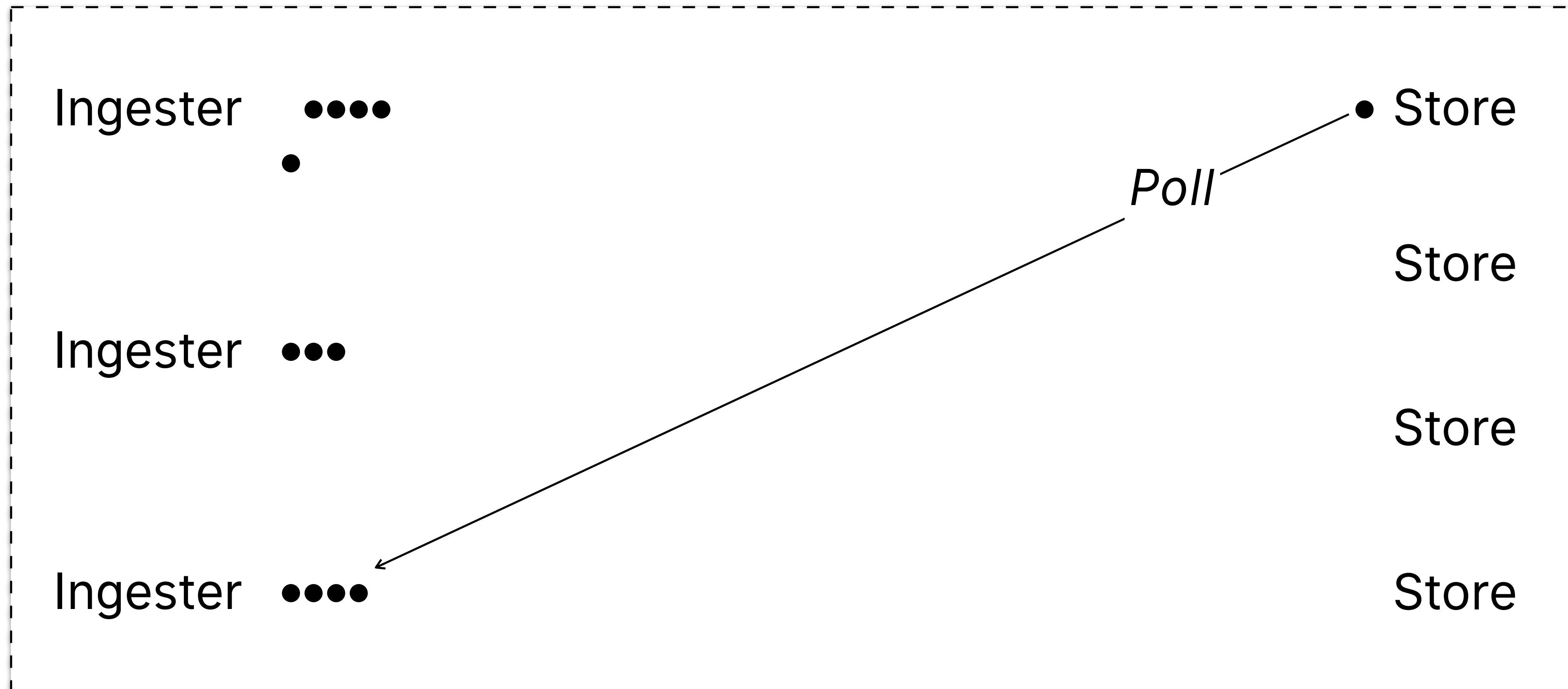
Replication



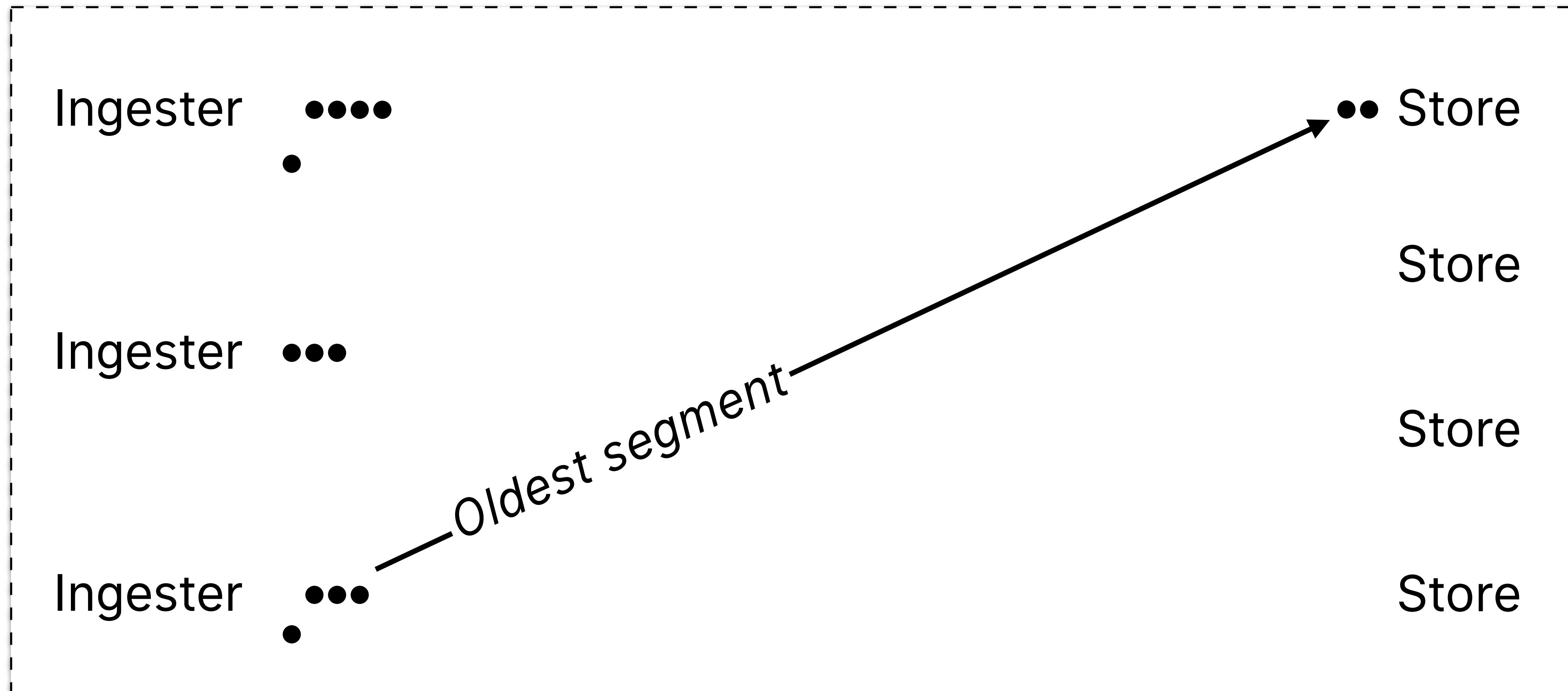
Replication



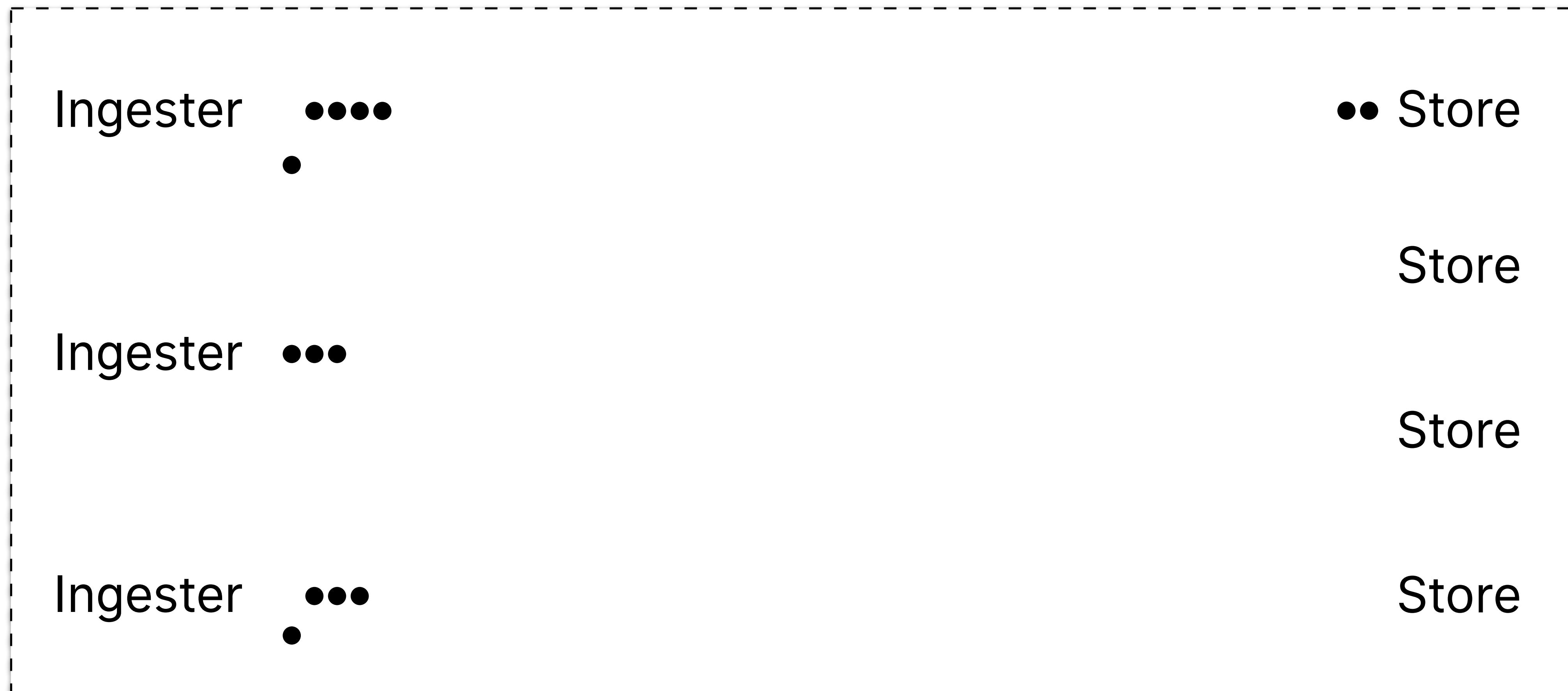
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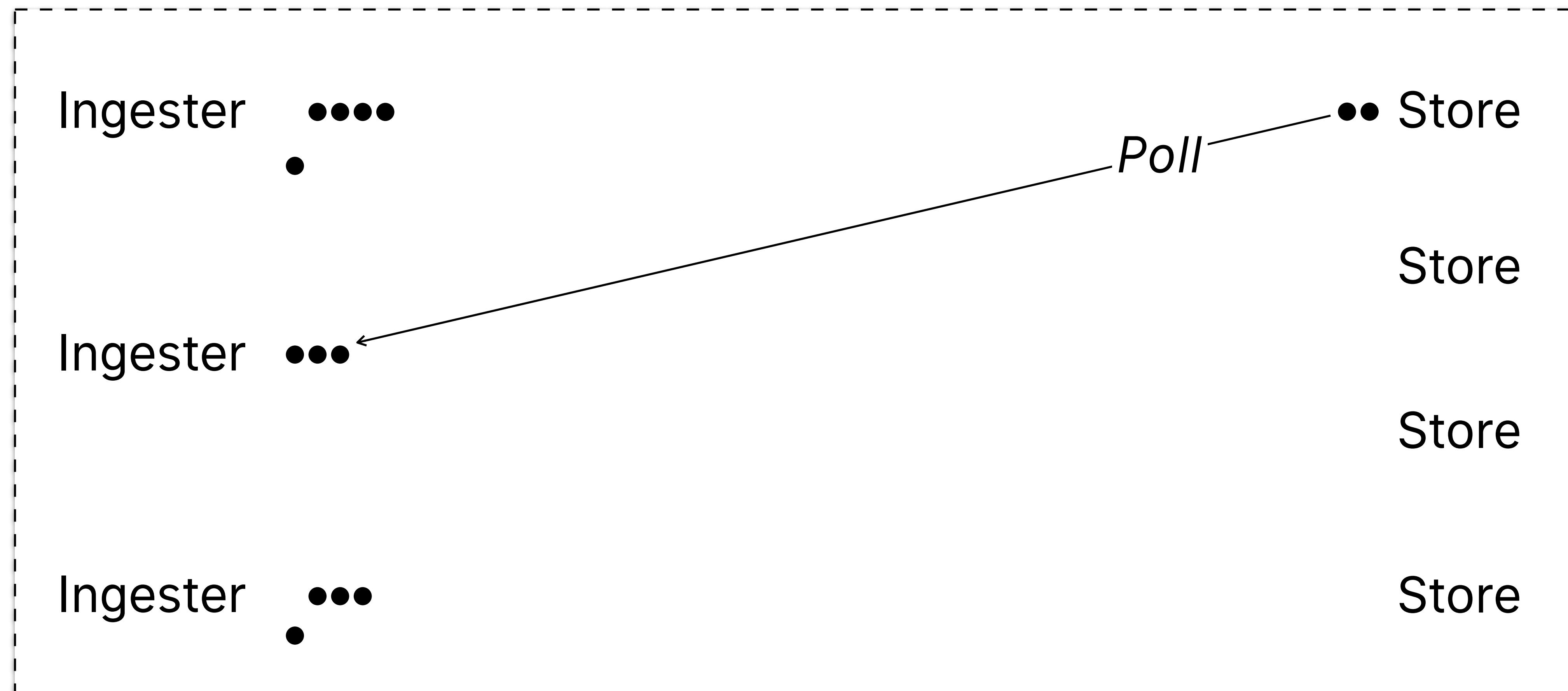
Replication



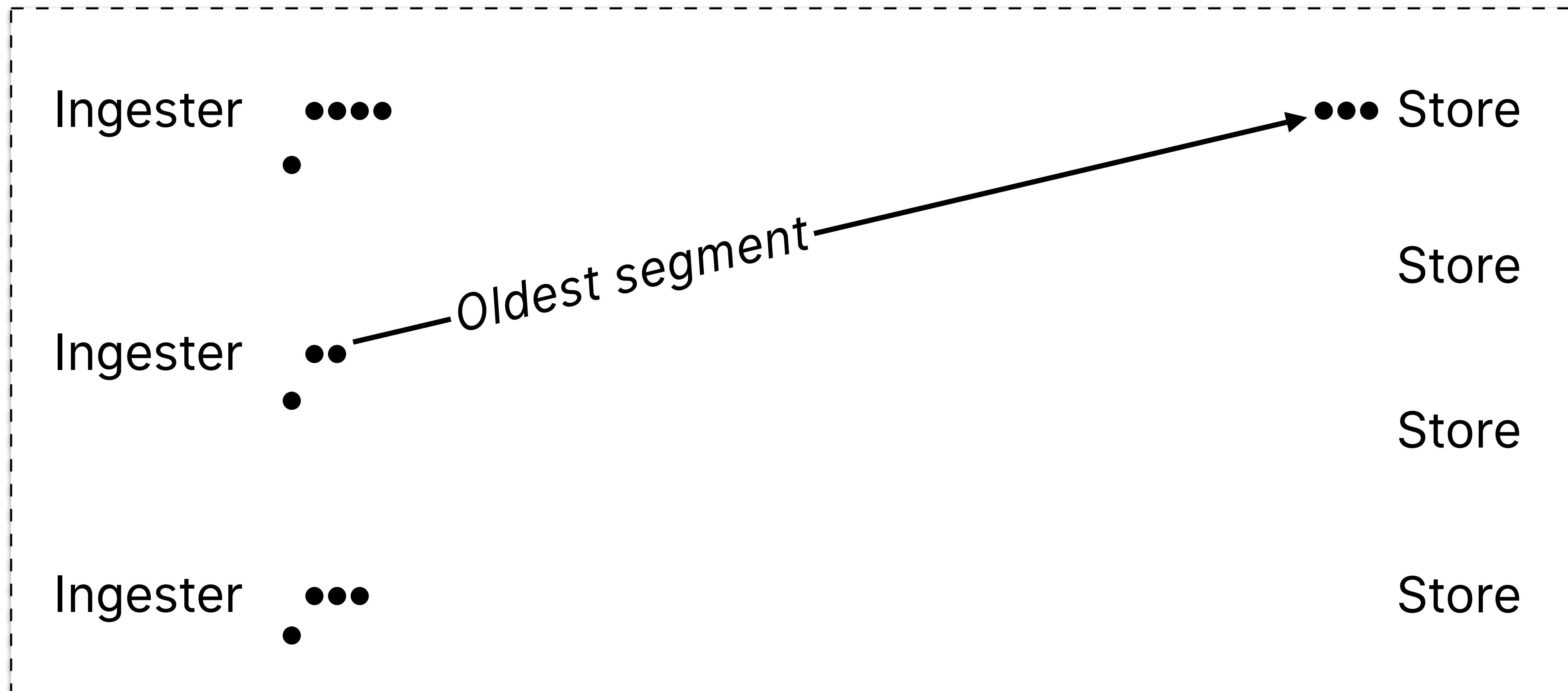
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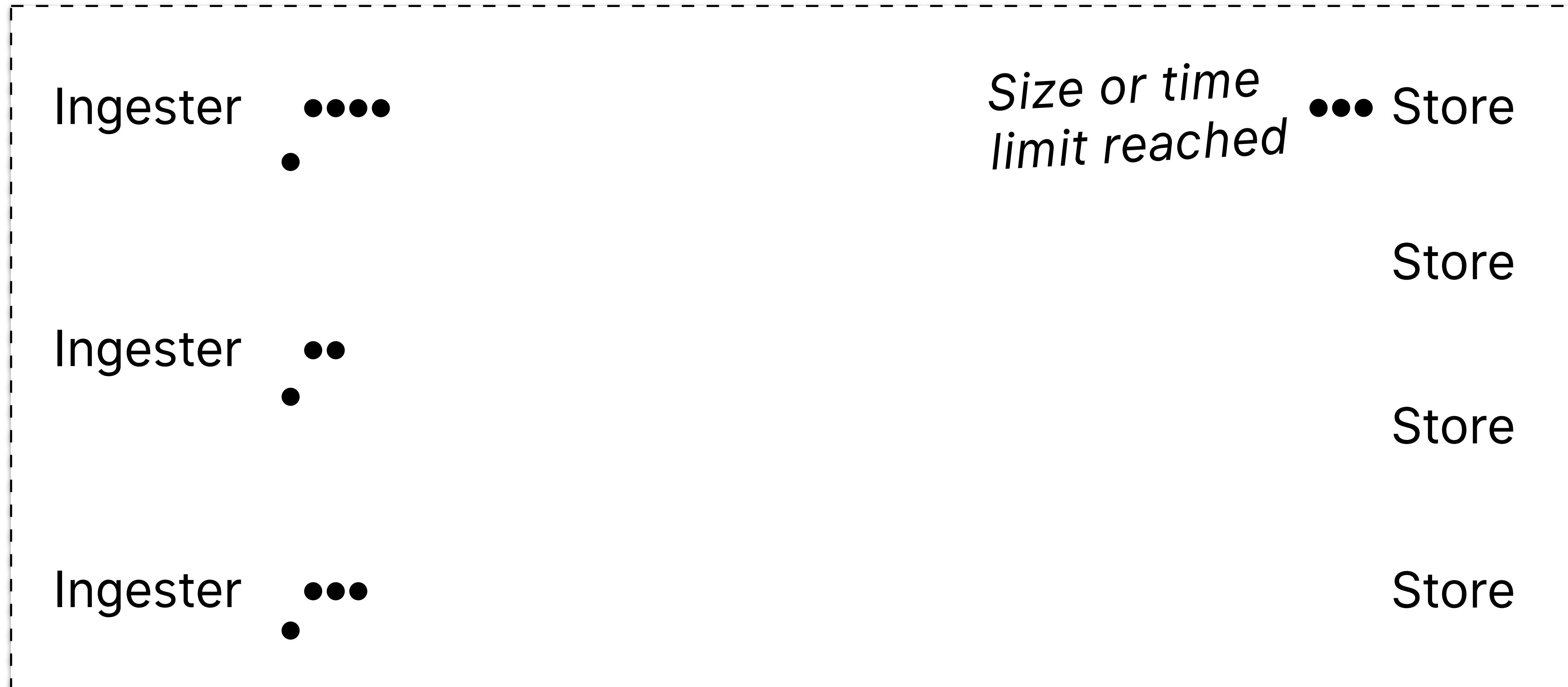
Replication



Replication



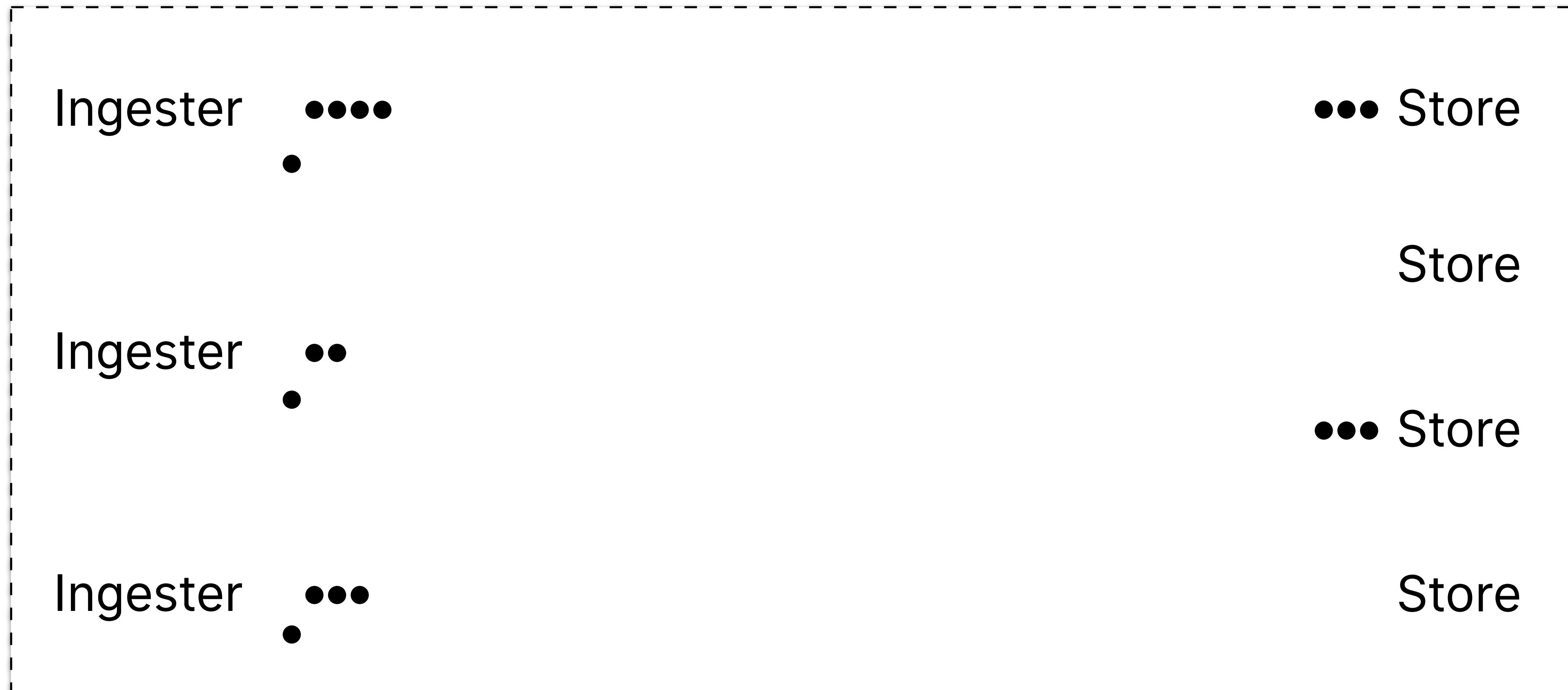
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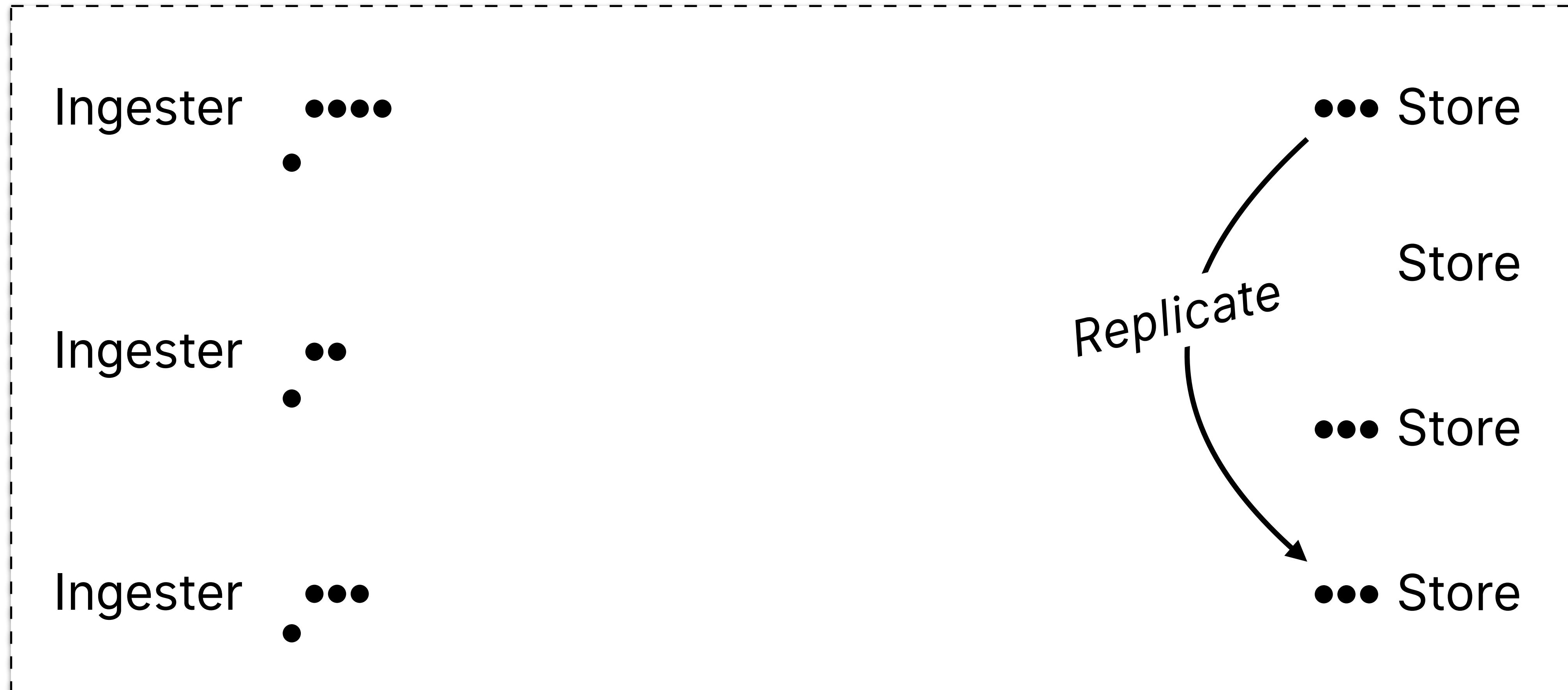
Replication



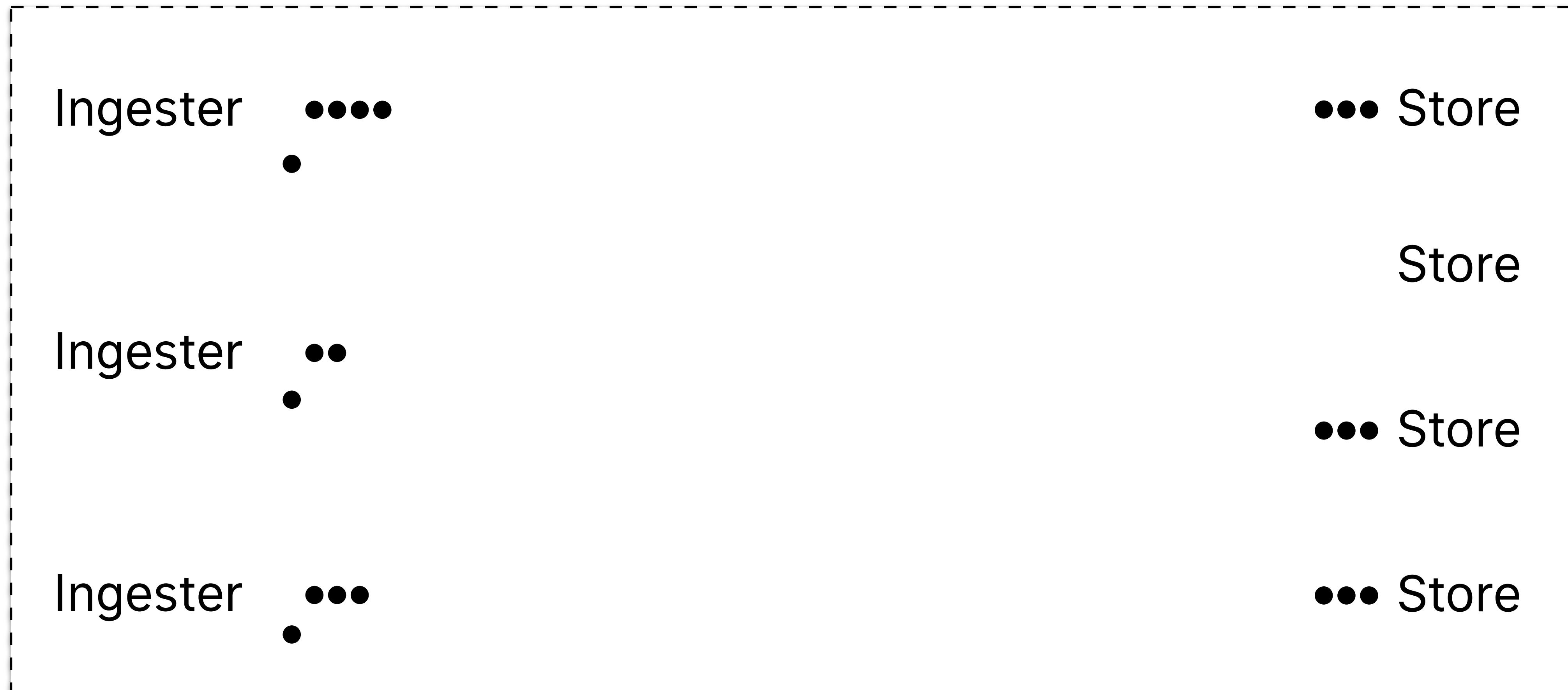
Replication



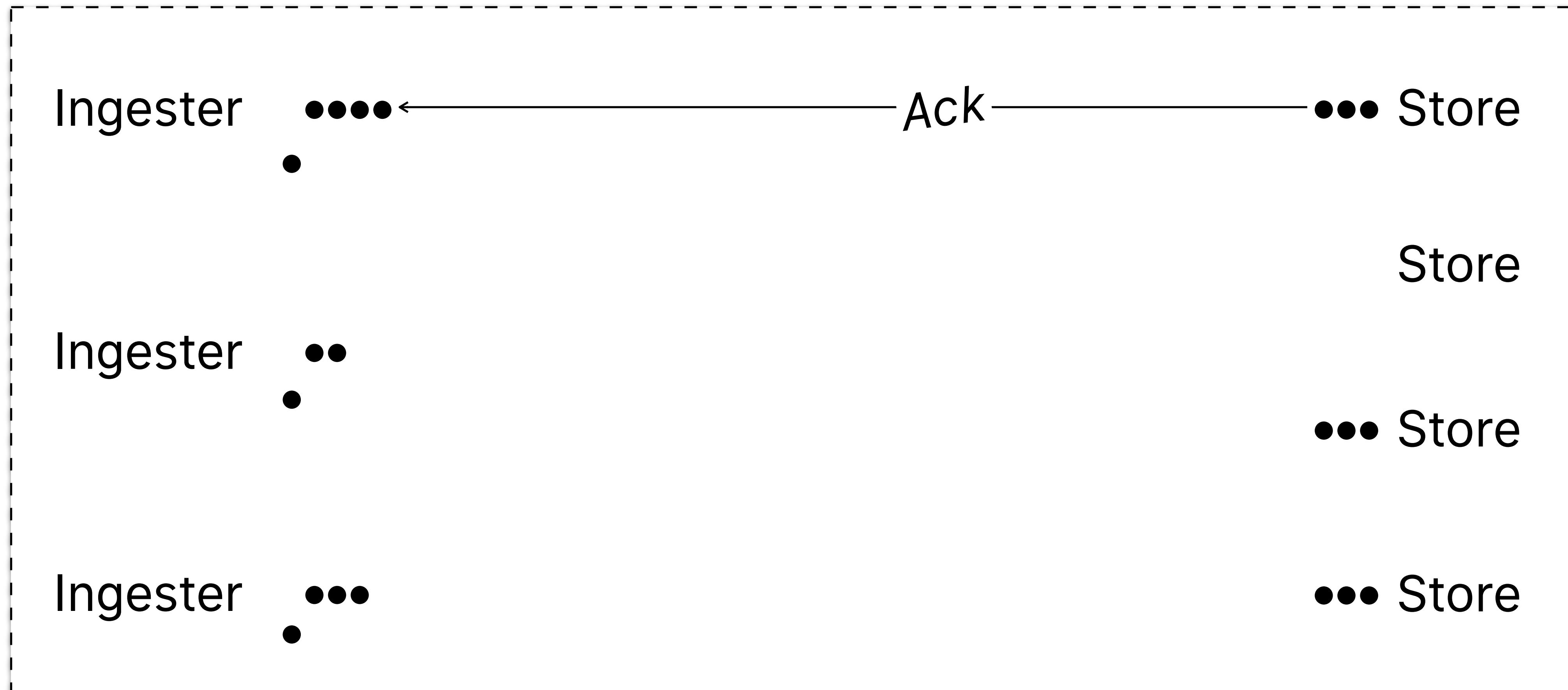
Replication



Replication



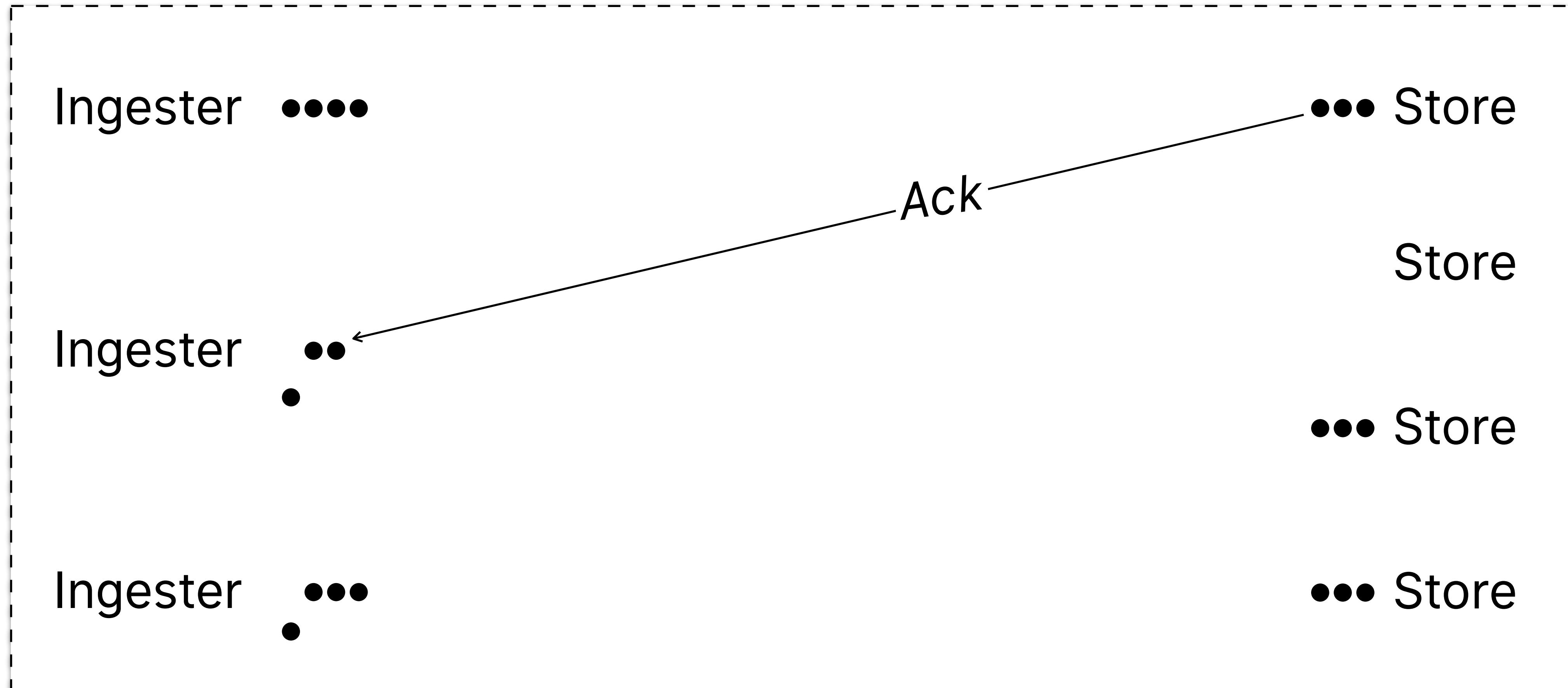
Replication



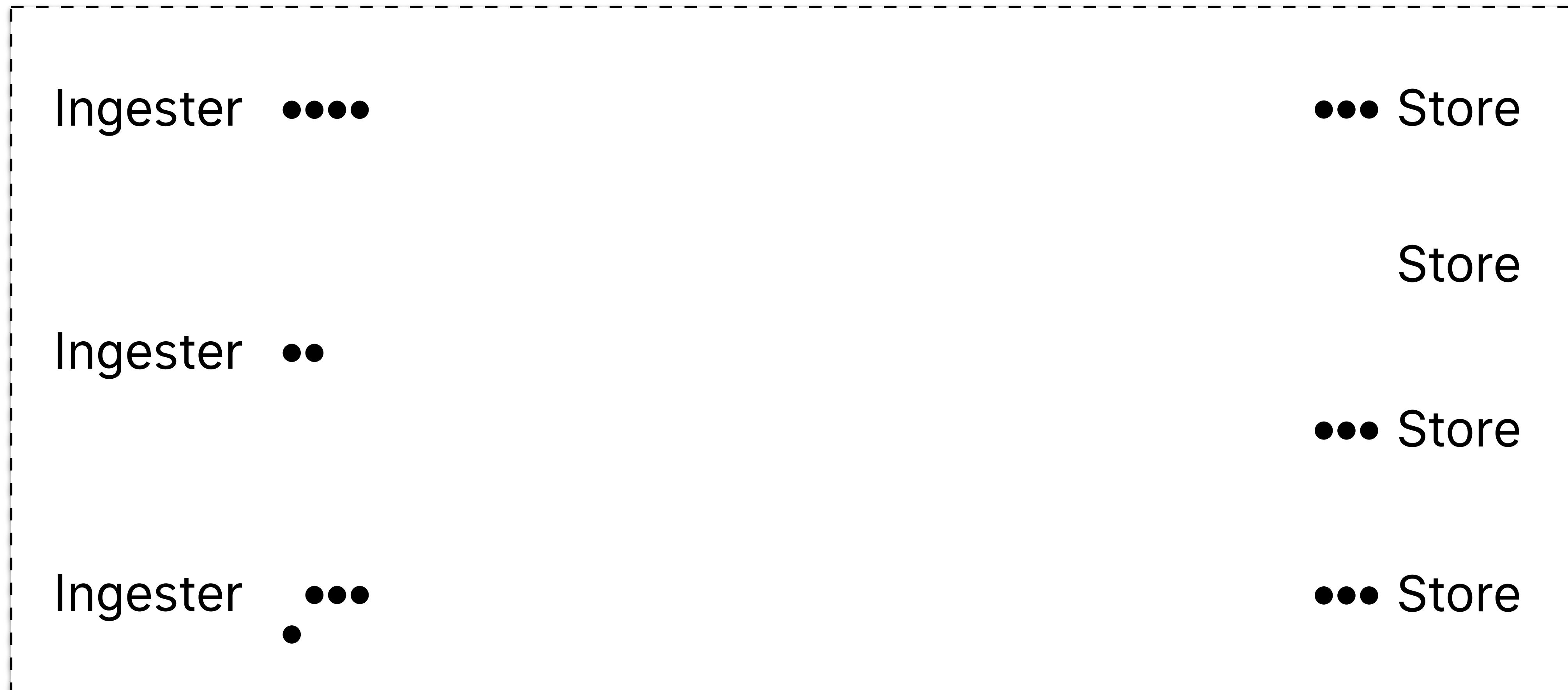
Replication



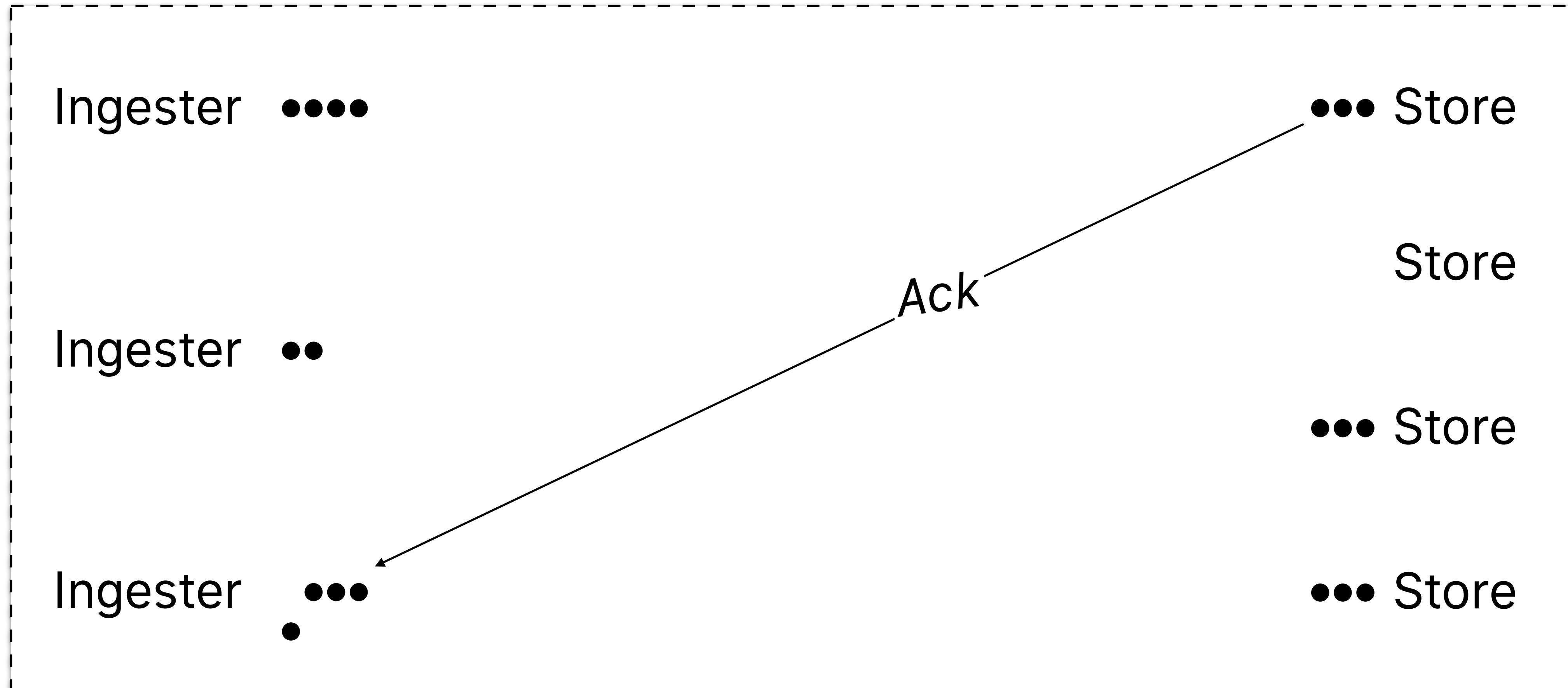
Replication



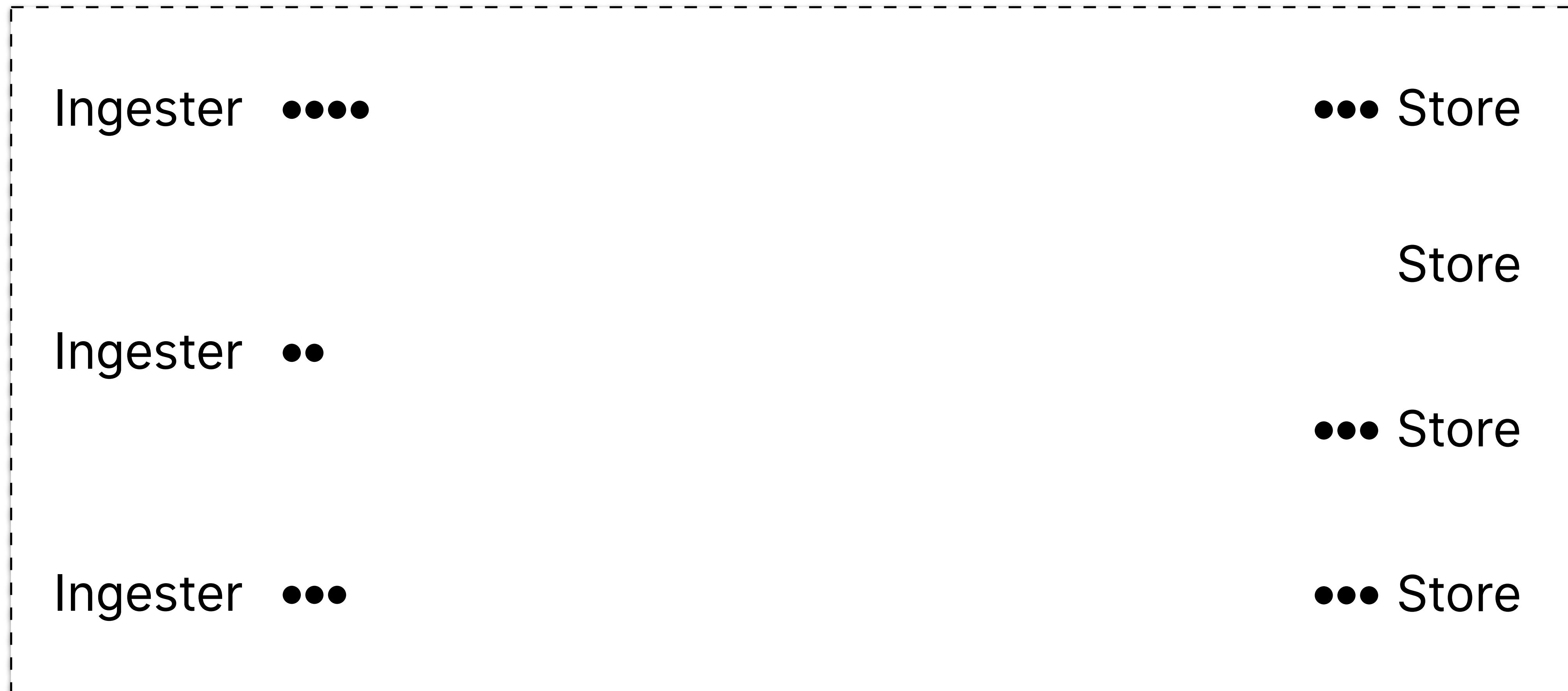
Replication



Replication



Replication



Transactional coördination

- Okay, it's not totally coördination-free 😐
- We have coördination within each gather/replicate transaction
- But those transactions are short-lived and tolerate failure well



Failure during gather

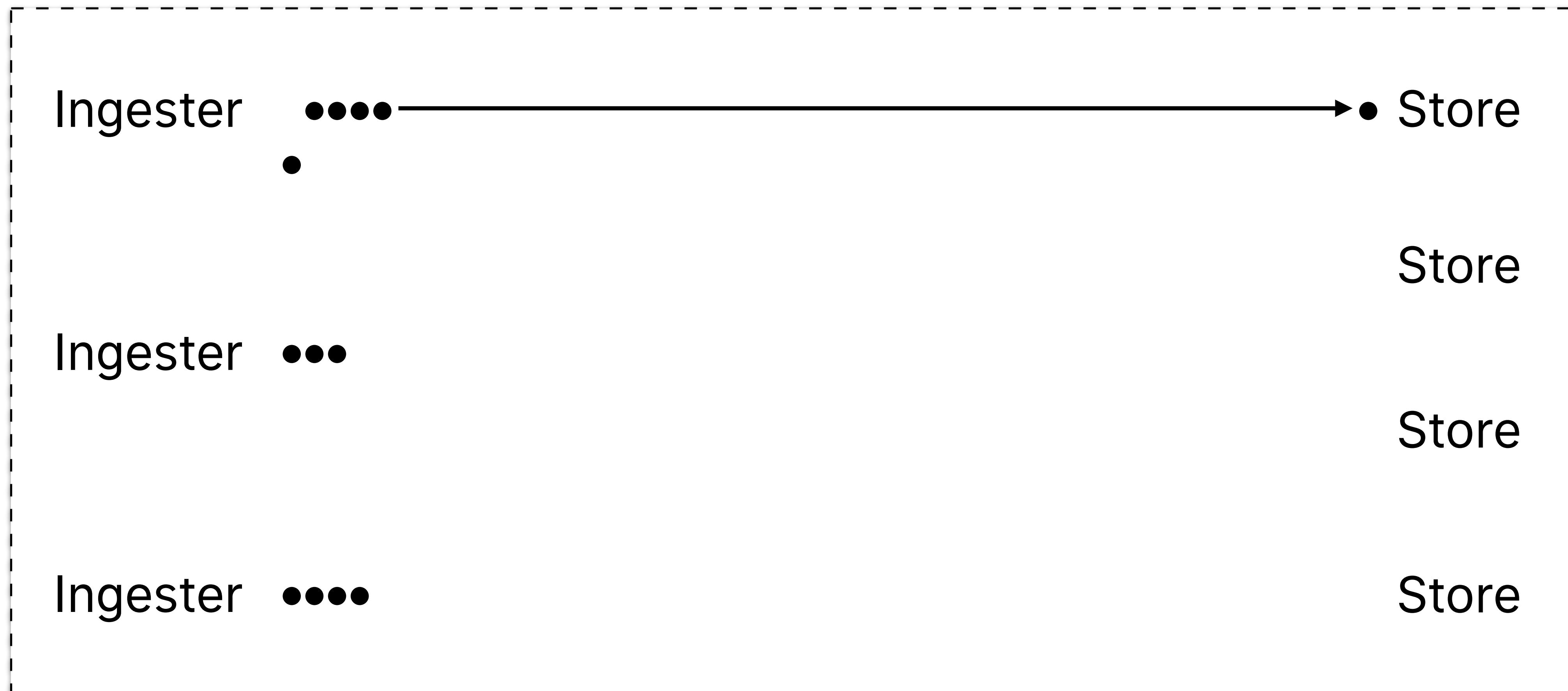
Timeouts to the rescue



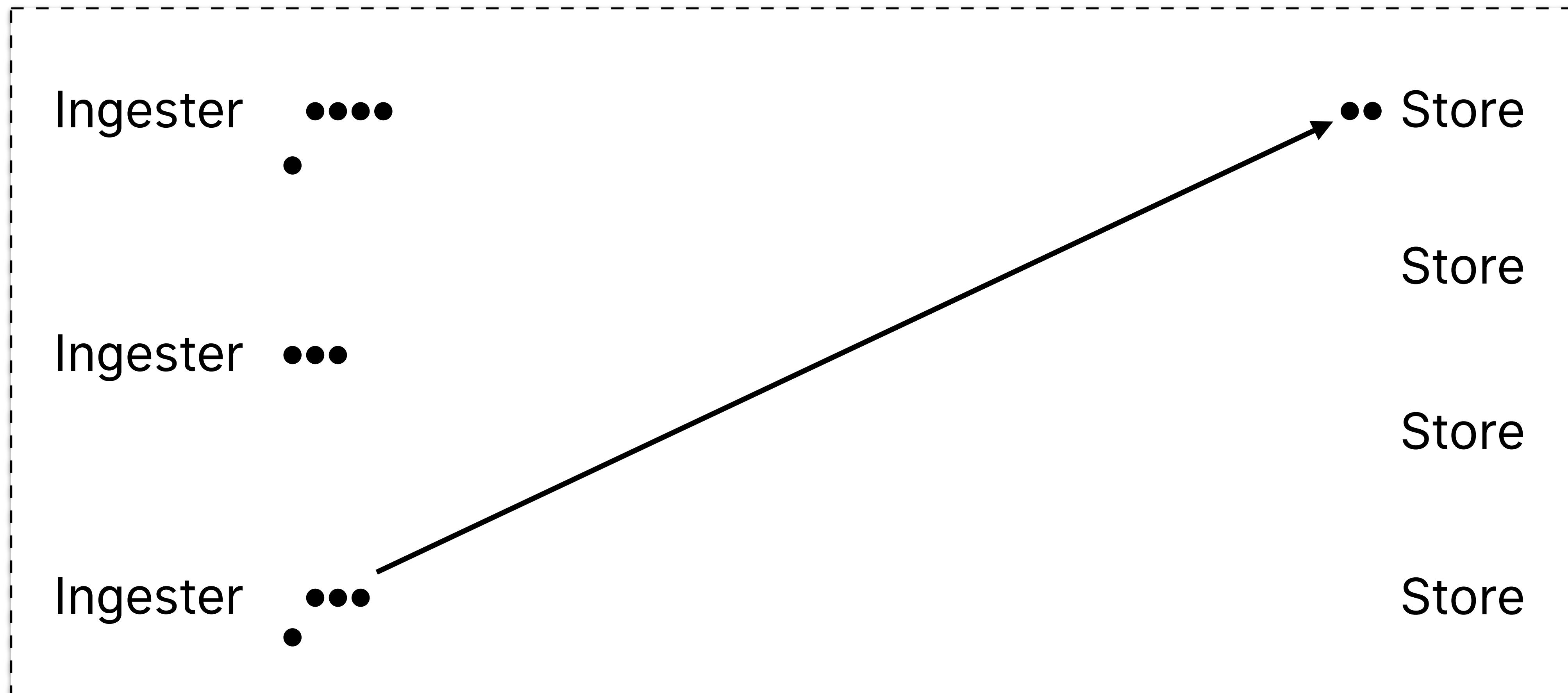
Failure during gather



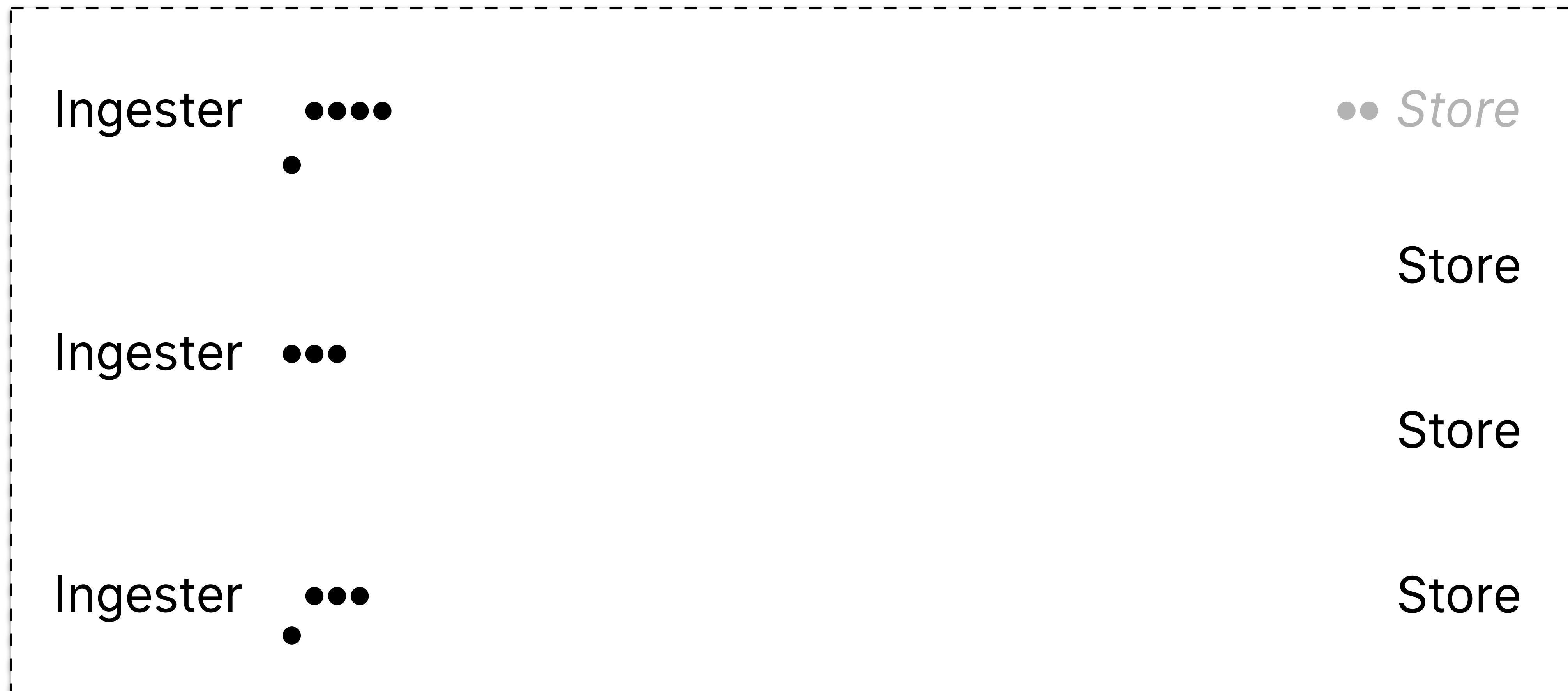
Failure during gather



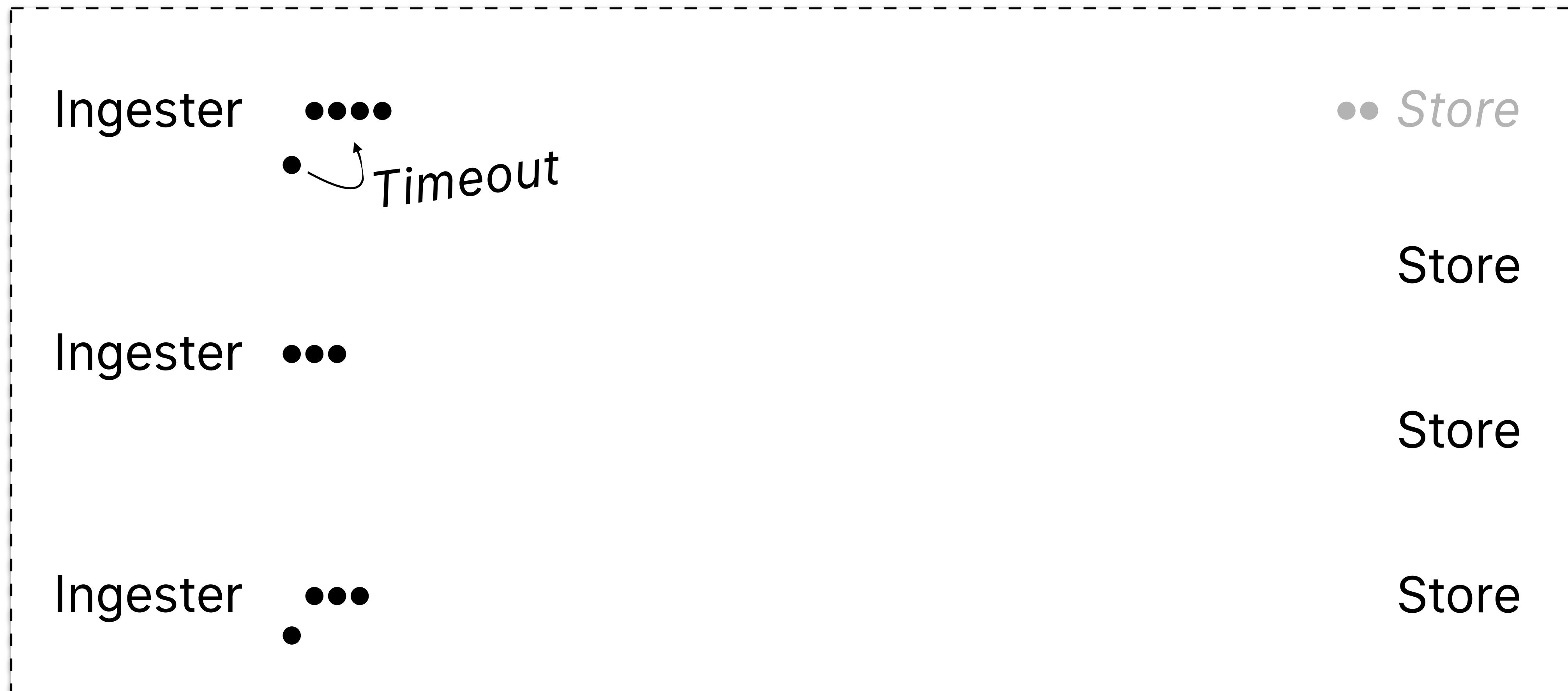
Failure during gather



Failure during gather



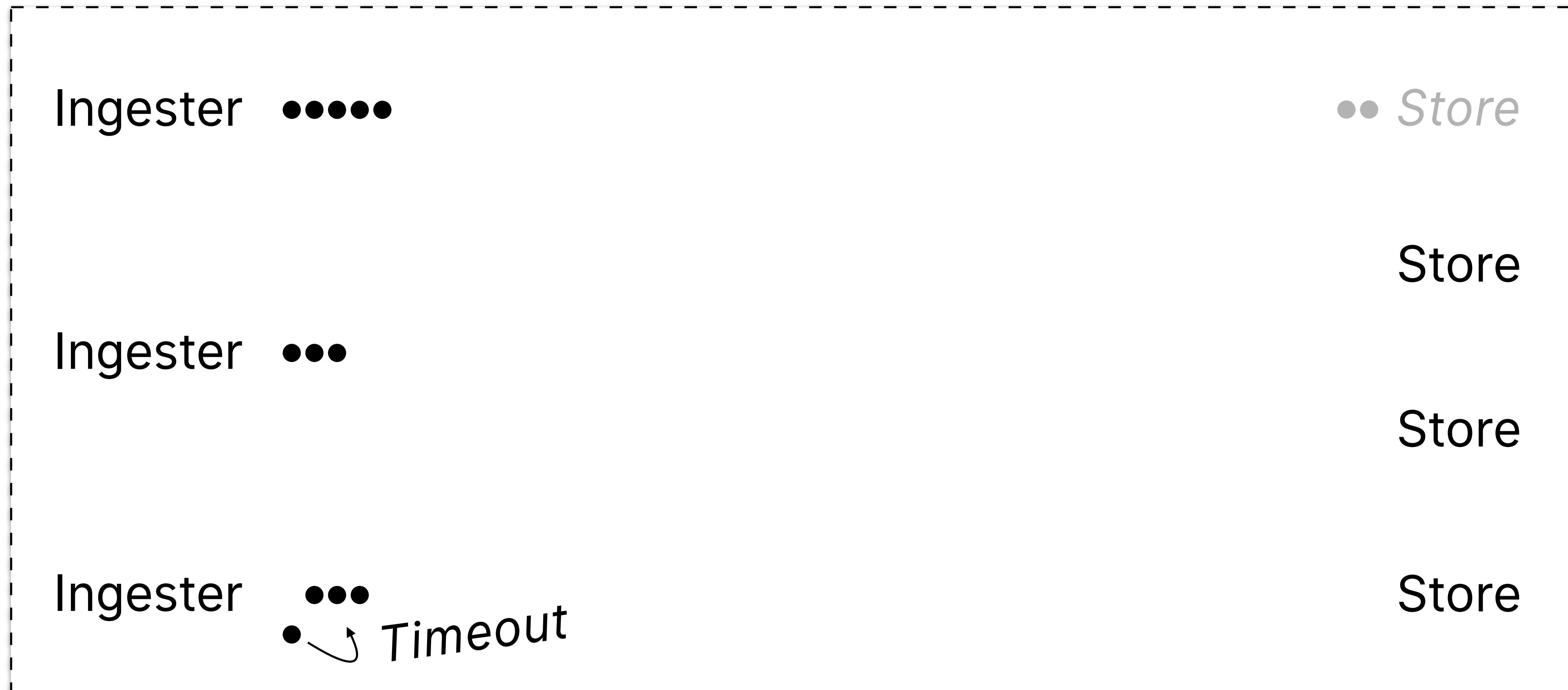
Failure during gather



Failure during gather



Failure during gather



Failure during gather

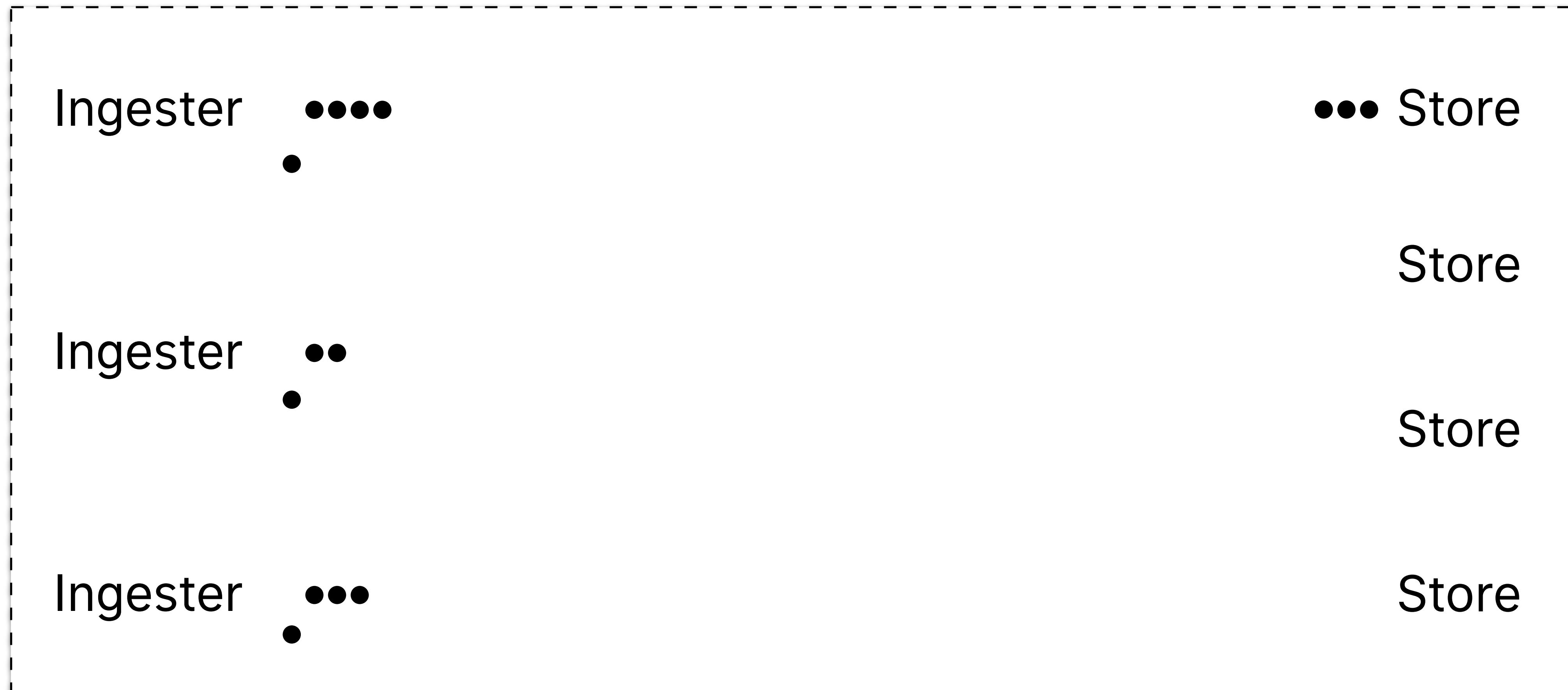


Failure during replicate

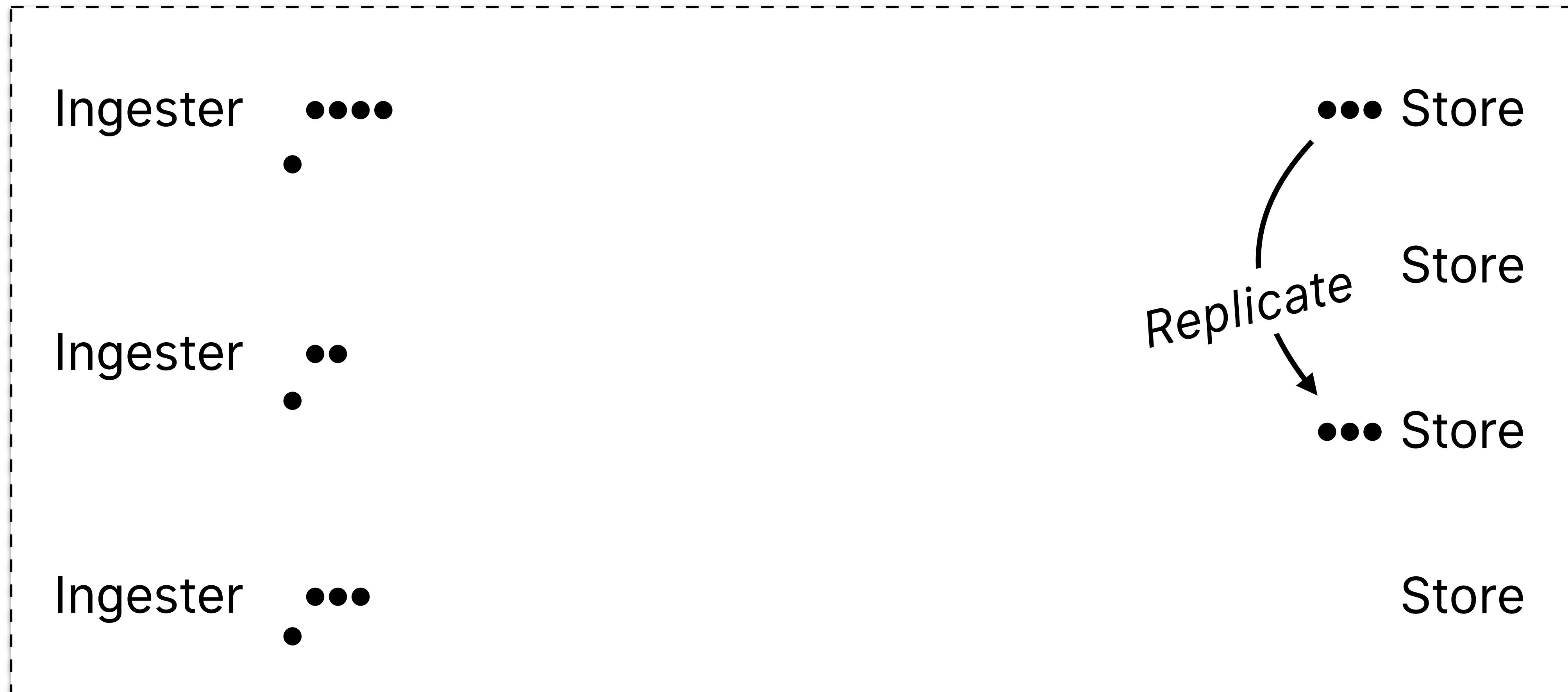
We've got dupes!



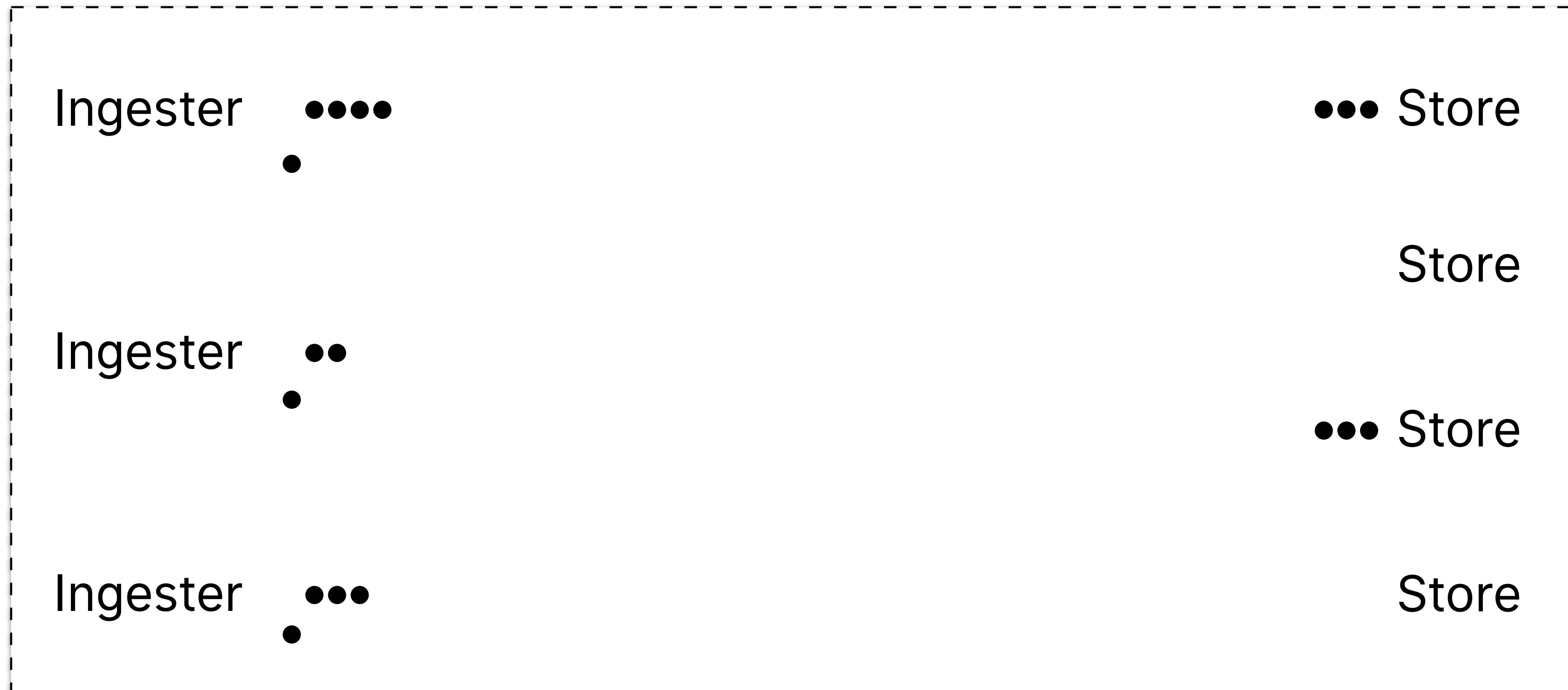
Failure during replicate



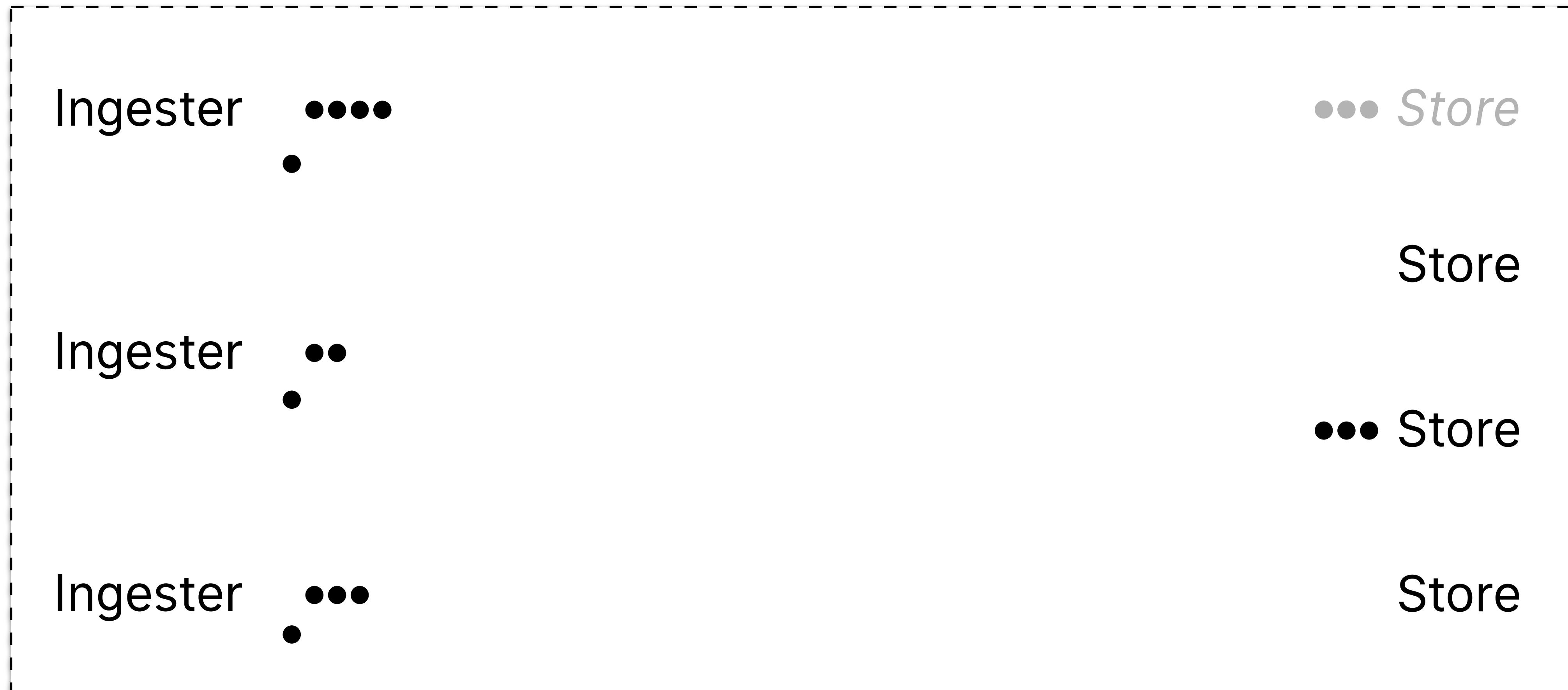
Failure during replicate



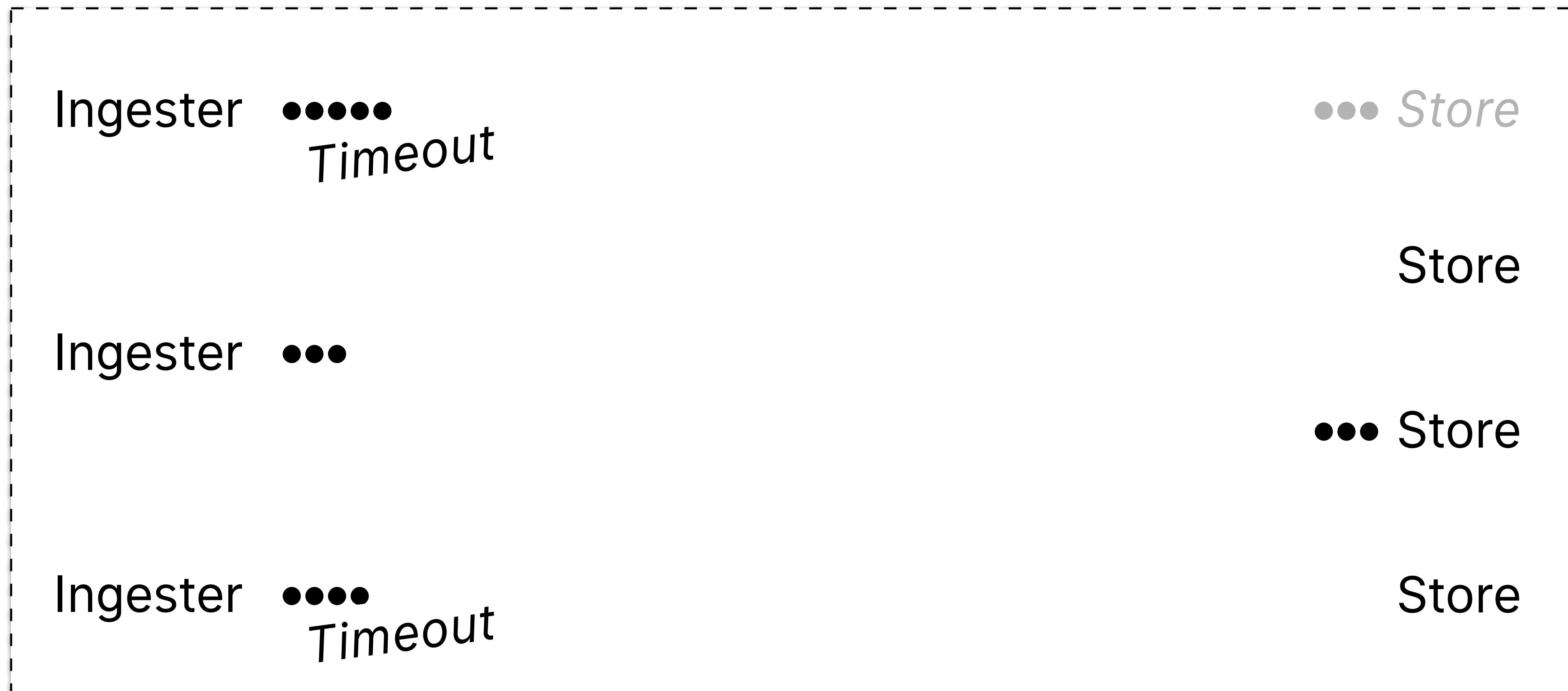
Failure during replicate



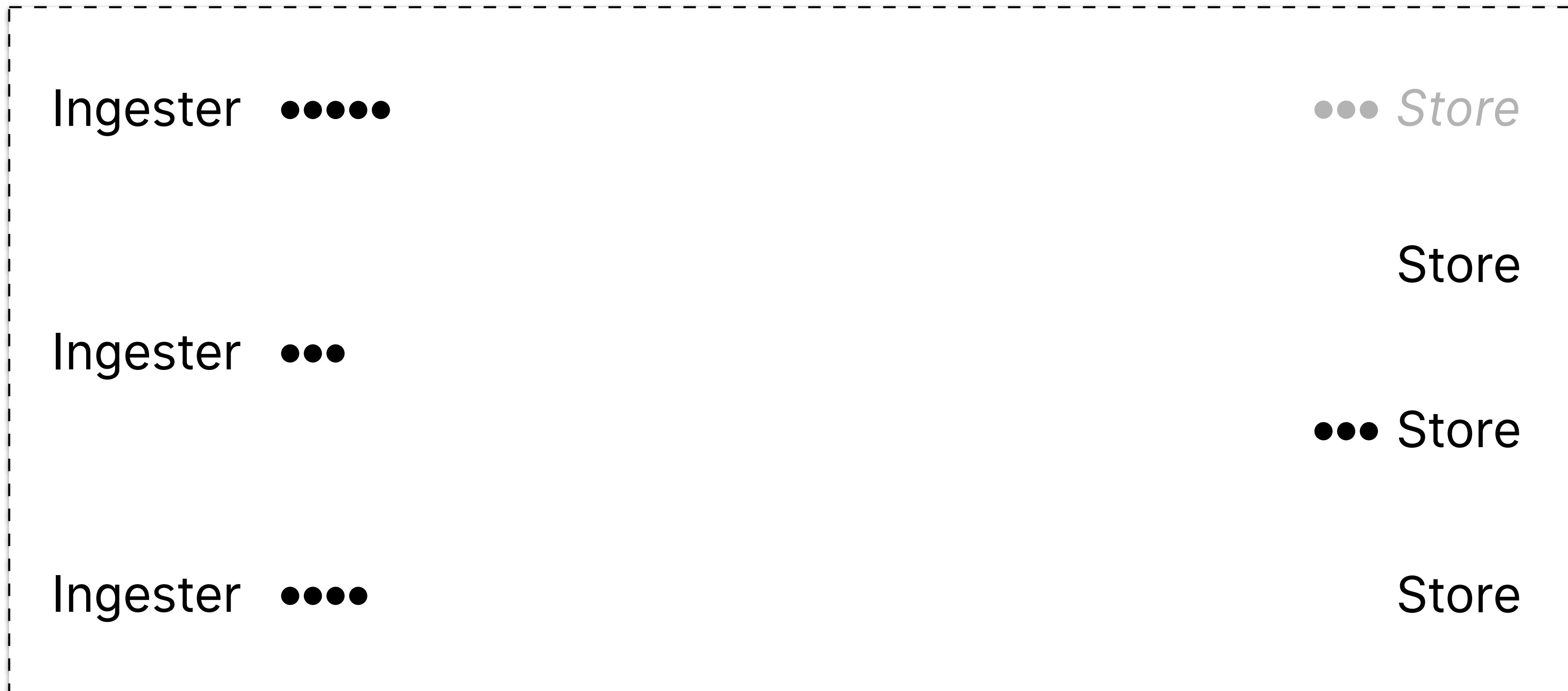
Failure during replicate



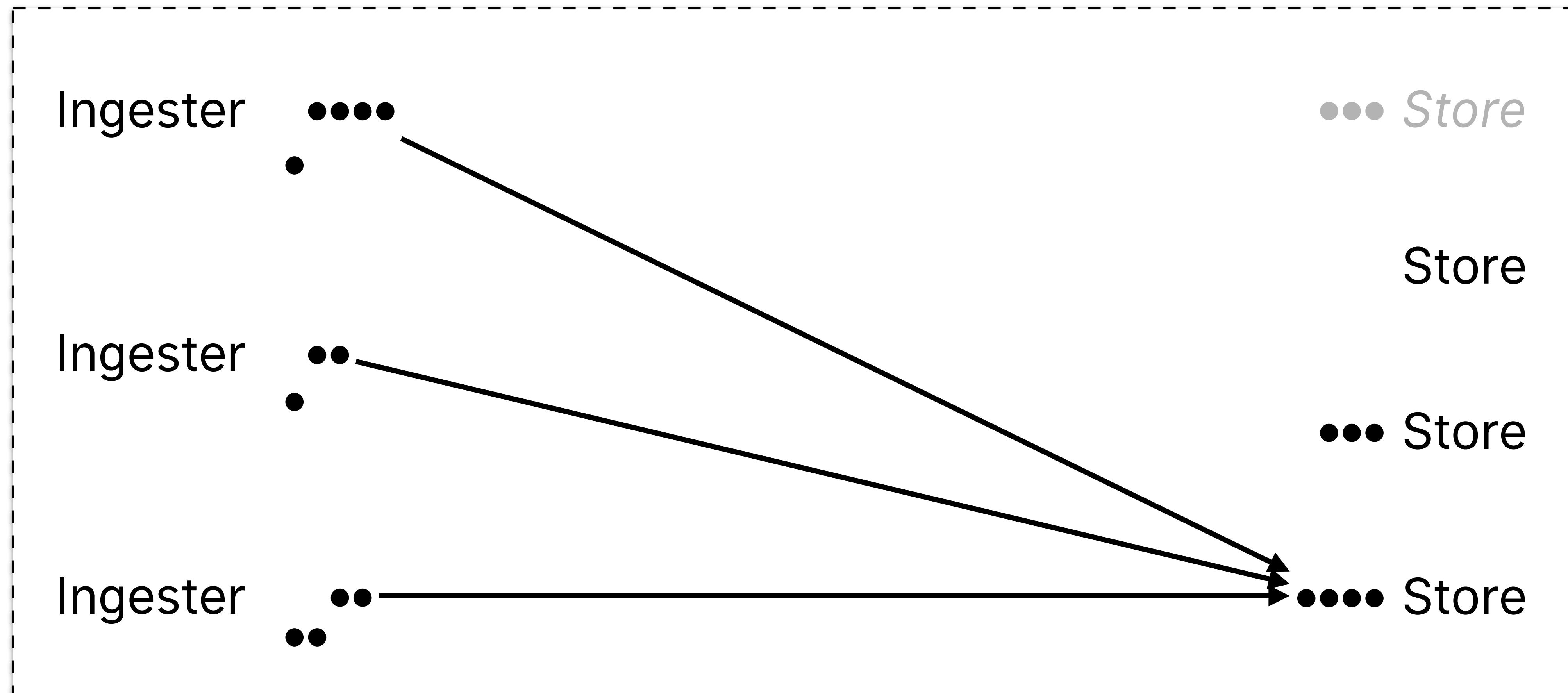
Failure during replicate



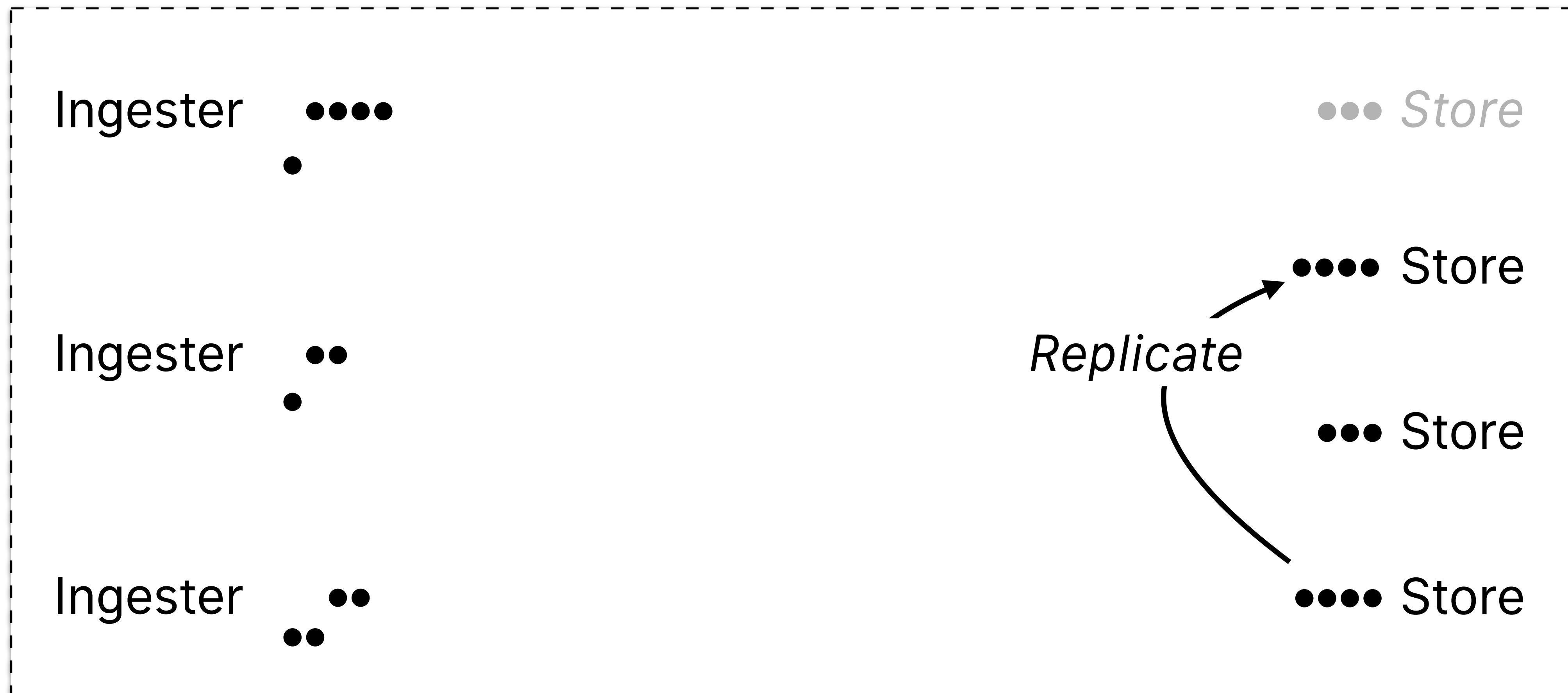
Failure during replicate



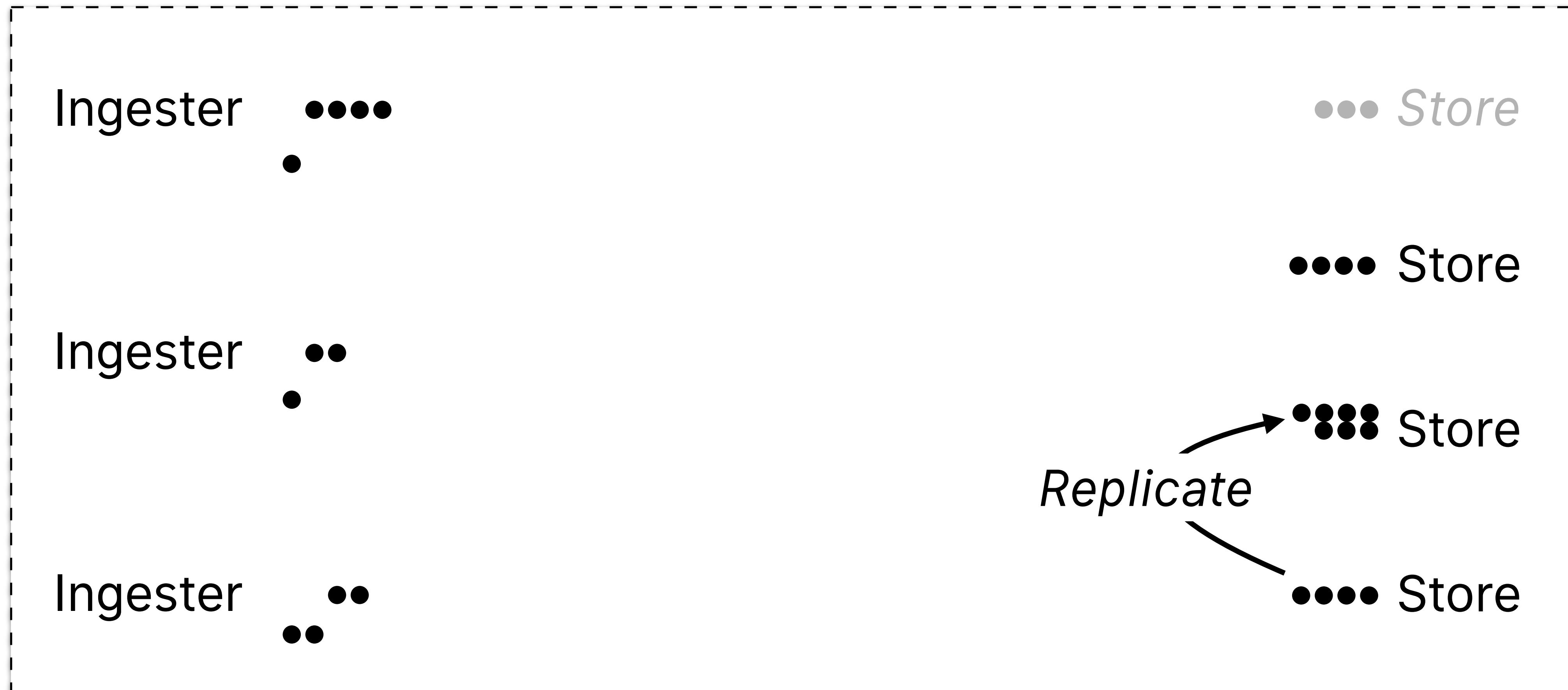
Failure during replicate



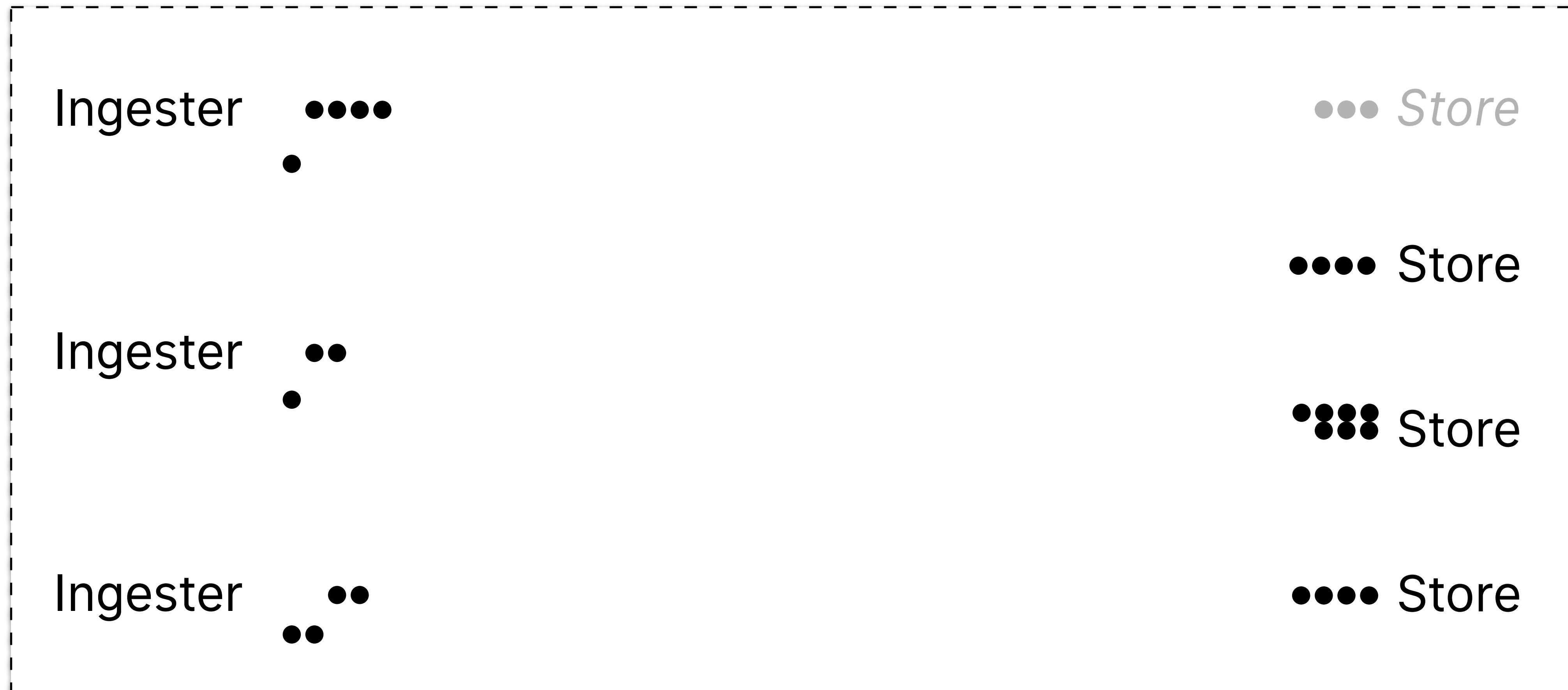
Failure during replicate



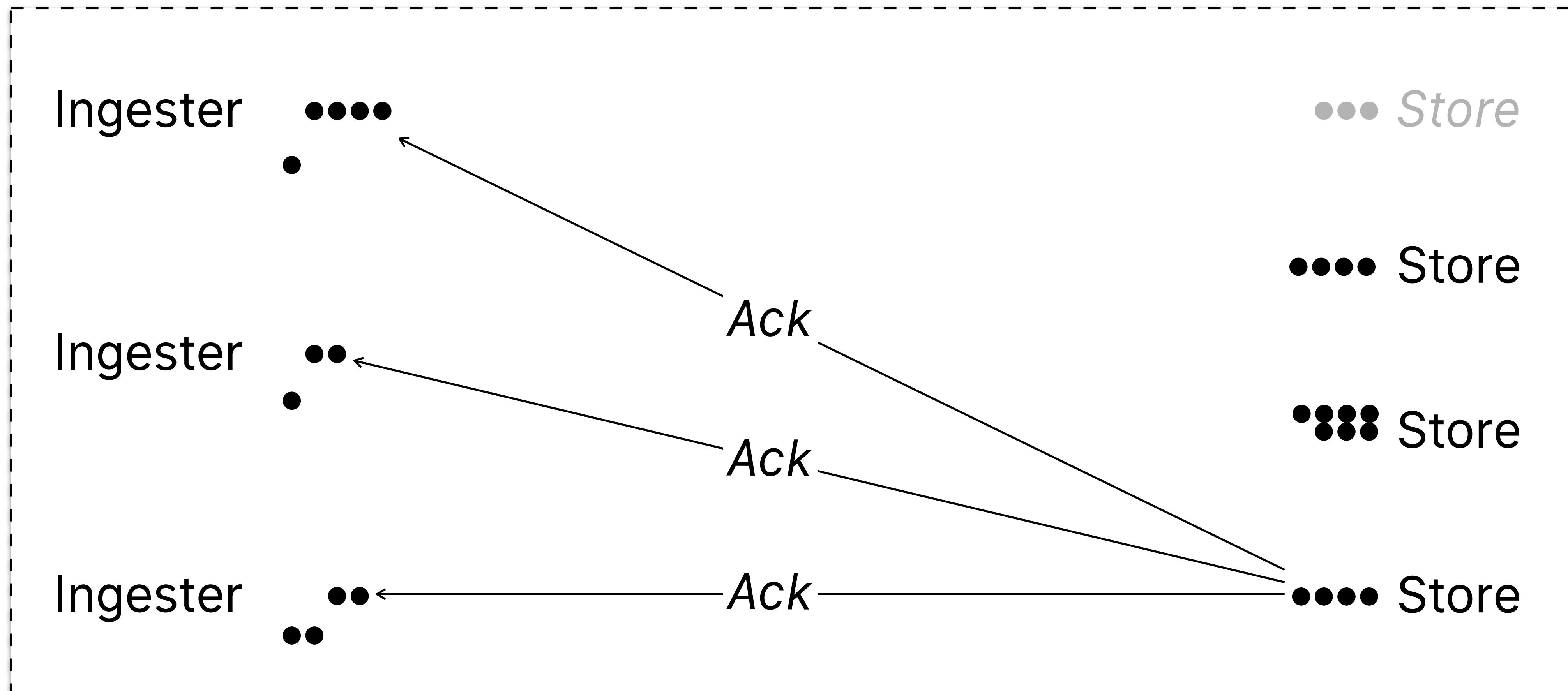
Failure during replicate



Failure during replicate



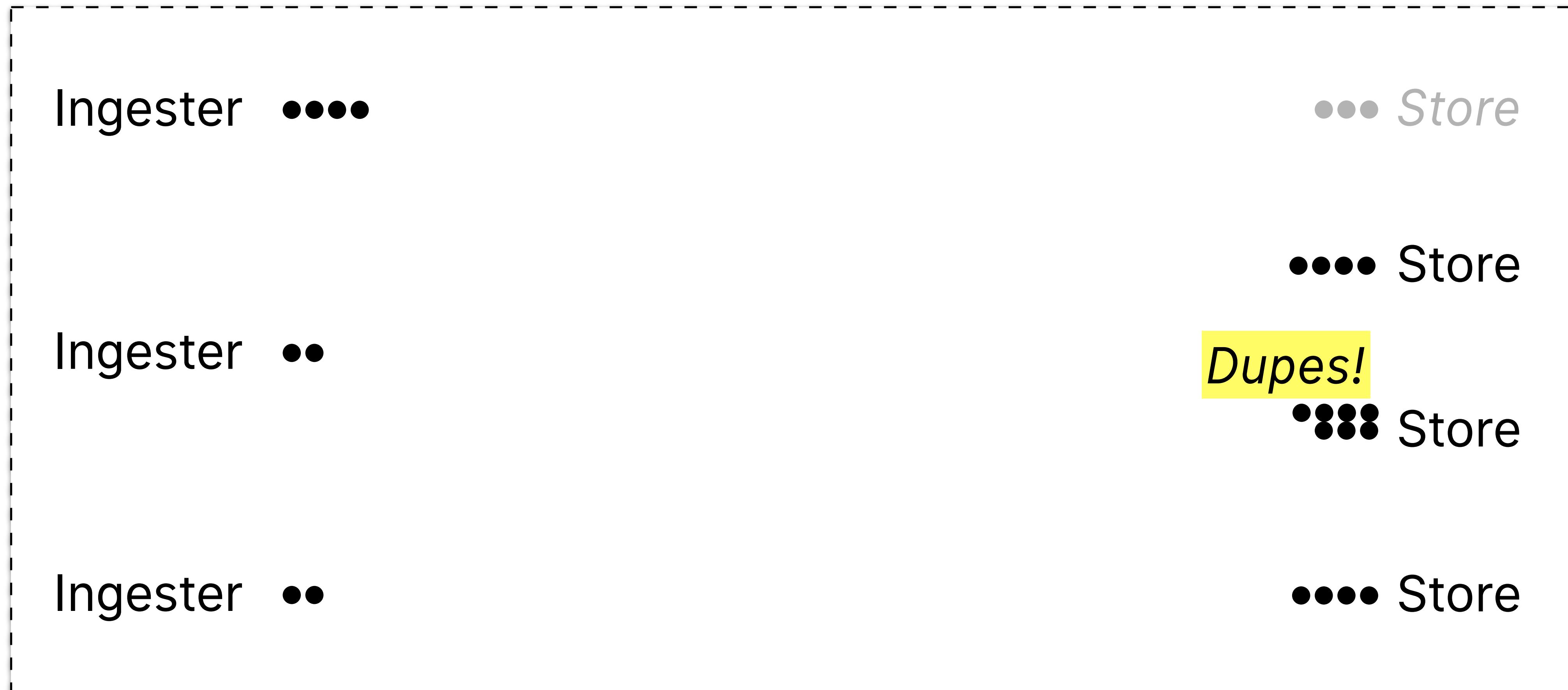
Failure during replicate



Failure during replicate



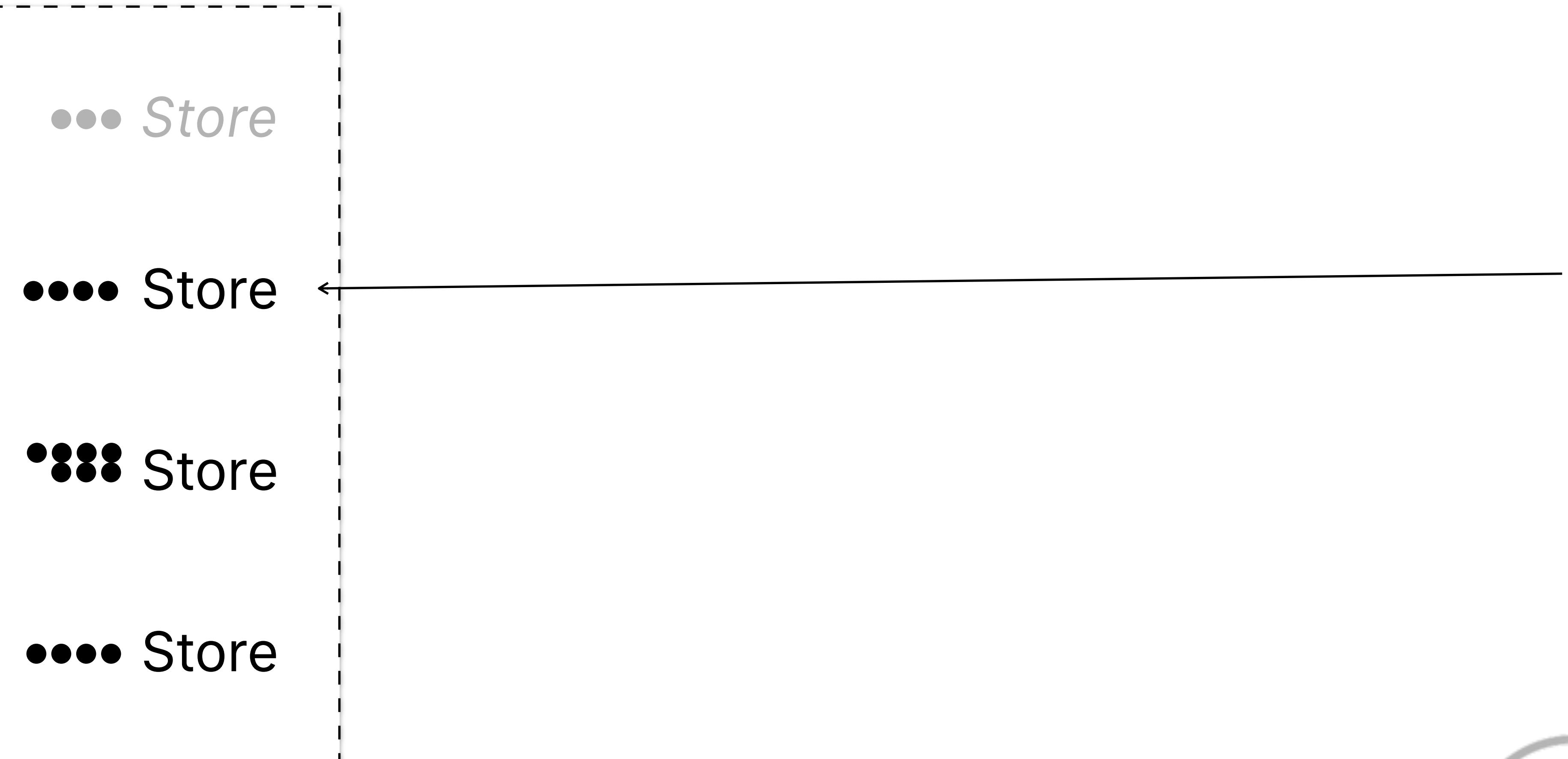
Failure during replicate



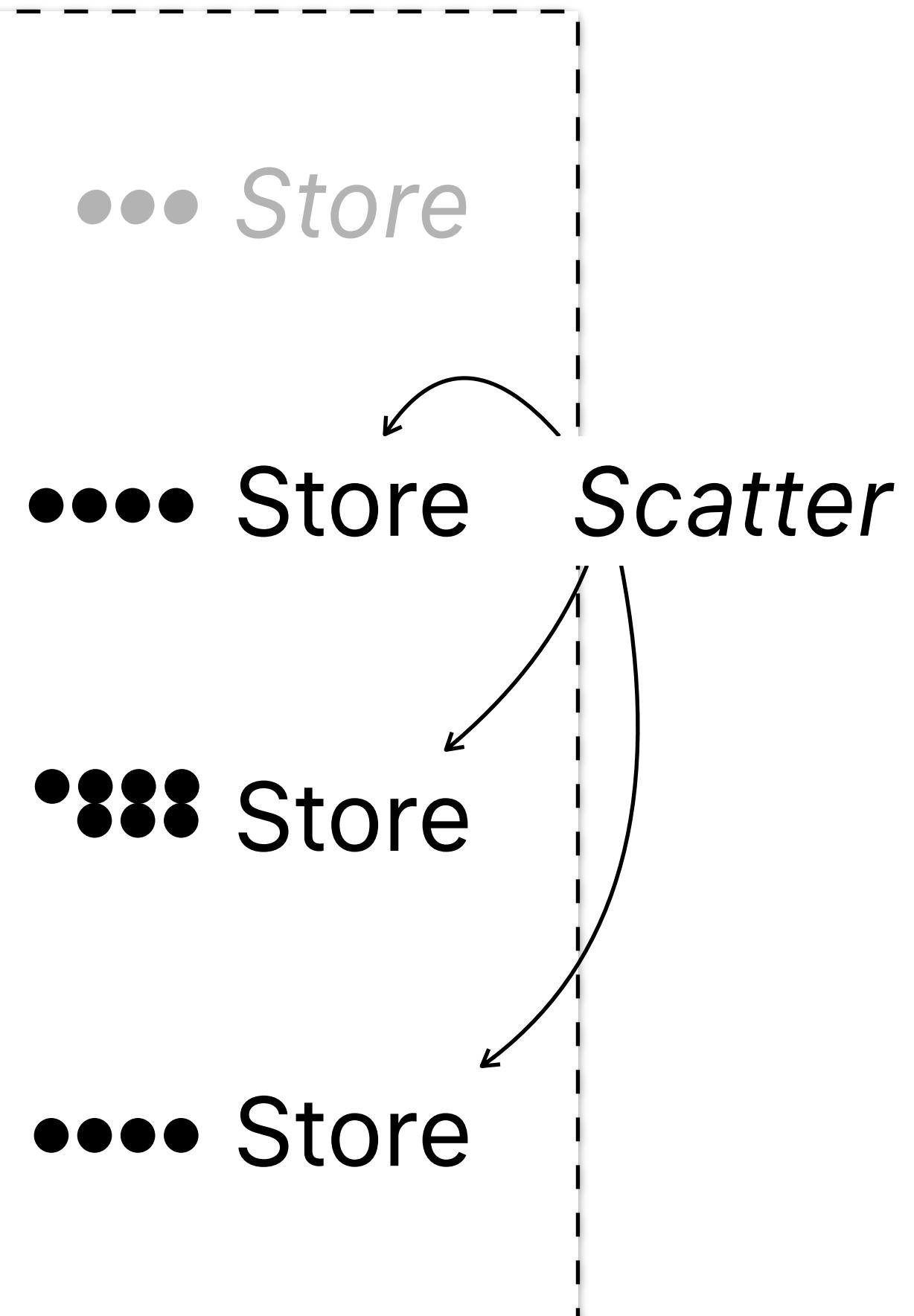
Dupes don't matter



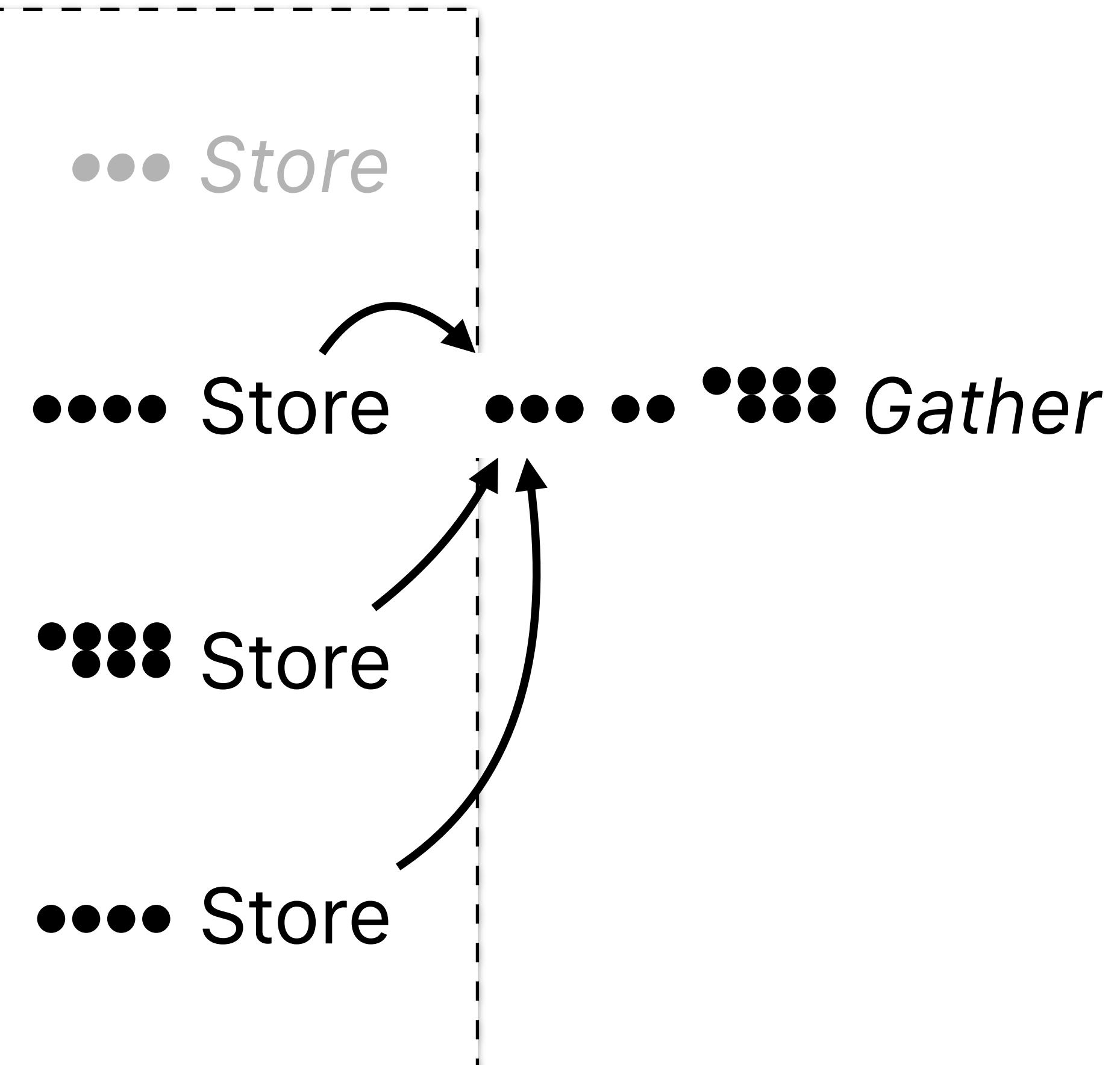
Dupes don't matter



Dupes don't matter



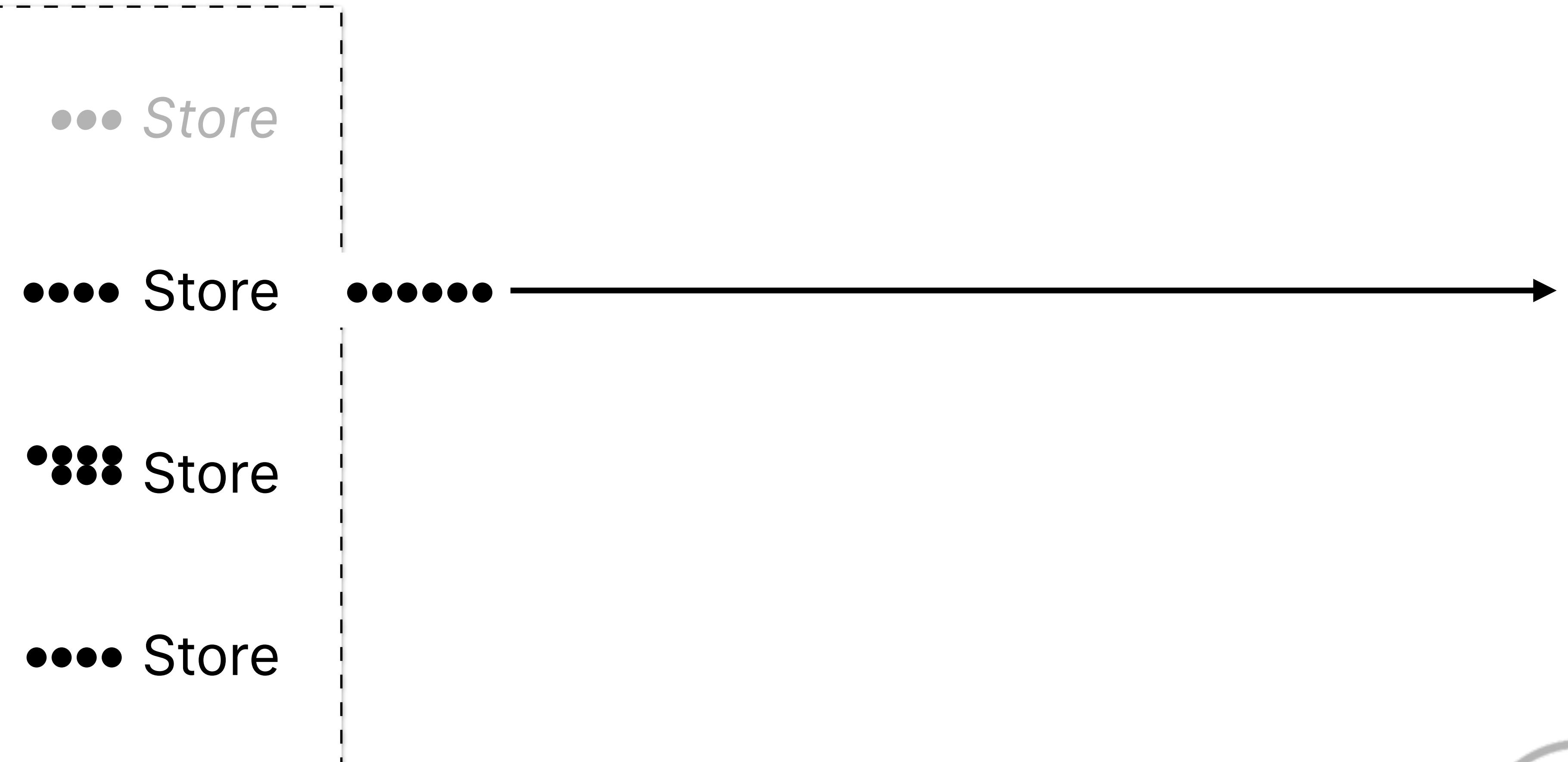
Dupes don't matter



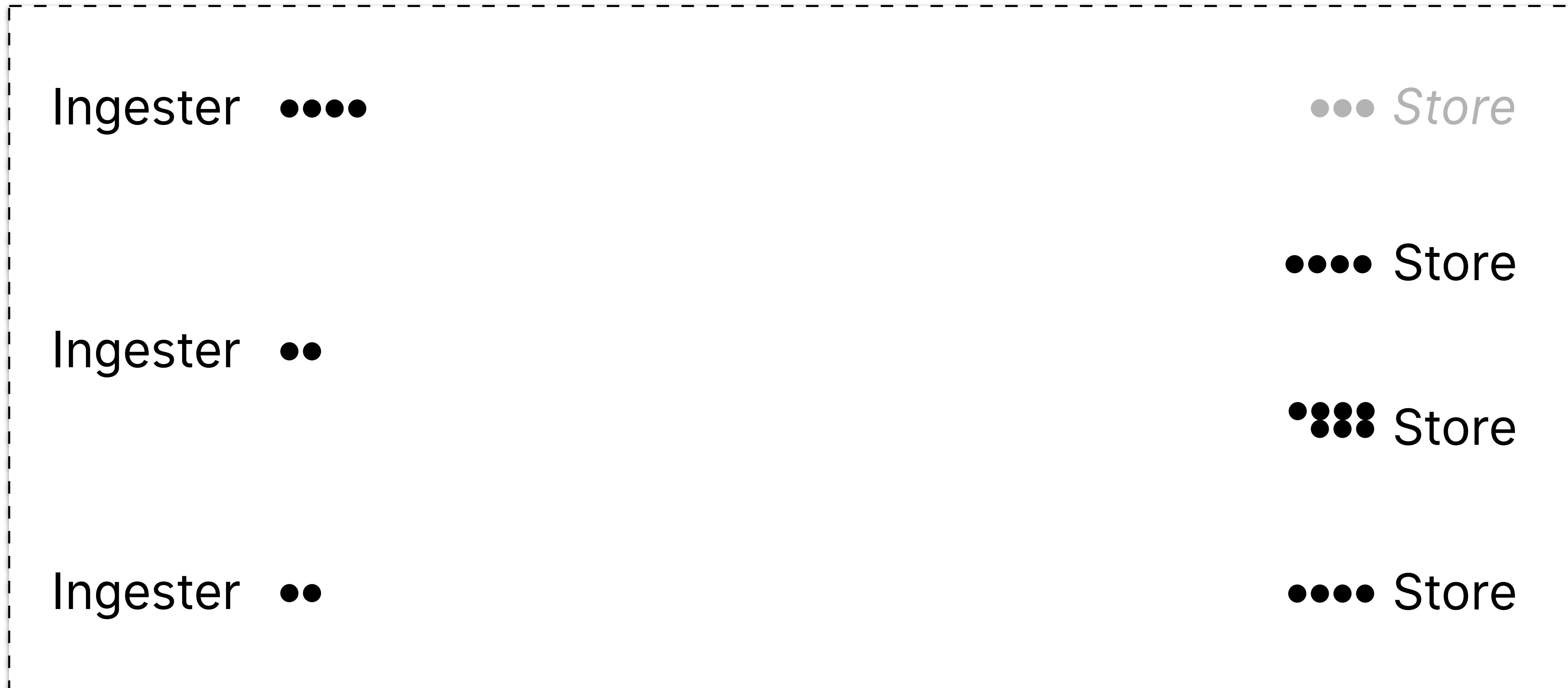
Dupes don't matter



Dupes don't matter



Dupes get compacted



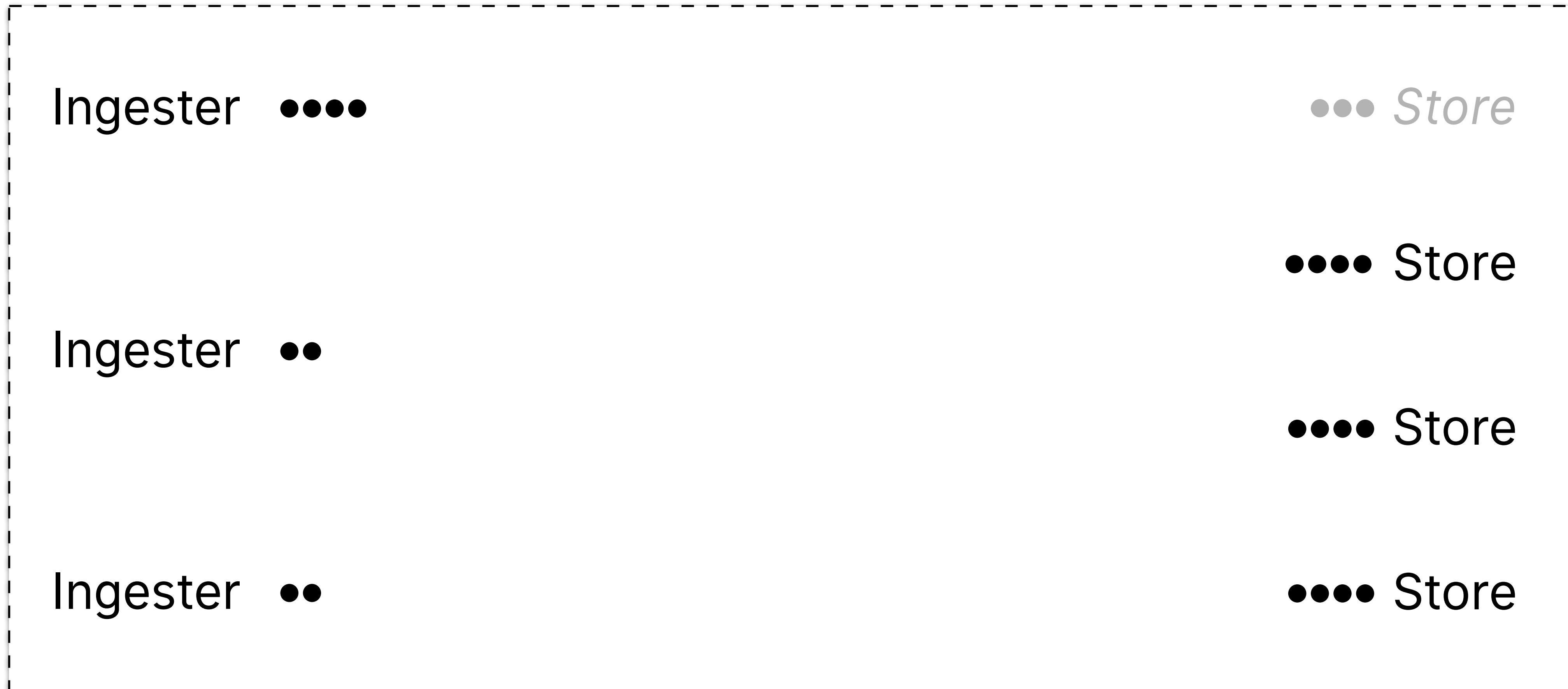
Dupes get compacted



Dupes get compacted



Dupes get compacted

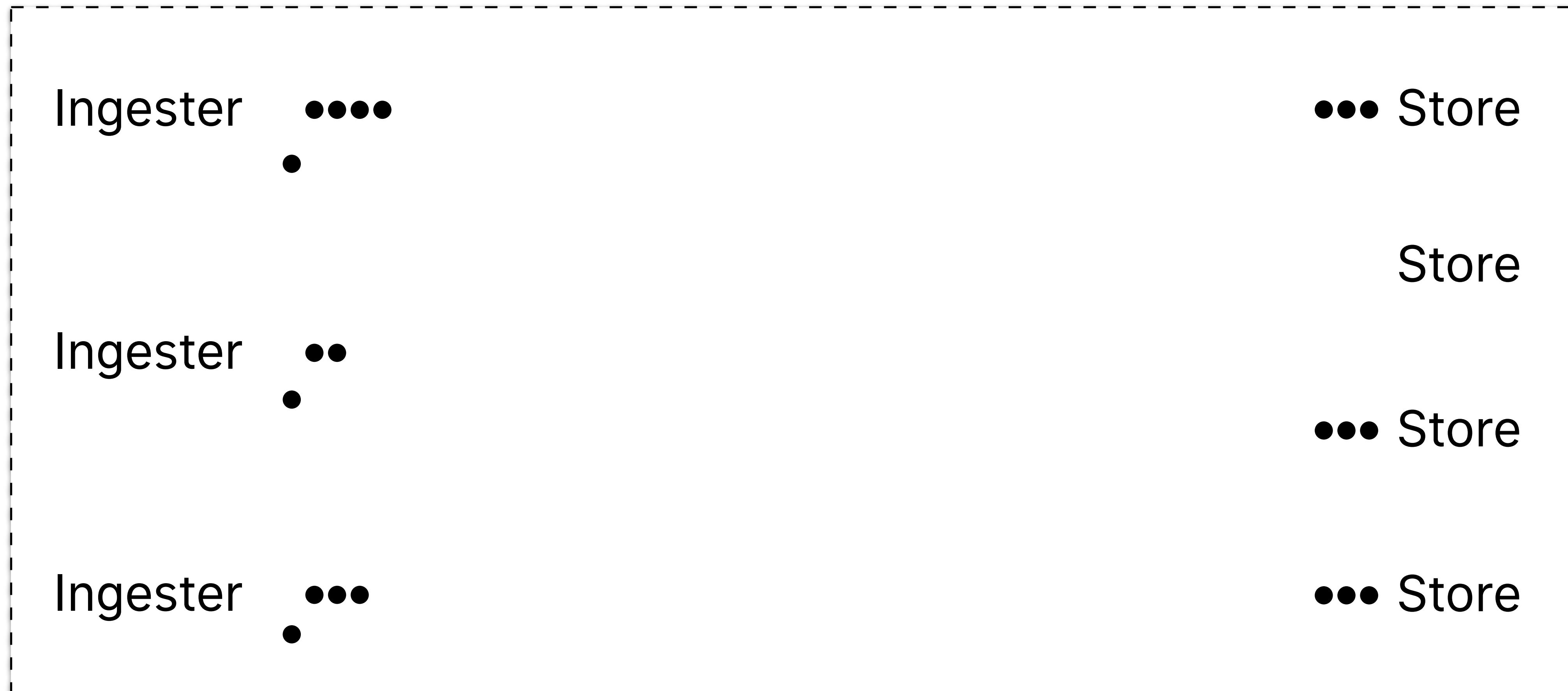


Failure during acks

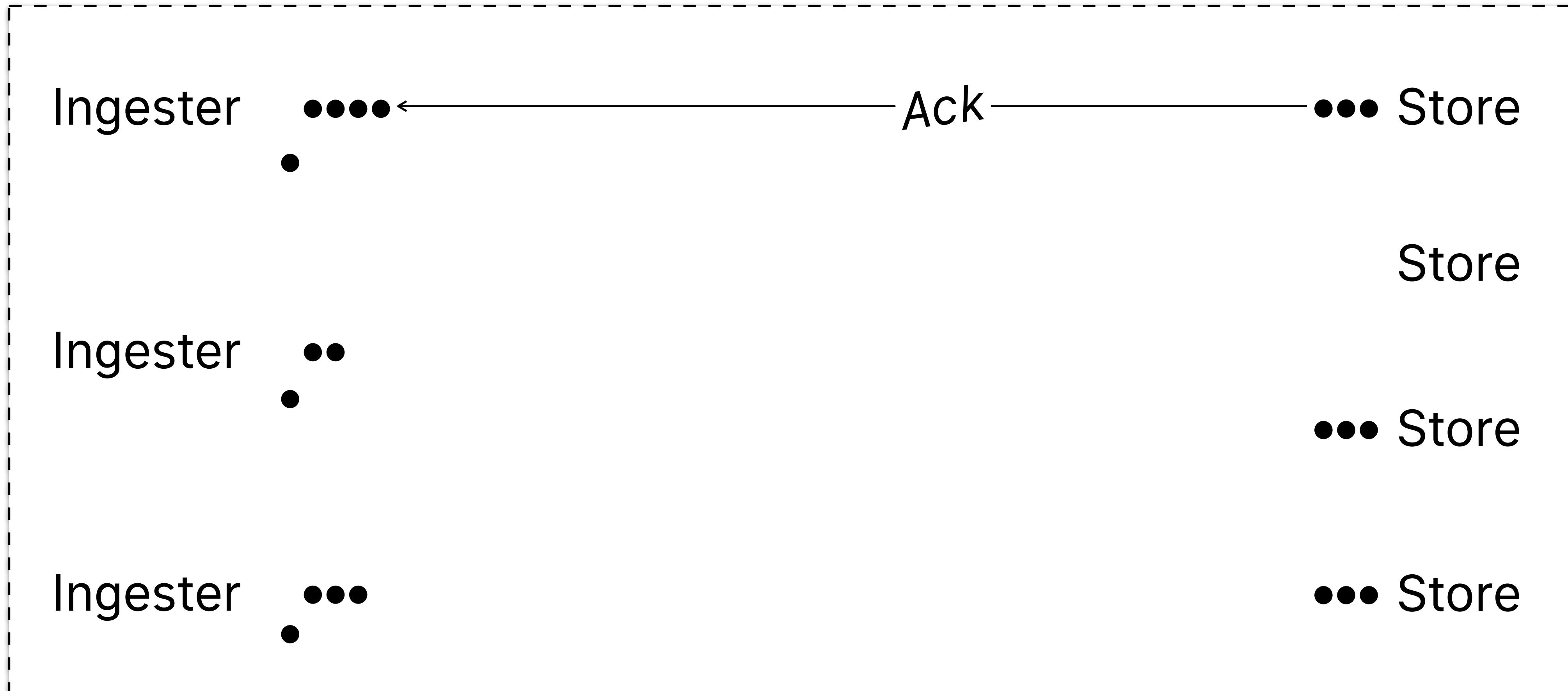
Just a degenerate case



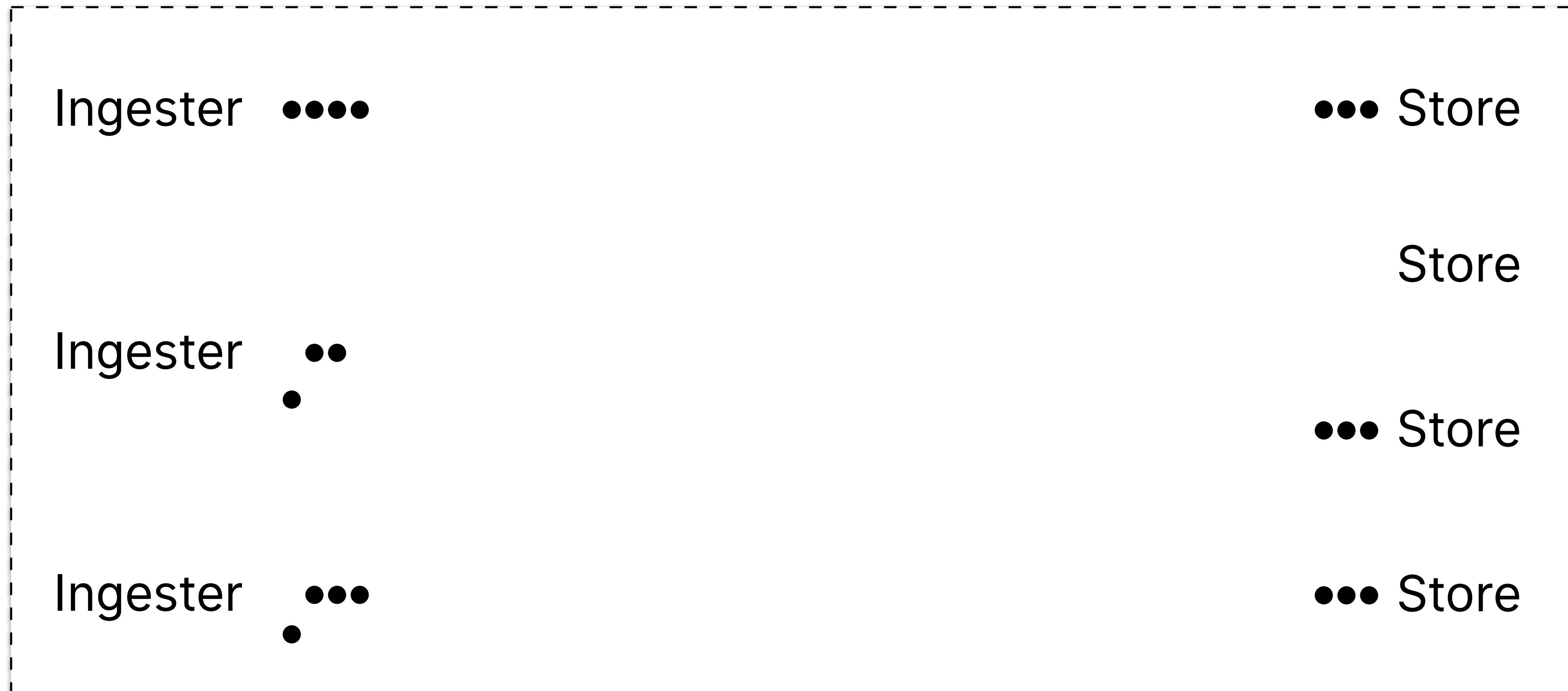
Failure during acks



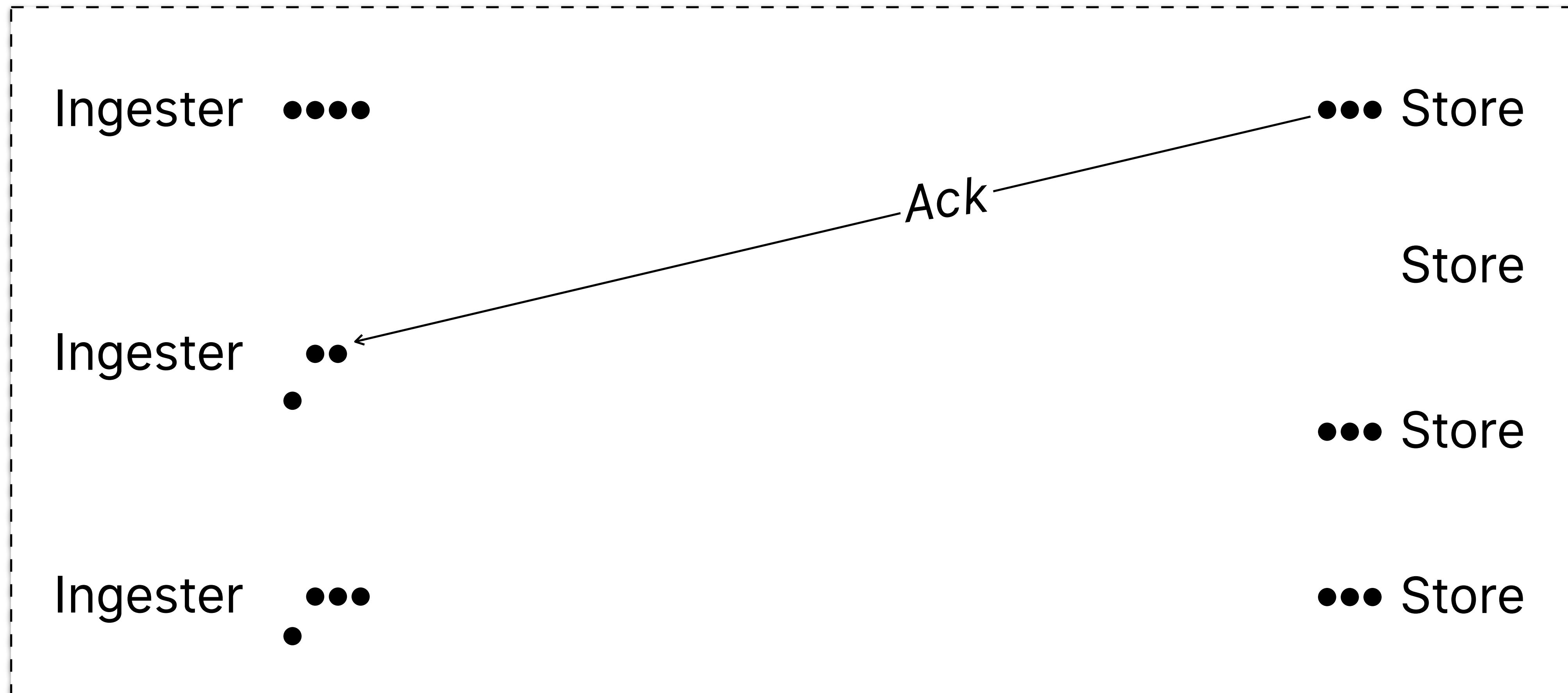
Failure during acks



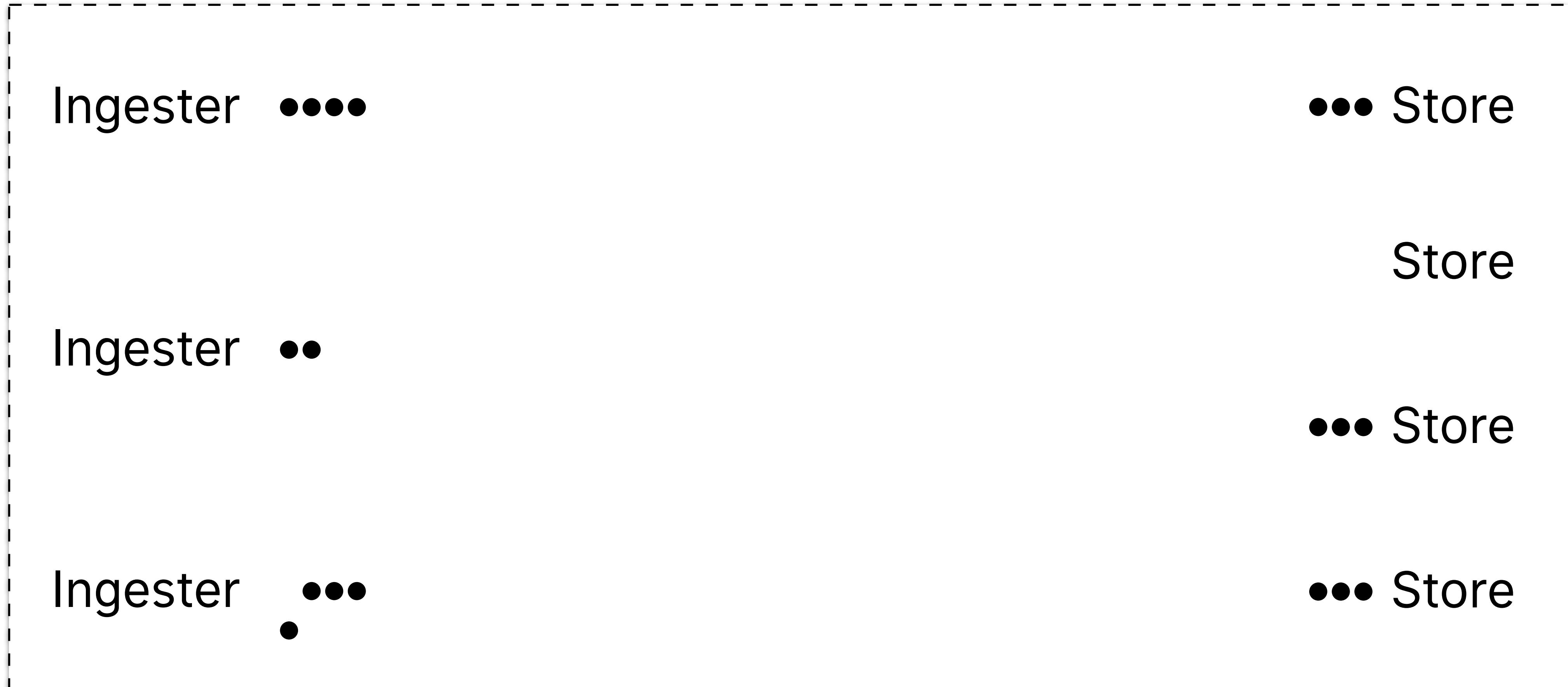
Failure during acks



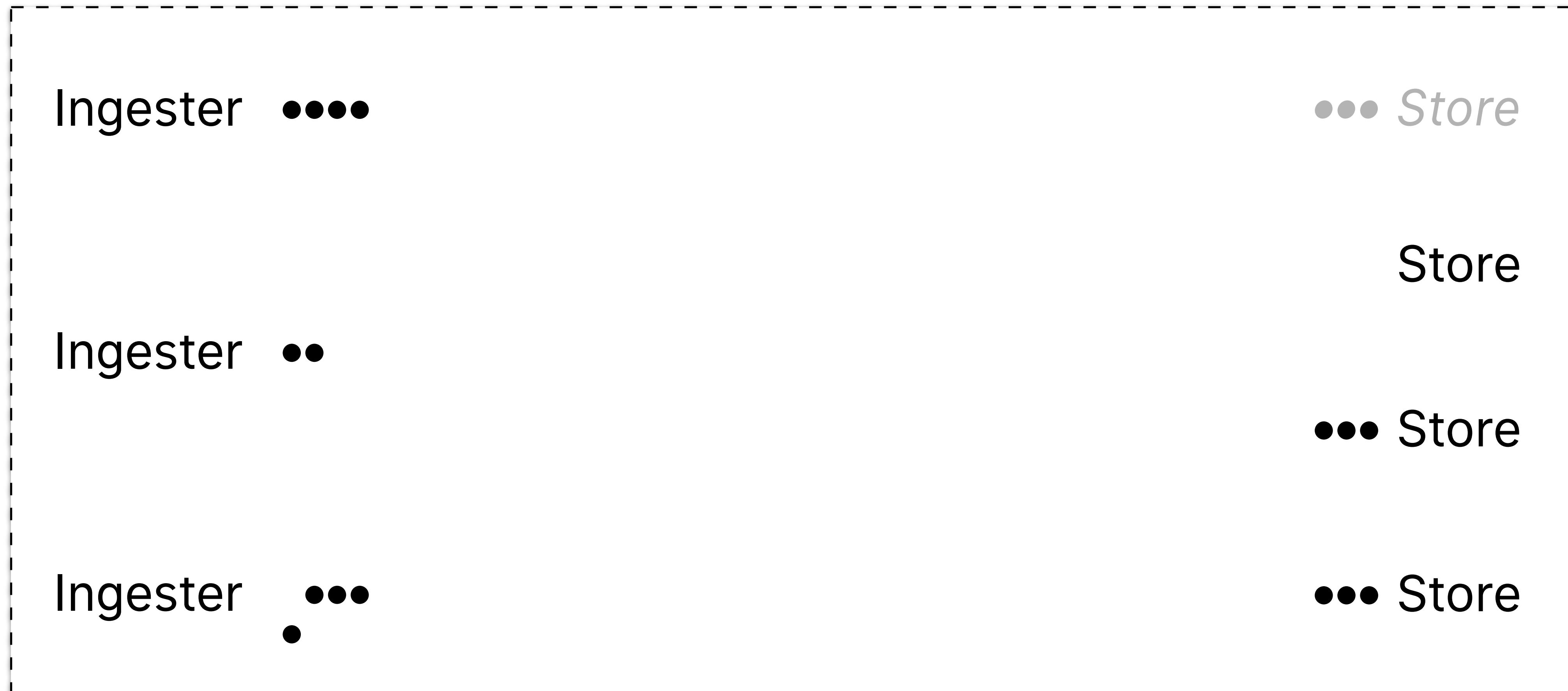
Failure during acks



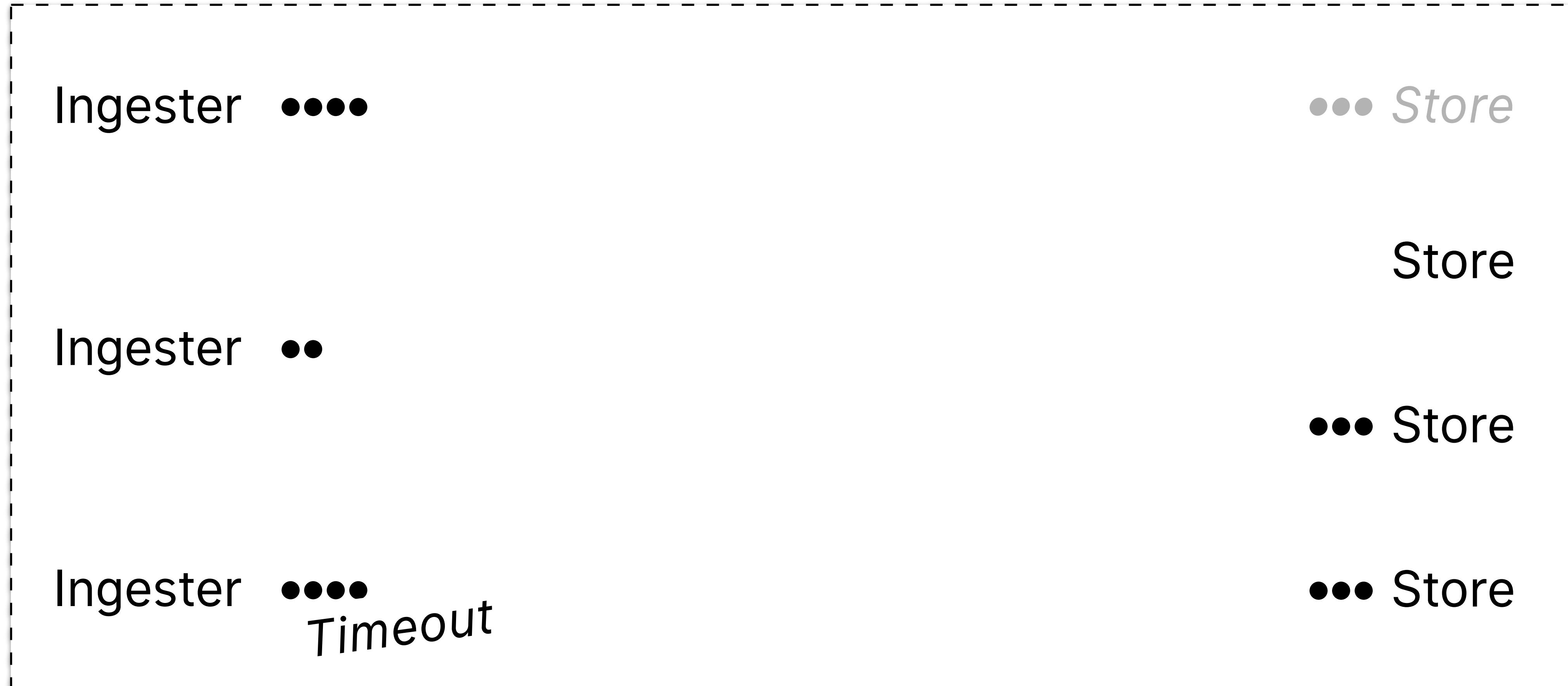
Failure during acks



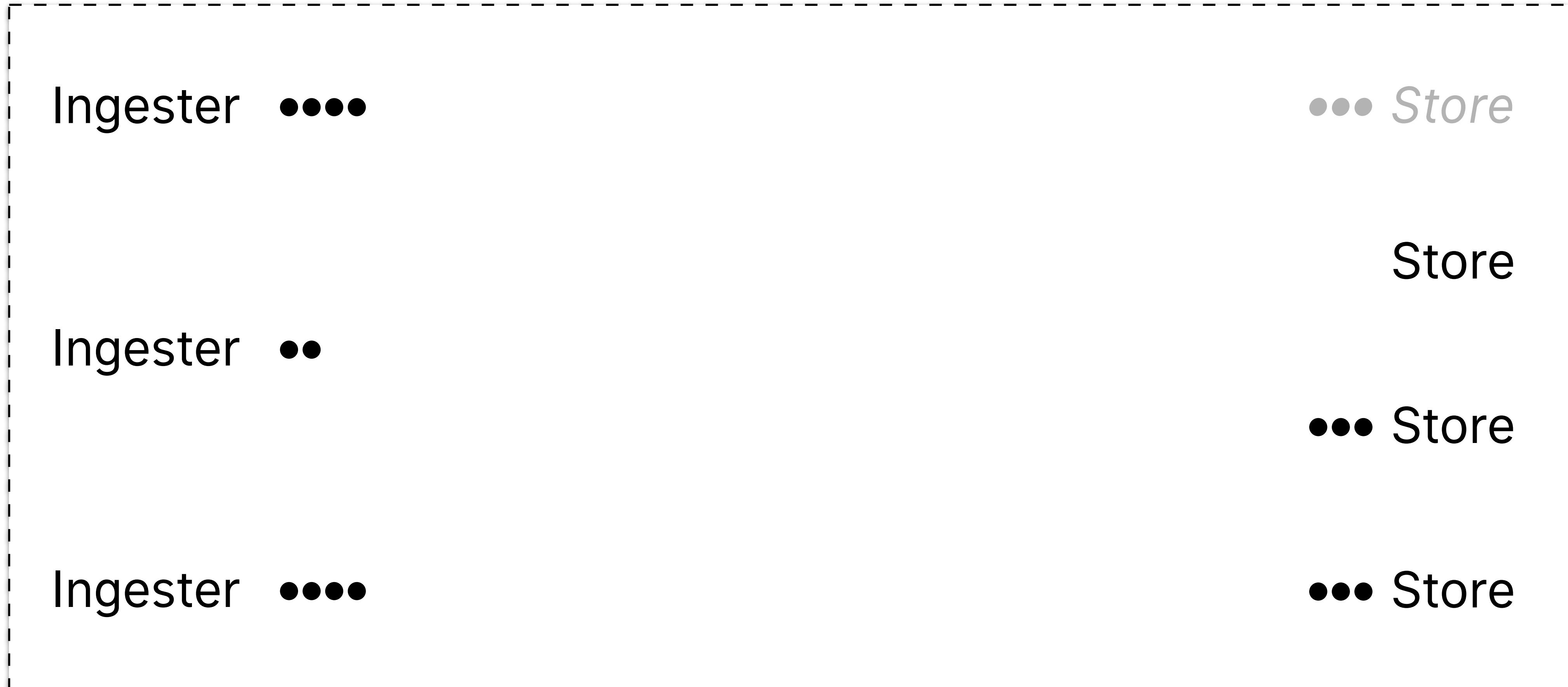
Failure during acks



Failure during acks



Failure during acks

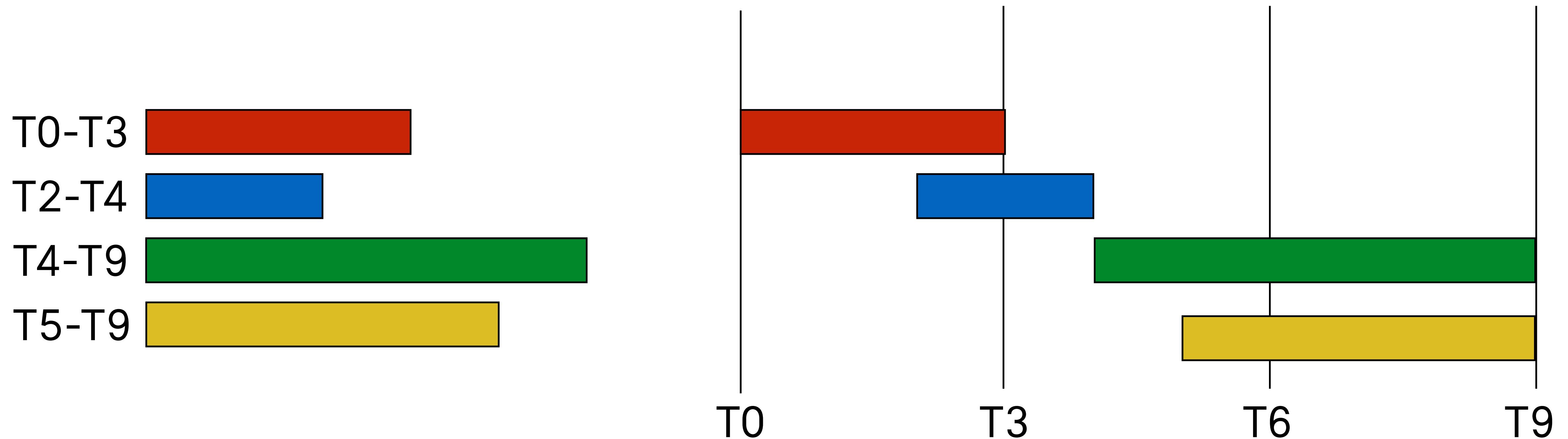


Compaction

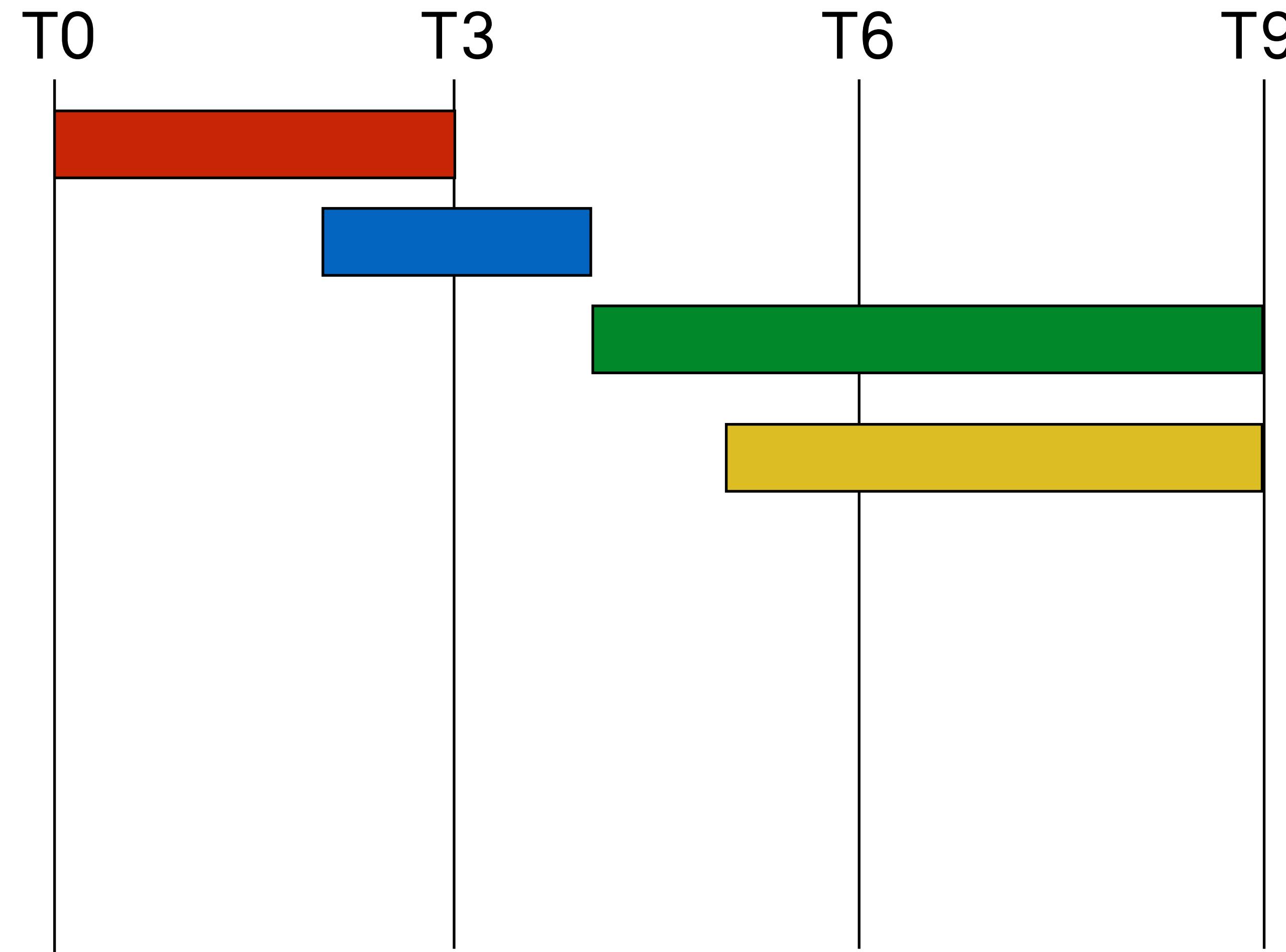
So stupid, so good



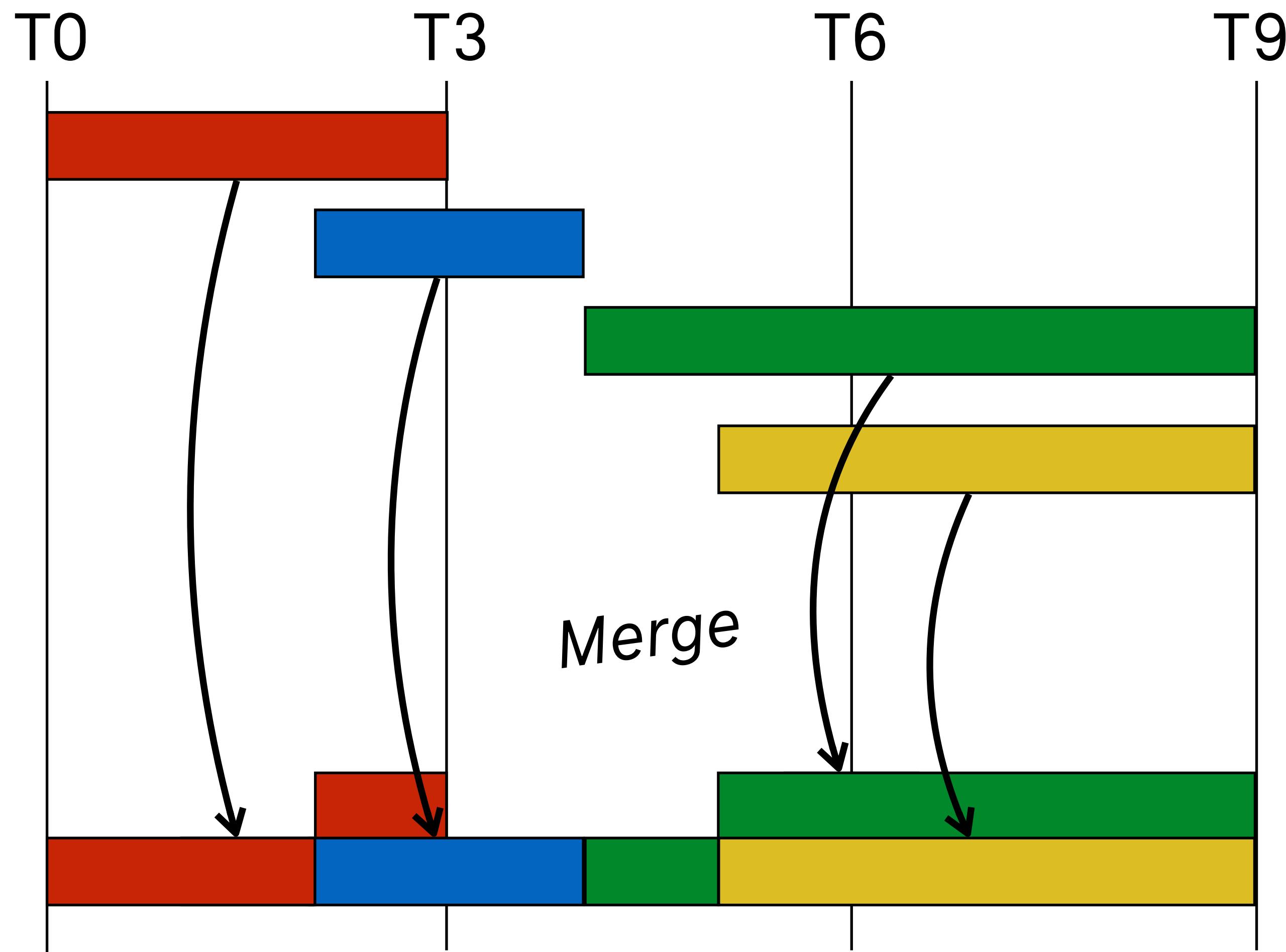
Segment files



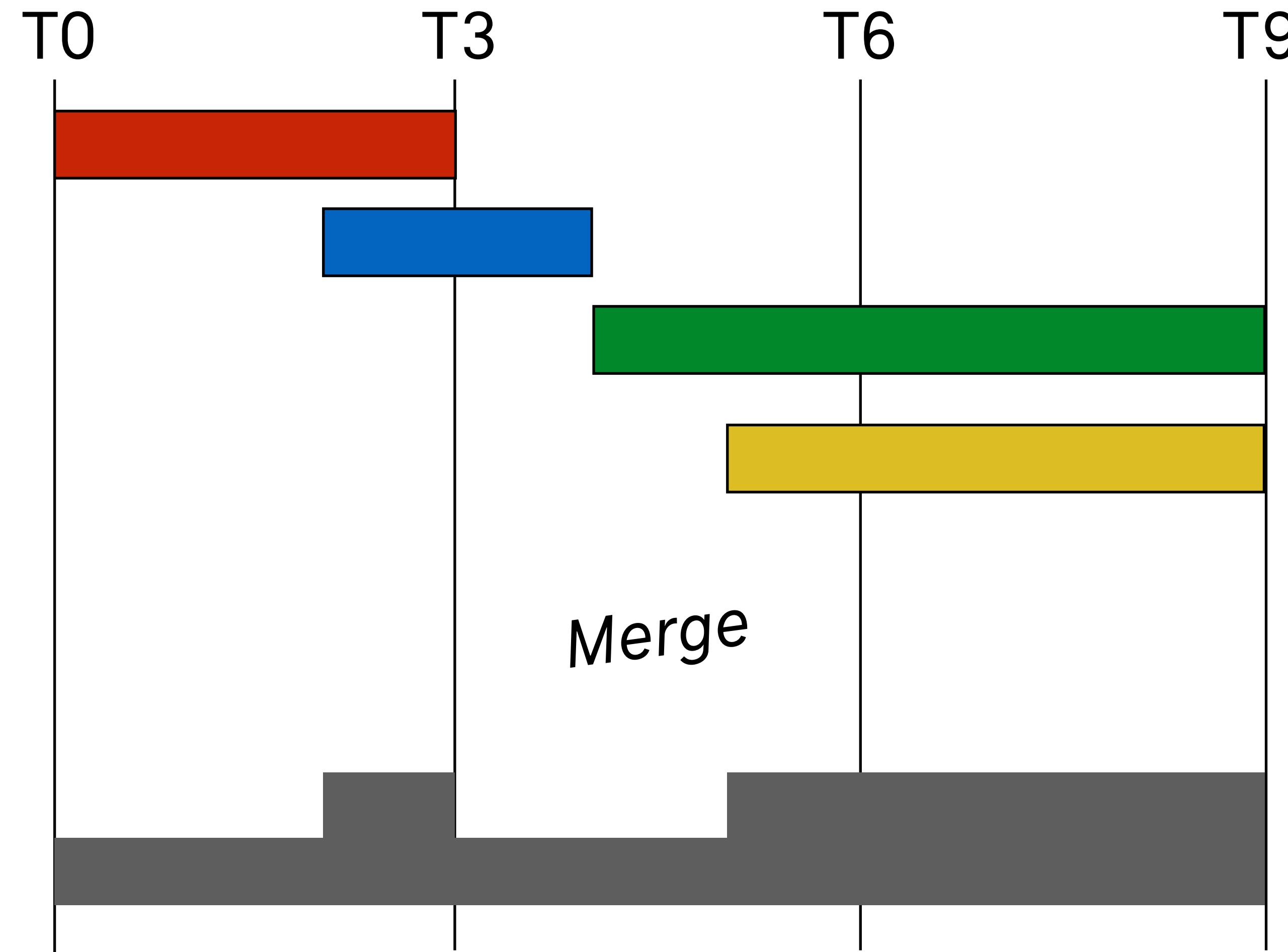
Segment files



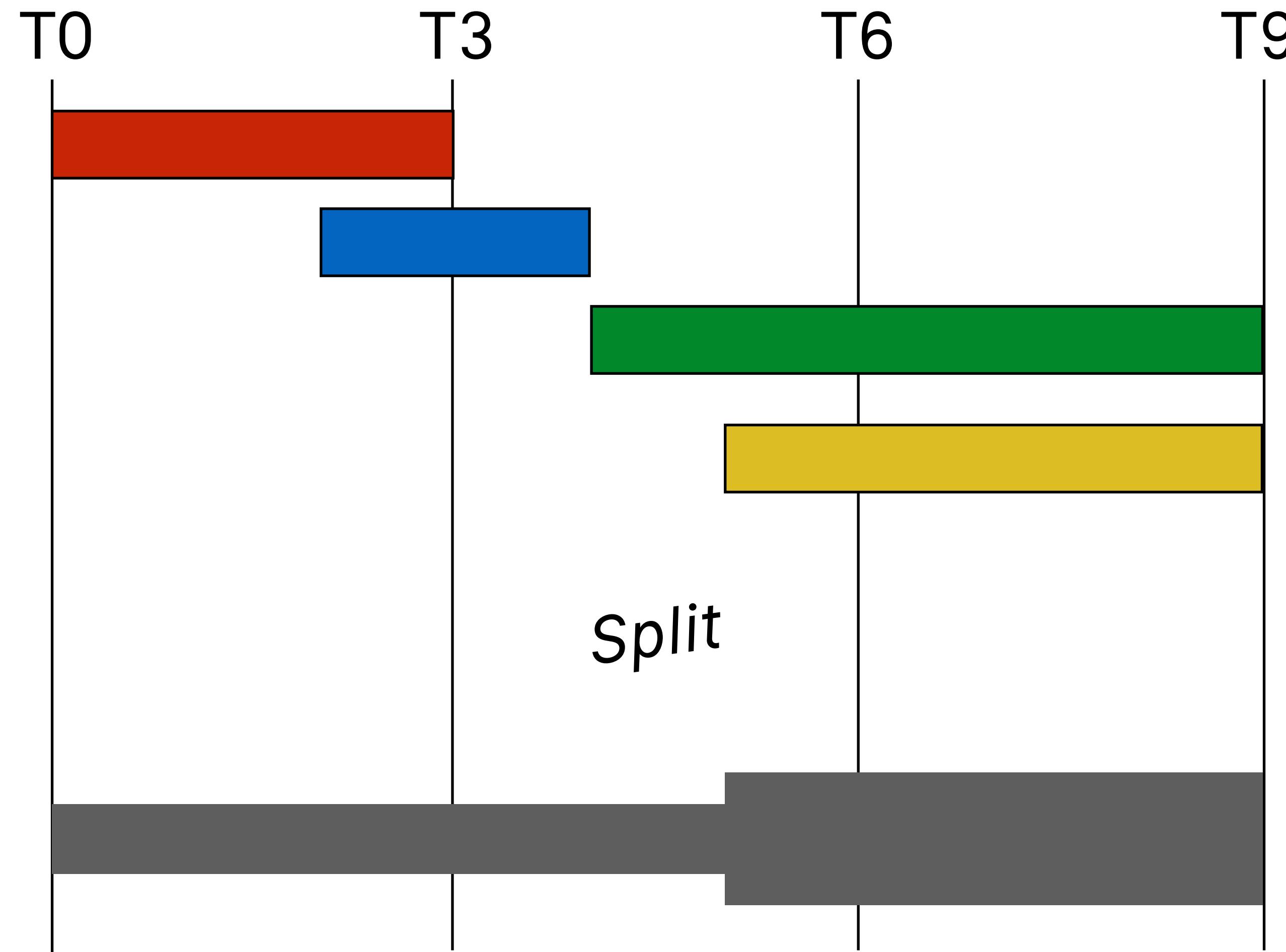
Segment files



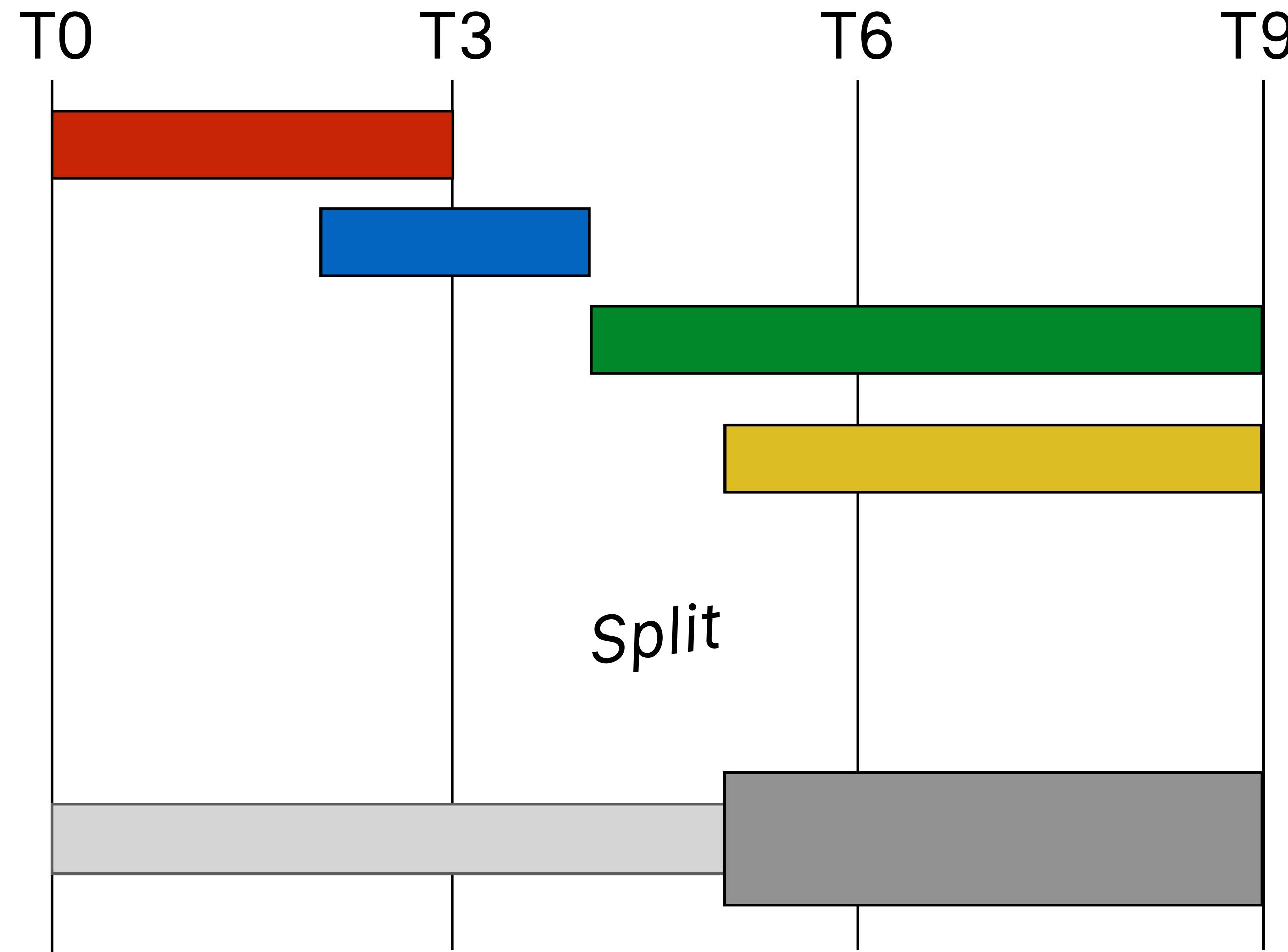
Segment files



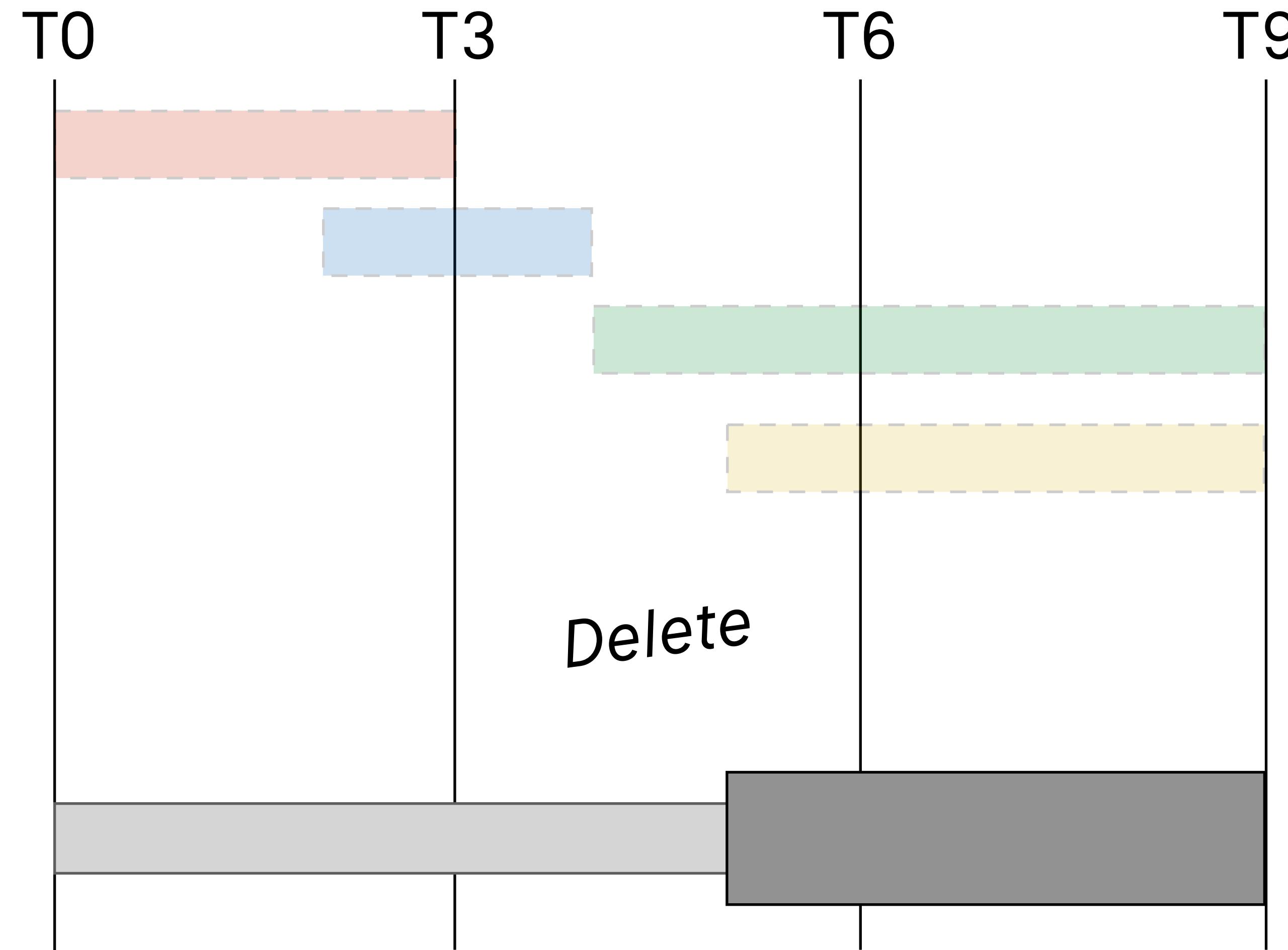
Segment files



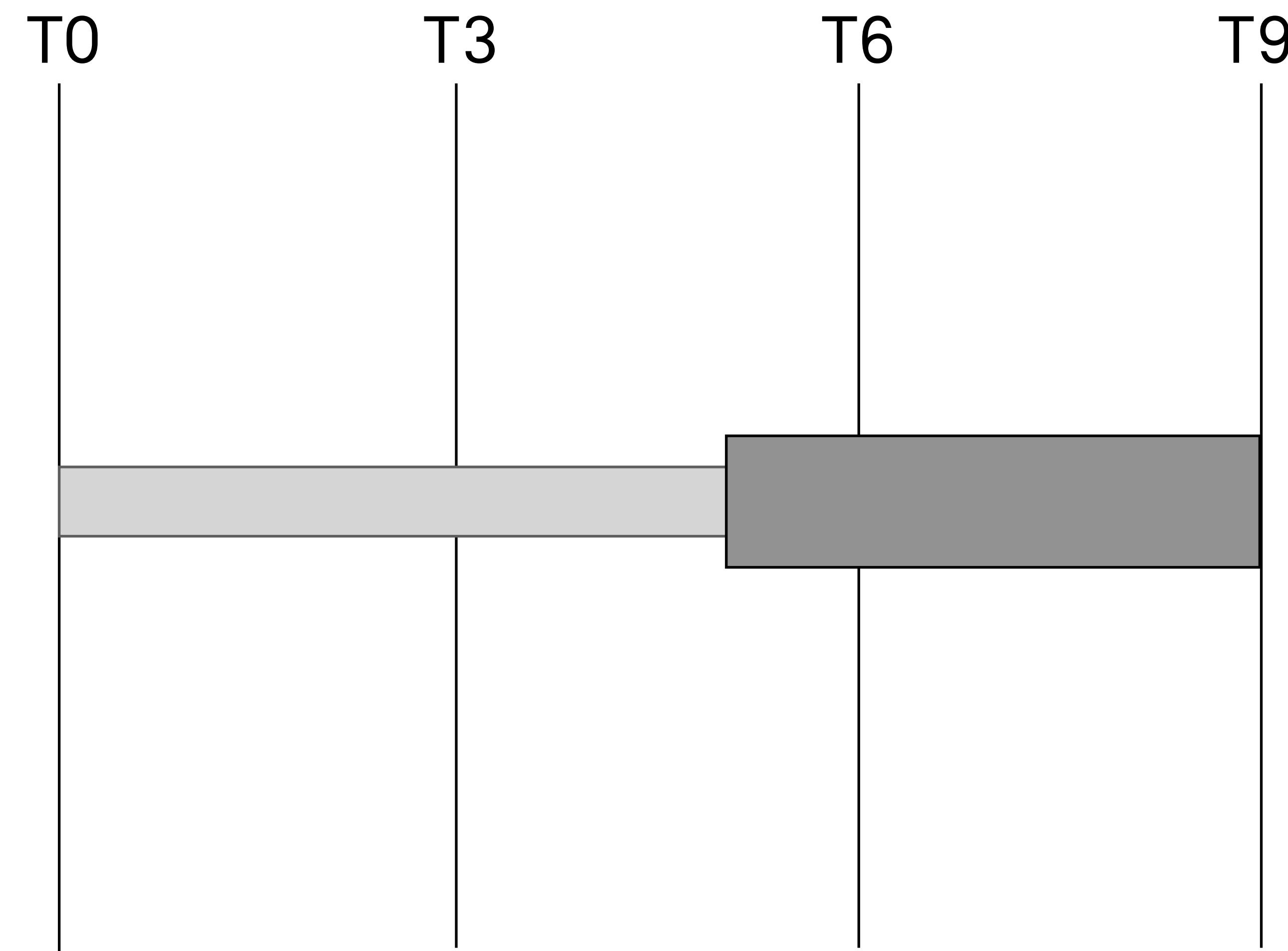
Segment files



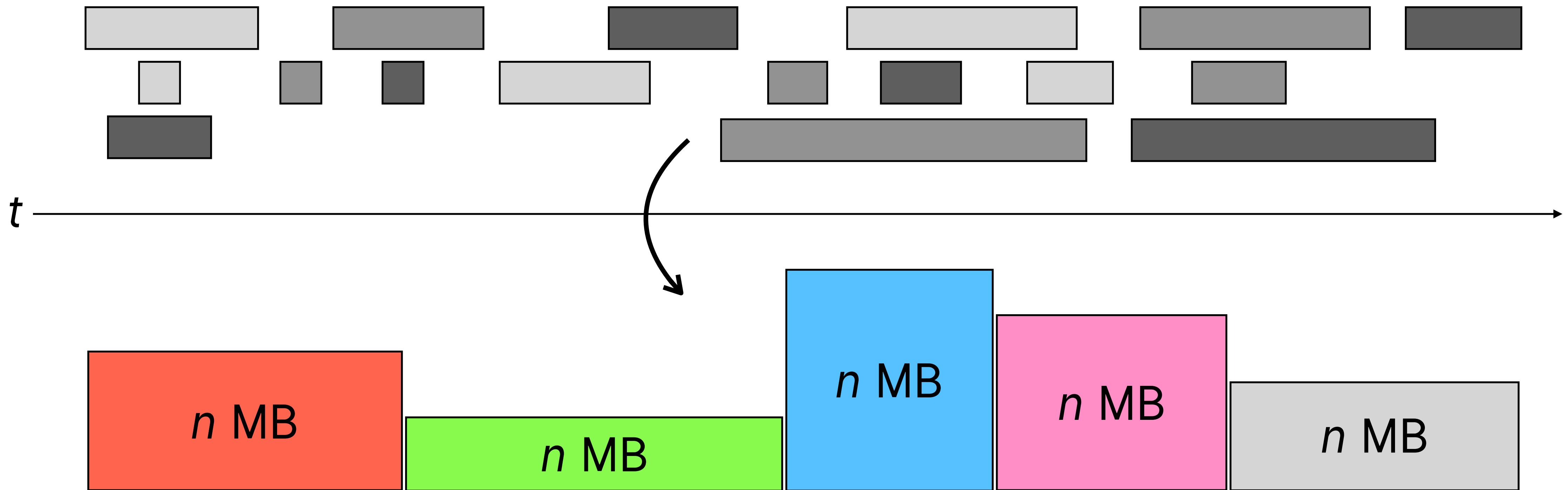
Segment files



Segment files



Goal of compaction



Adding capacity

Nothing fancy

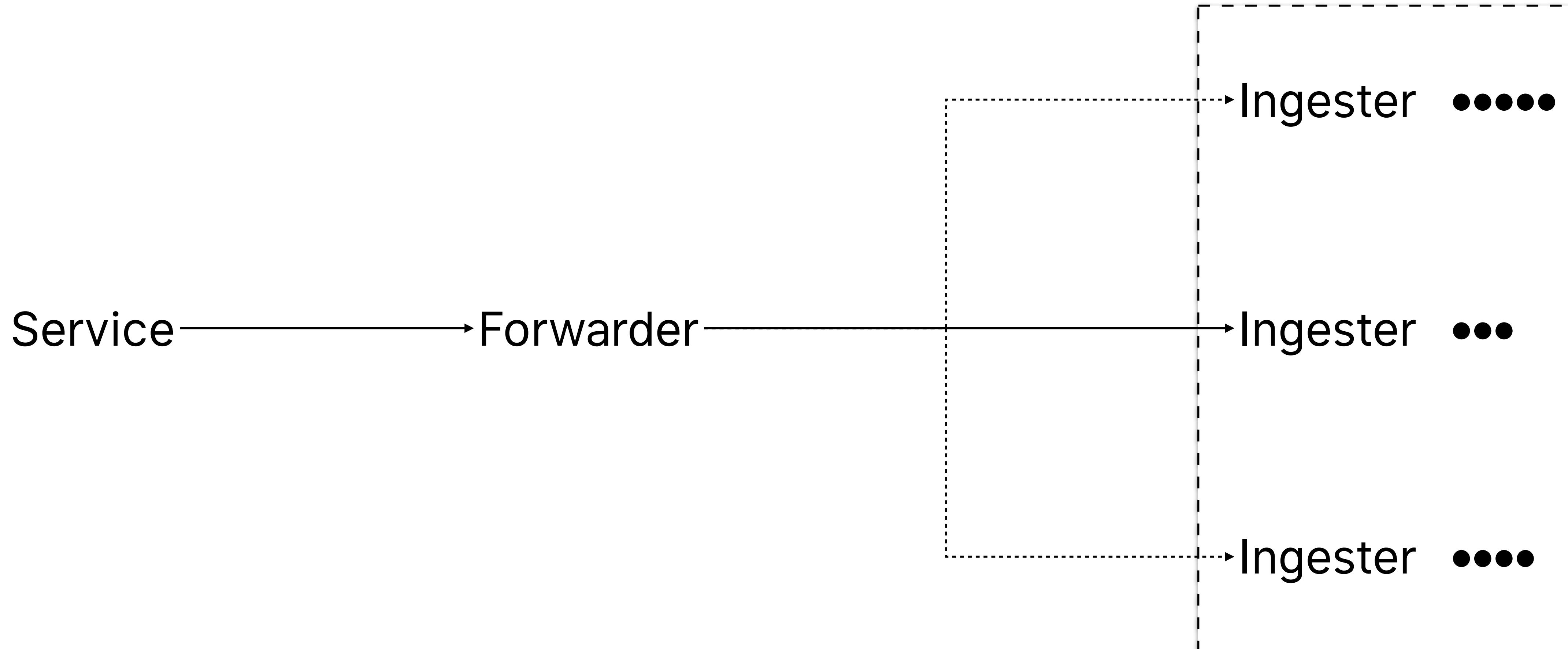


Adding ingest capacity

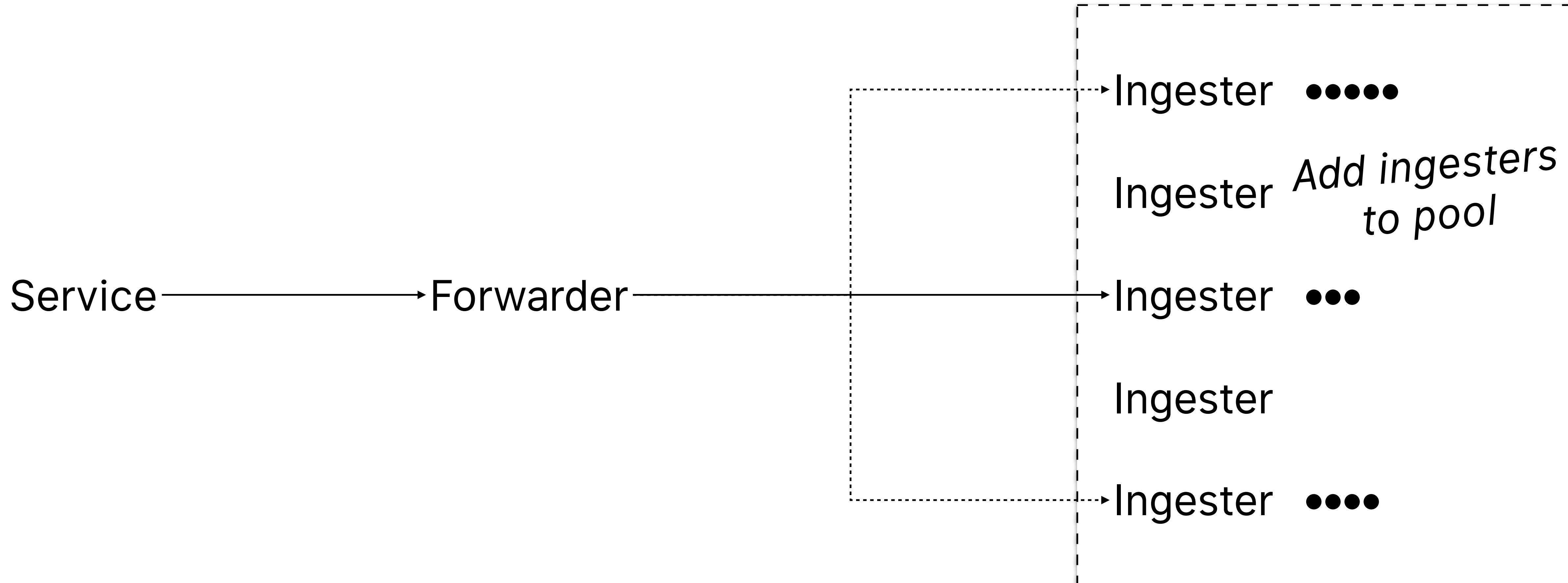
Literally just add ingestors



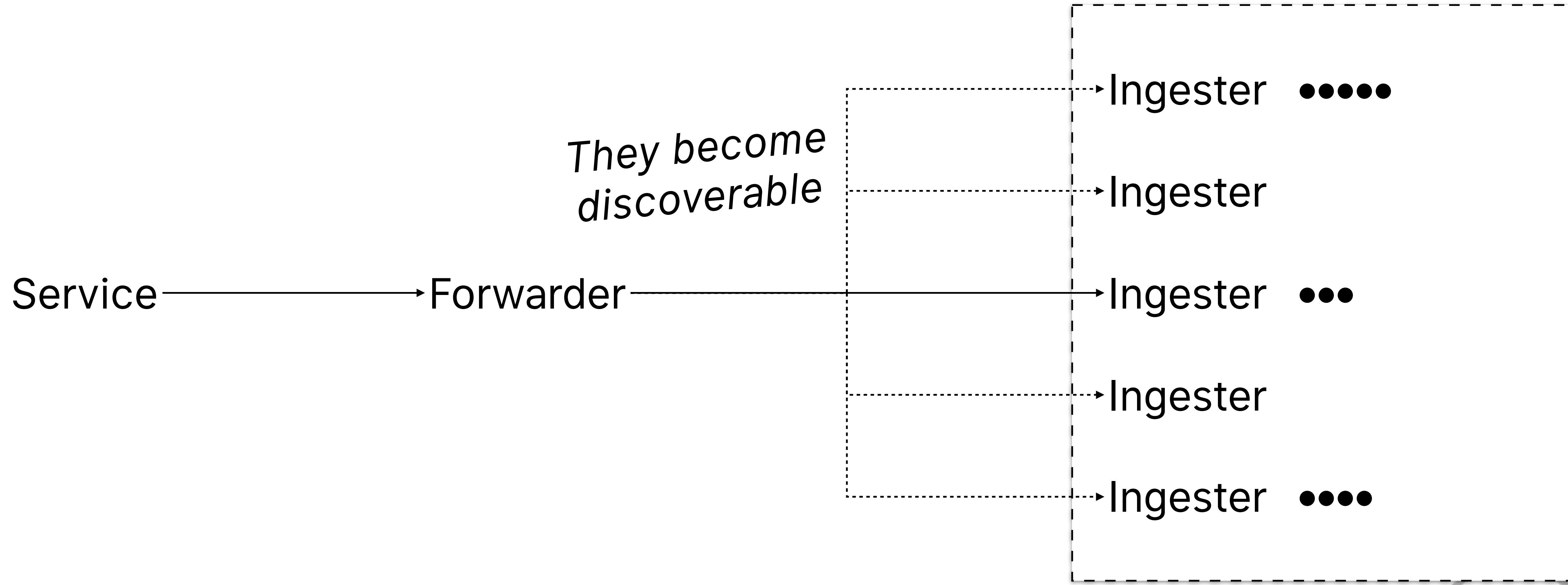
Adding ingest capacity



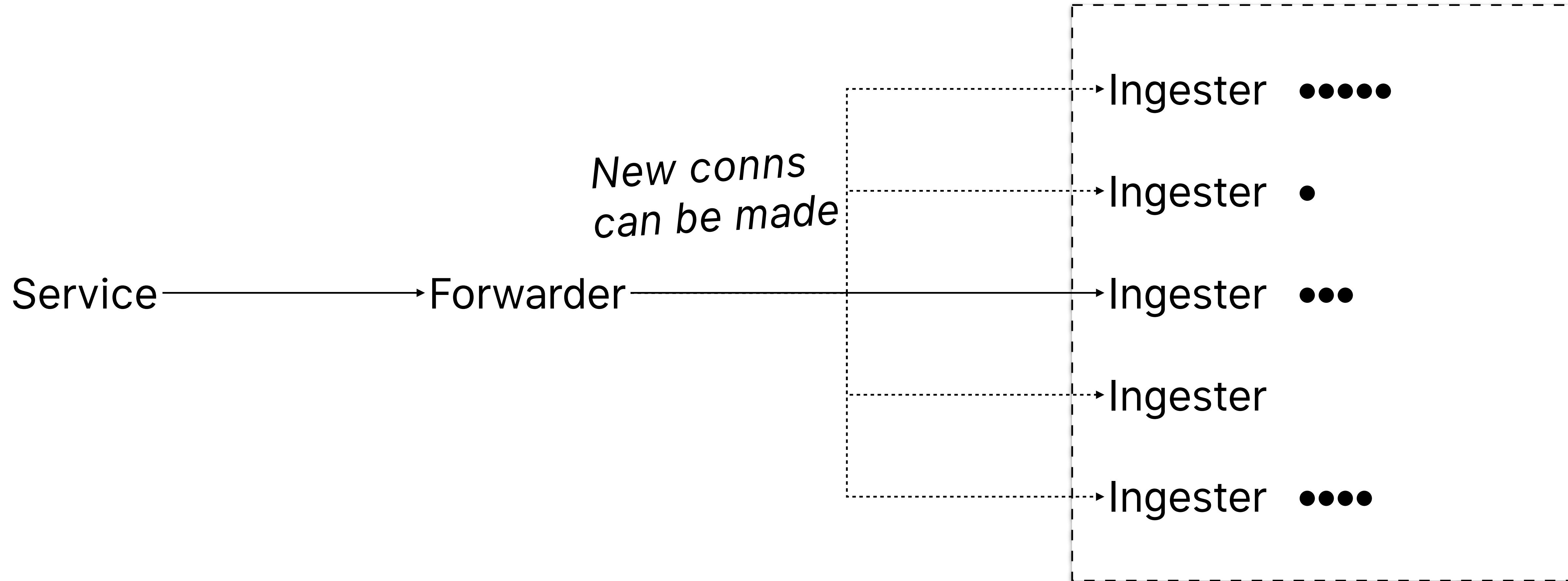
Adding ingest capacity



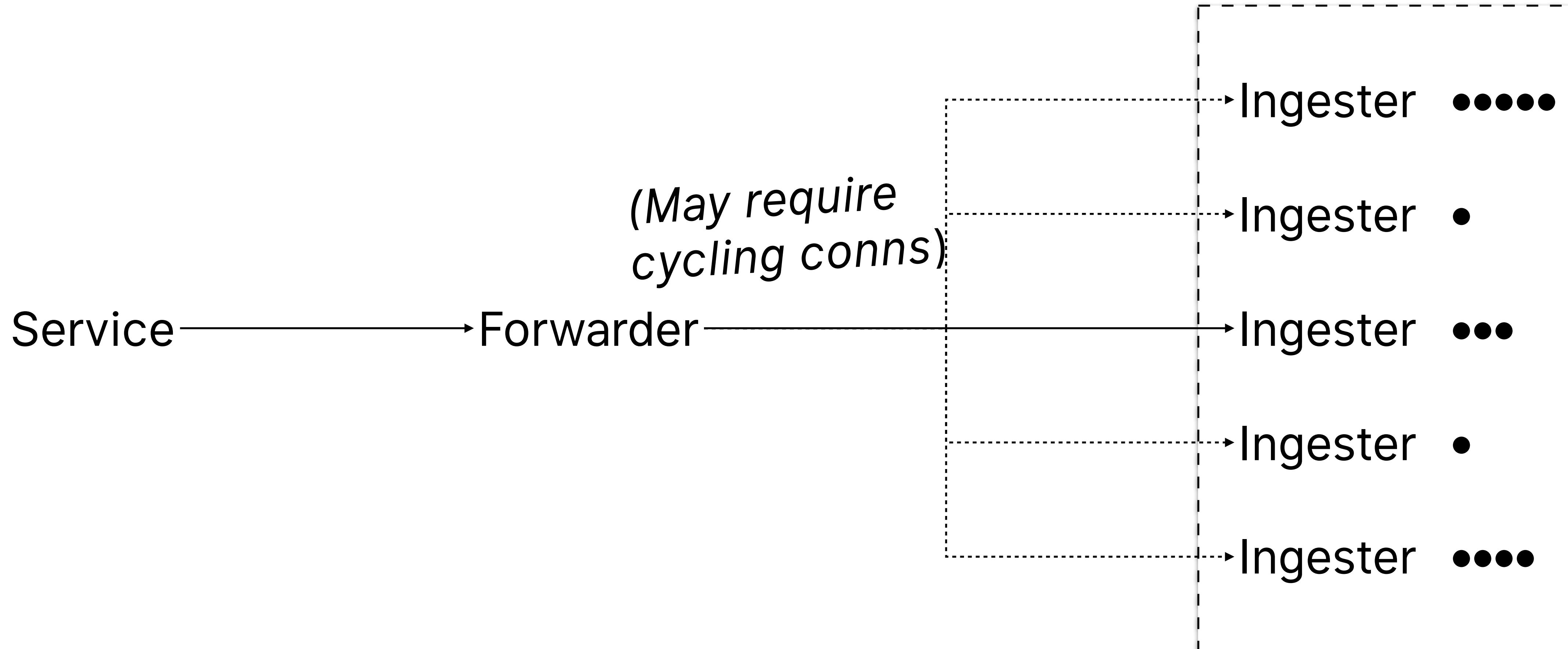
Adding ingest capacity



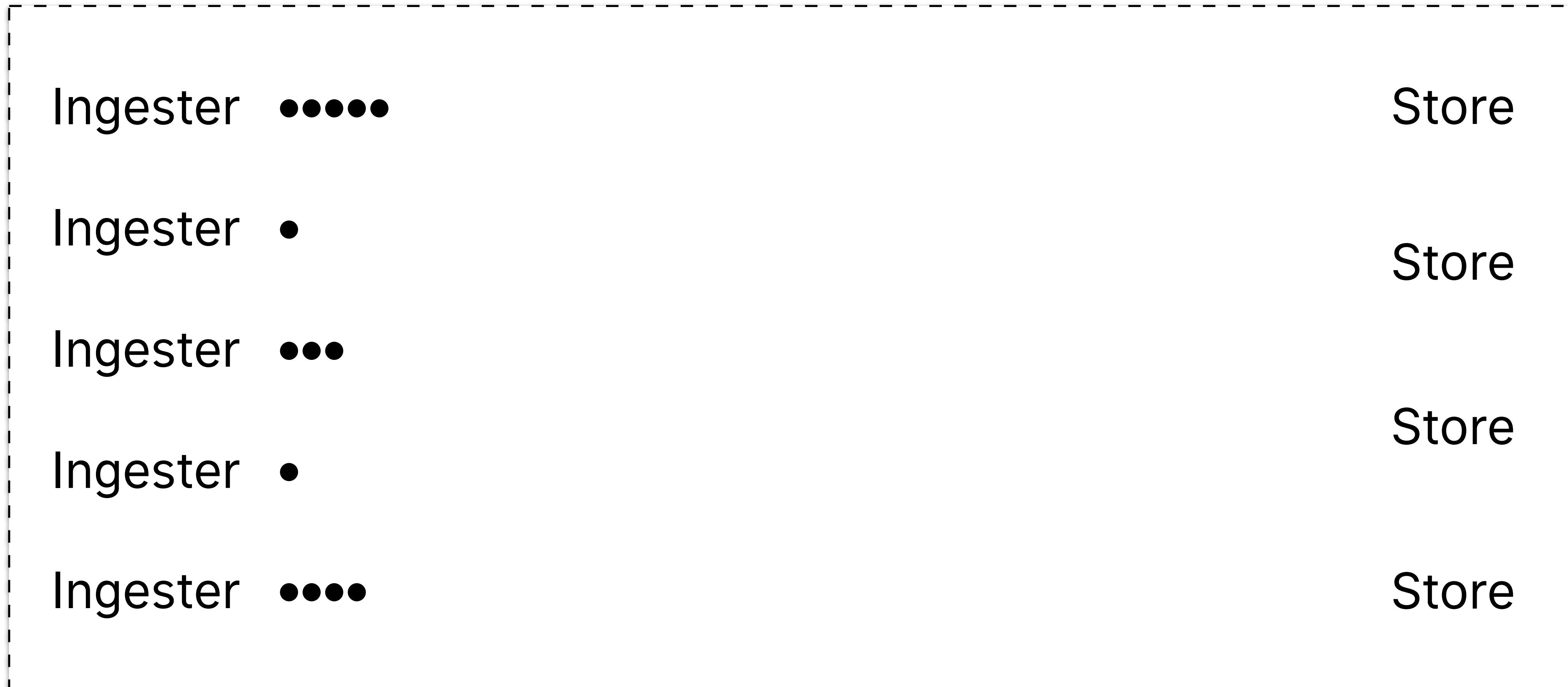
Adding ingest capacity



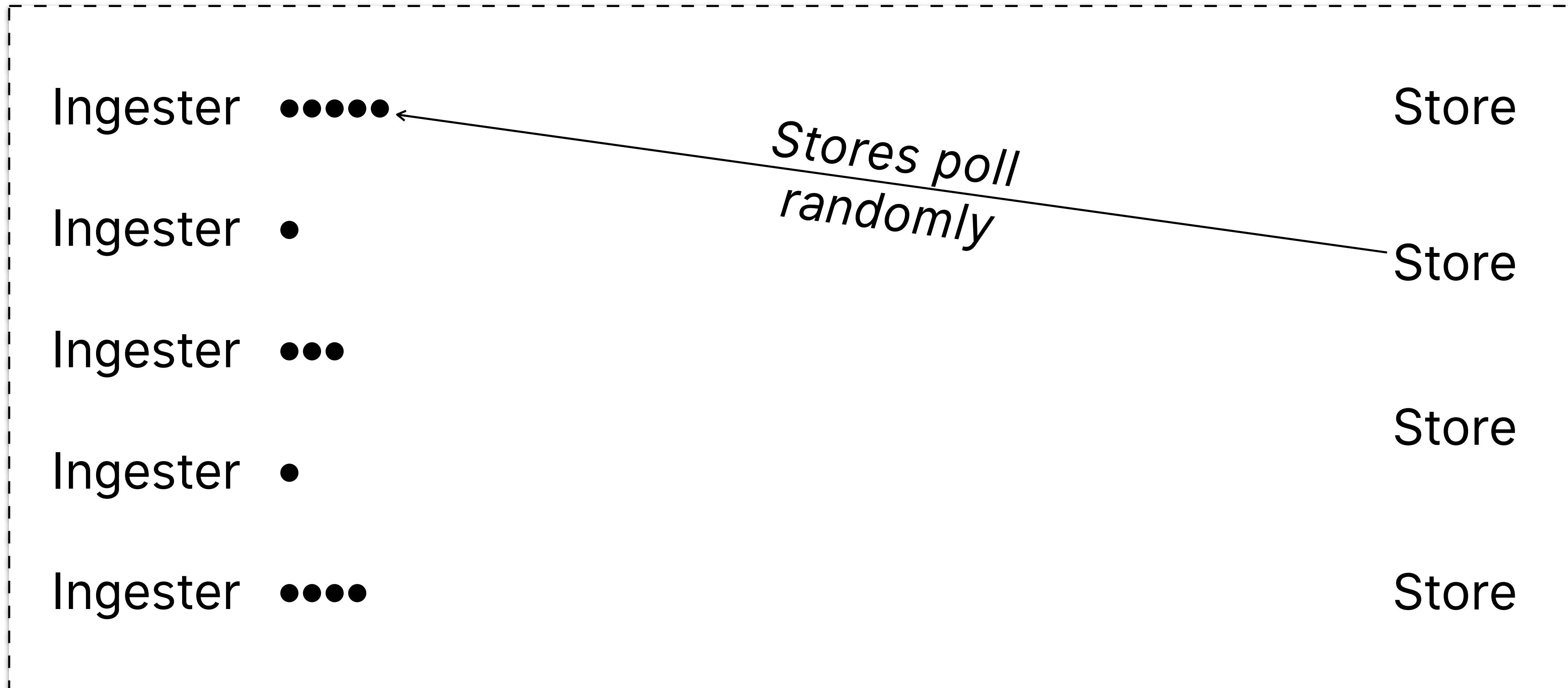
Adding ingest capacity



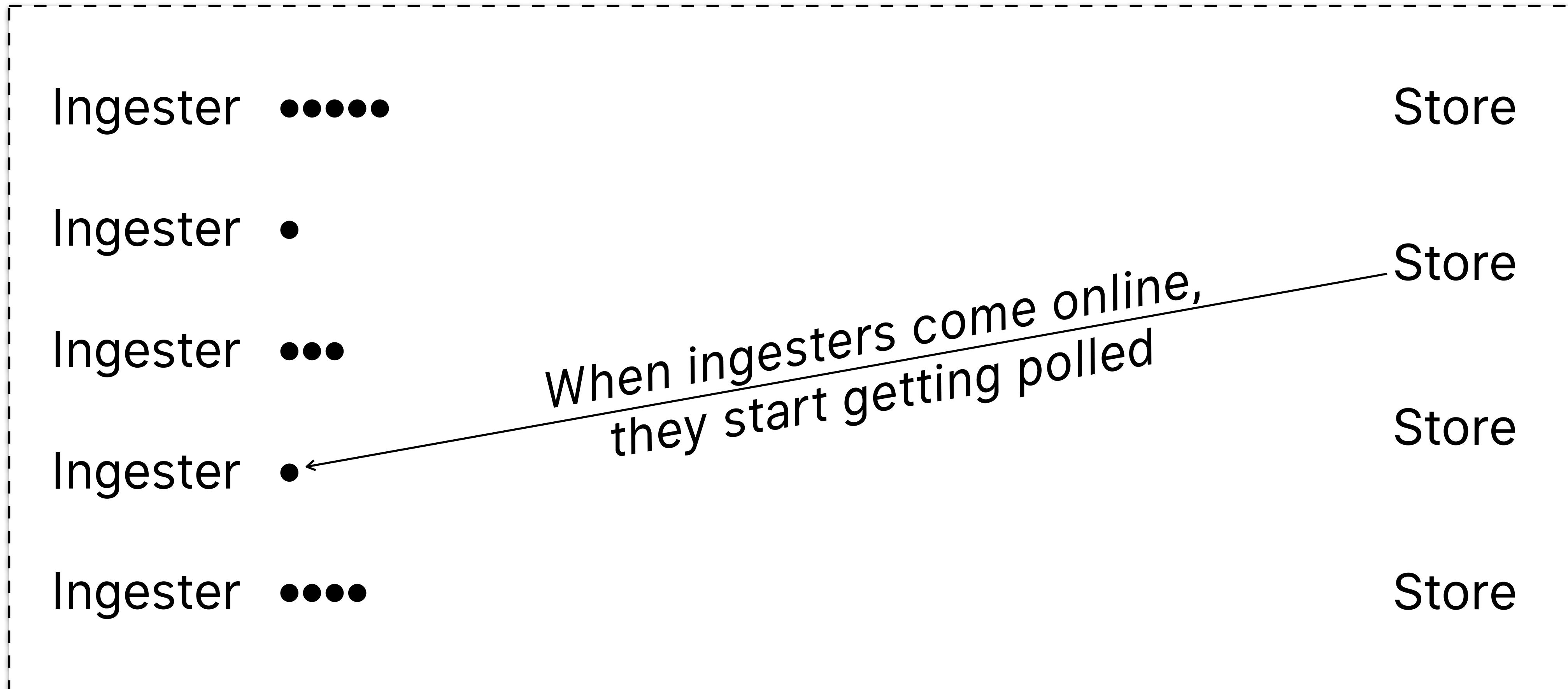
Adding ingest capacity



Adding ingest capacity



Adding ingest capacity



Adding query capacity

Add and wait



Adding query capacity



Adding query capacity



Adding query capacity

Ingestor •••••

Ingestor •

Ingestor •••

Ingestor •

Ingestor •••••

*And start getting equal
shares of segments*

••••• Store

● Store

••••••• Store

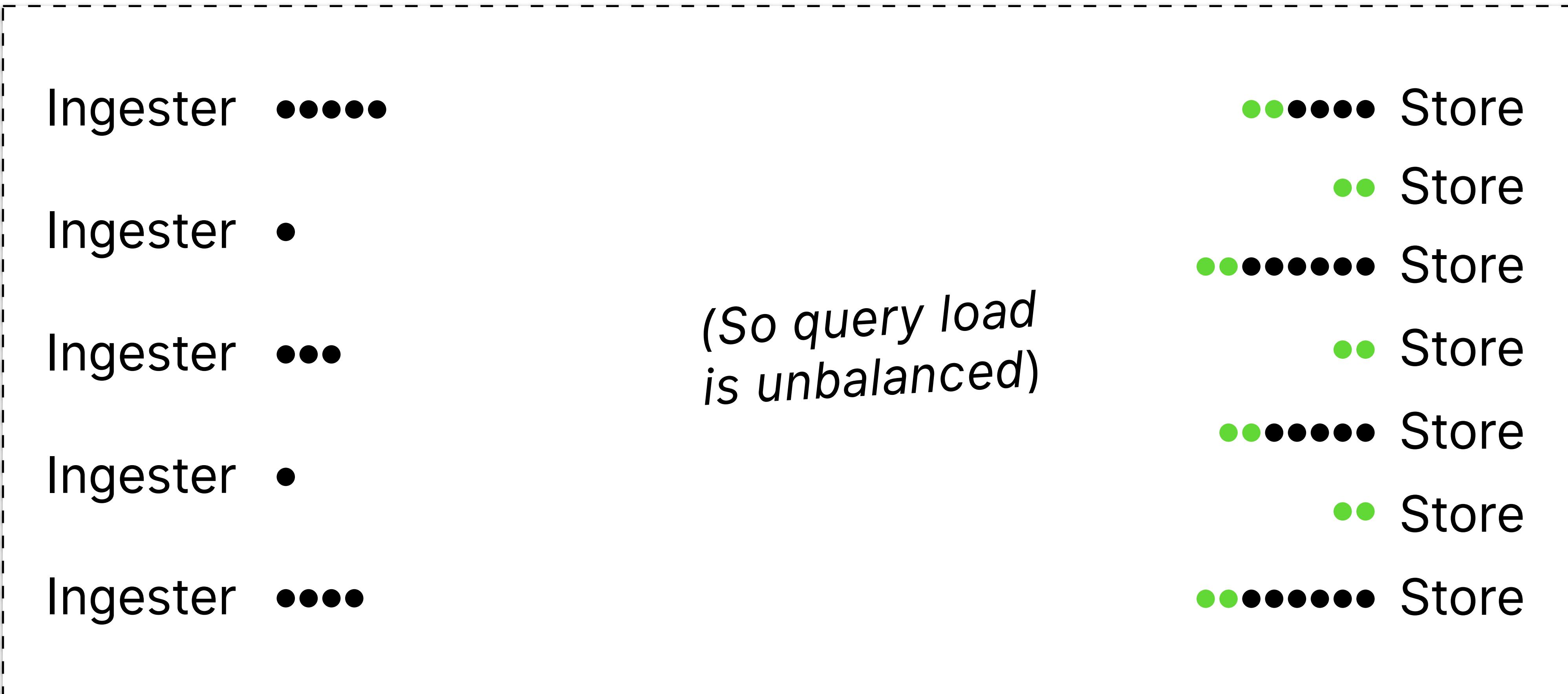
● Store

••••• Store

● Store

••••••• Store

Adding query capacity



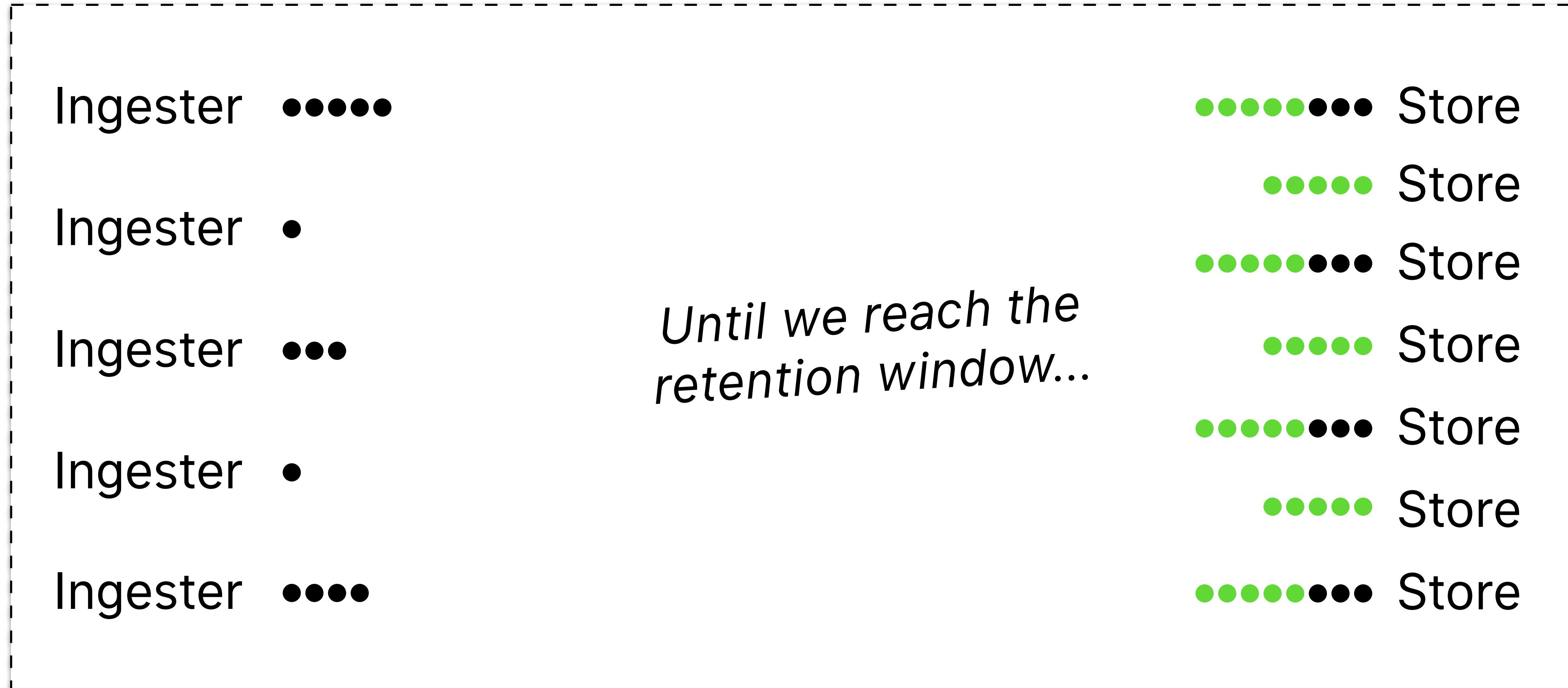
Adding query capacity



Adding query capacity



Adding query capacity



Adding query capacity

Ingestor •••••

Ingestor •

Ingestor •••

Ingestor •

Ingestor •••••

*Until we reach the
retention window...*

•••••••••• Store

••••••• Store

•••••••••• Store

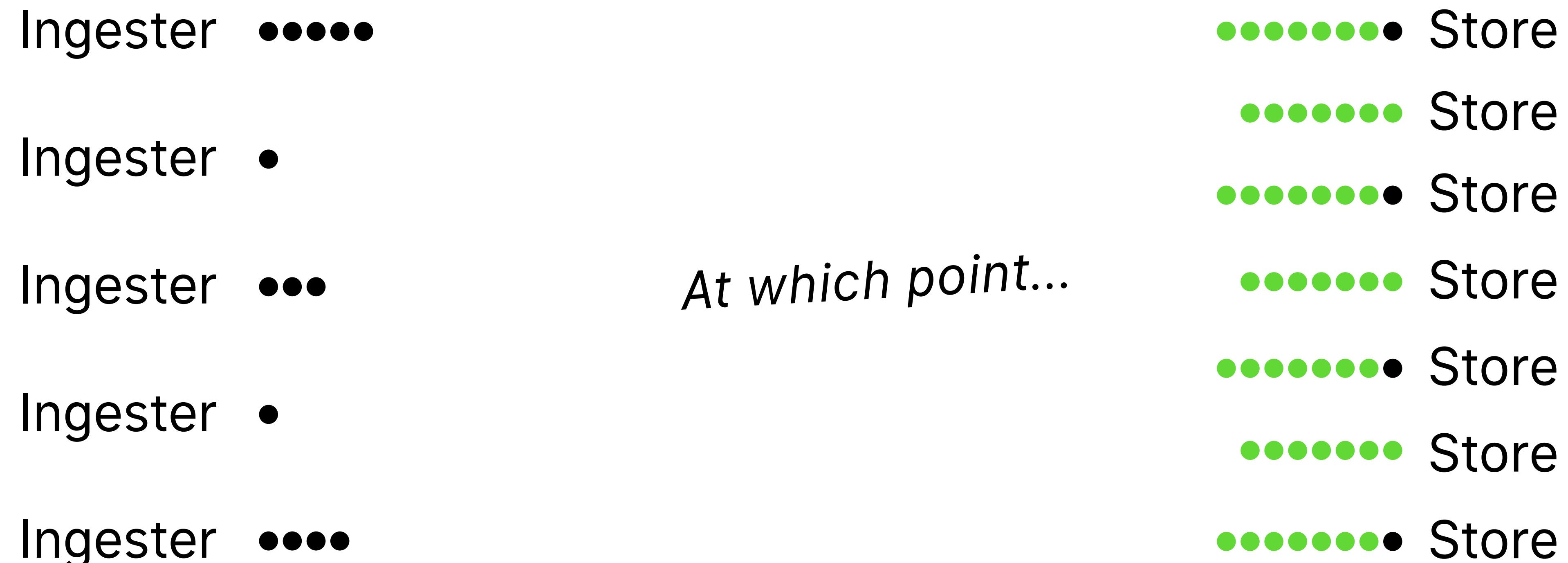
••••••• Store

•••••••••• Store

••••••• Store

•••••••••• Store

Adding query capacity



Adding query capacity

Ingestor ••••

Ingestor •

Ingestor •••

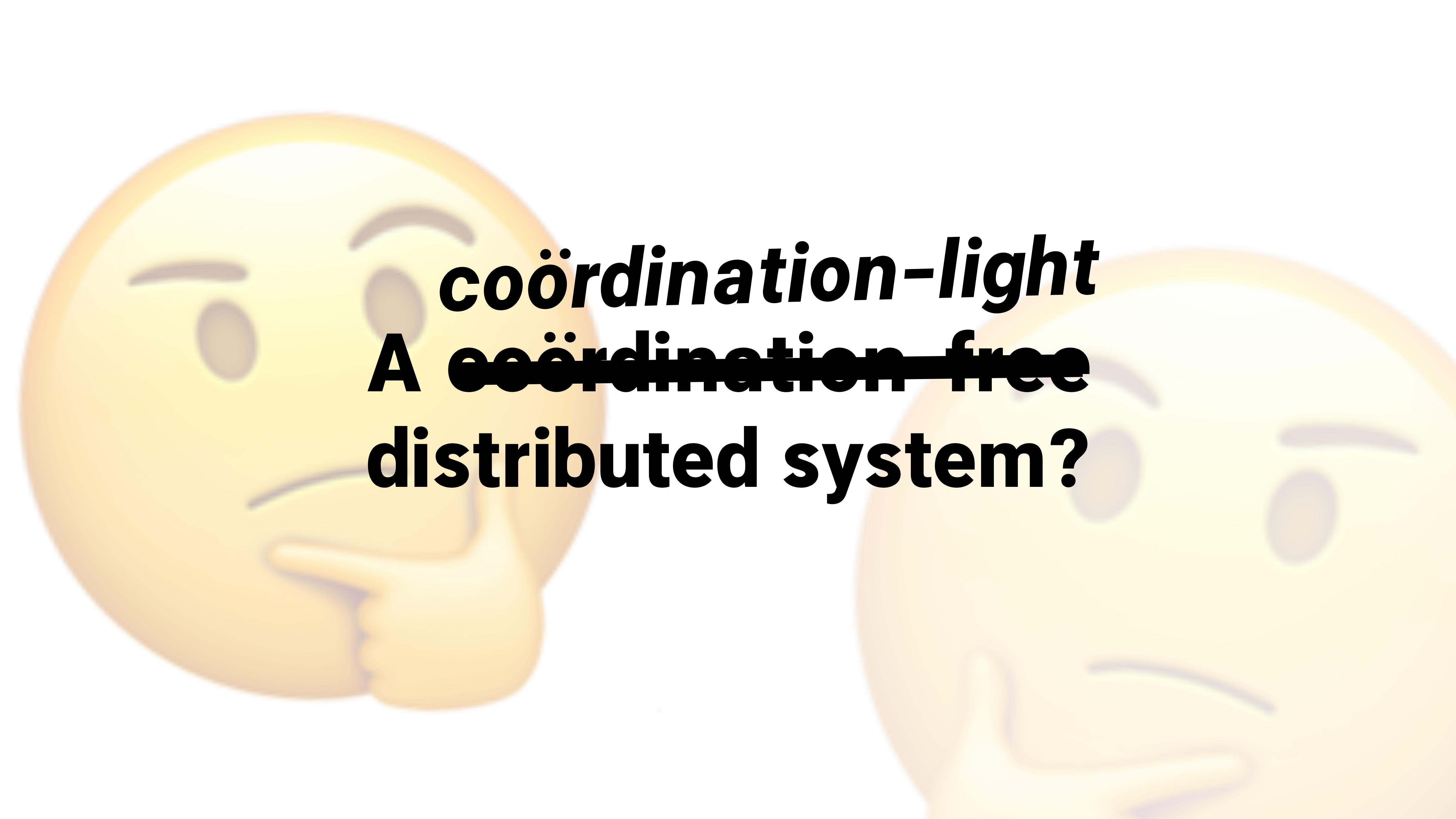
Ingestor •

Ingestor ••••

We're co-equal again

••••••• Store





coördination-light
A coördination-free
distributed system?

Conclusions

Solve smaller problems

It's OK to write code

Truth is unknowable



I crave social validation
Follow me on Twitter
@peterbourgon

OK Log 

Distributed and
coördination-*light*
logging