

# Incrementally Updateable Honey Password Vaults

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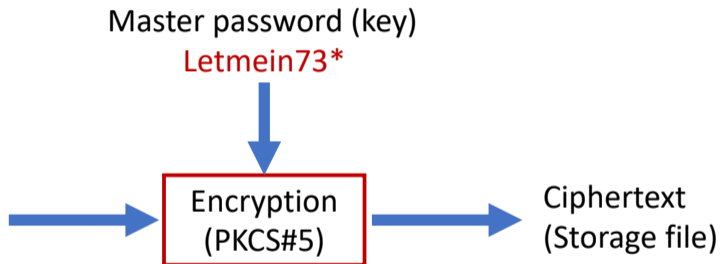


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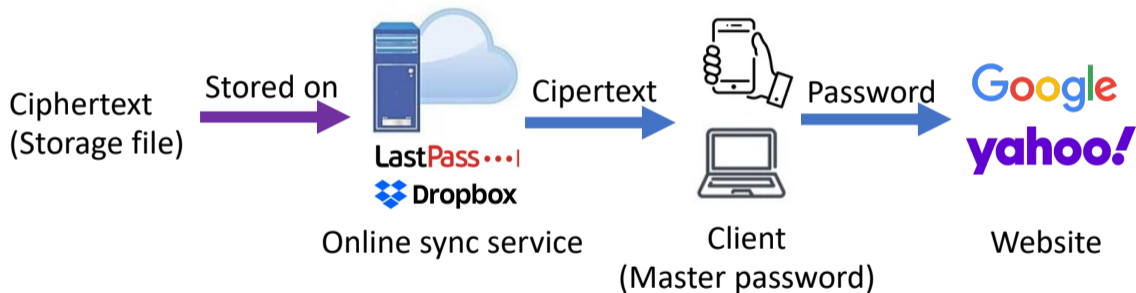


# Password vaults

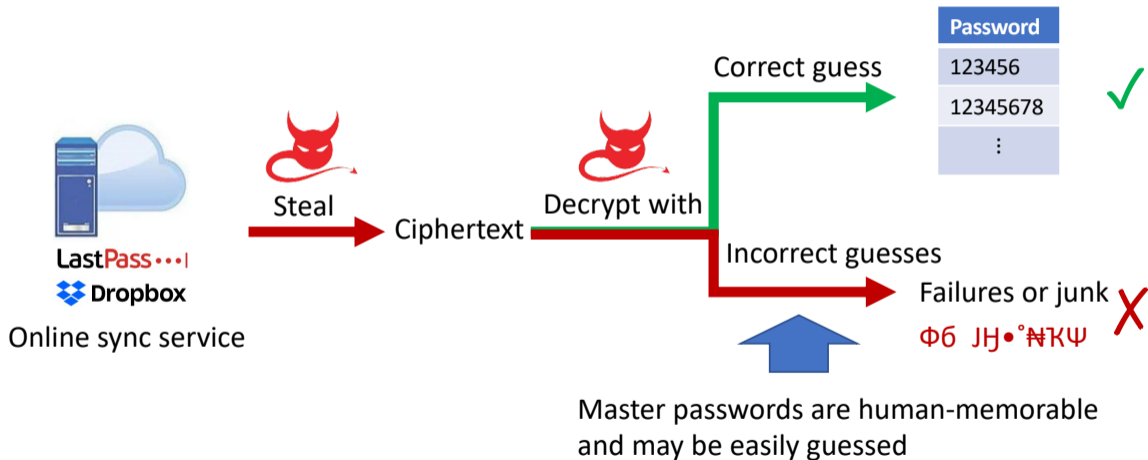
Website	Username	Password
Google	Aaron	123456
yahoo!	Aaron1	12345678
⋮	⋮	⋮
⋮	⋮	⋮



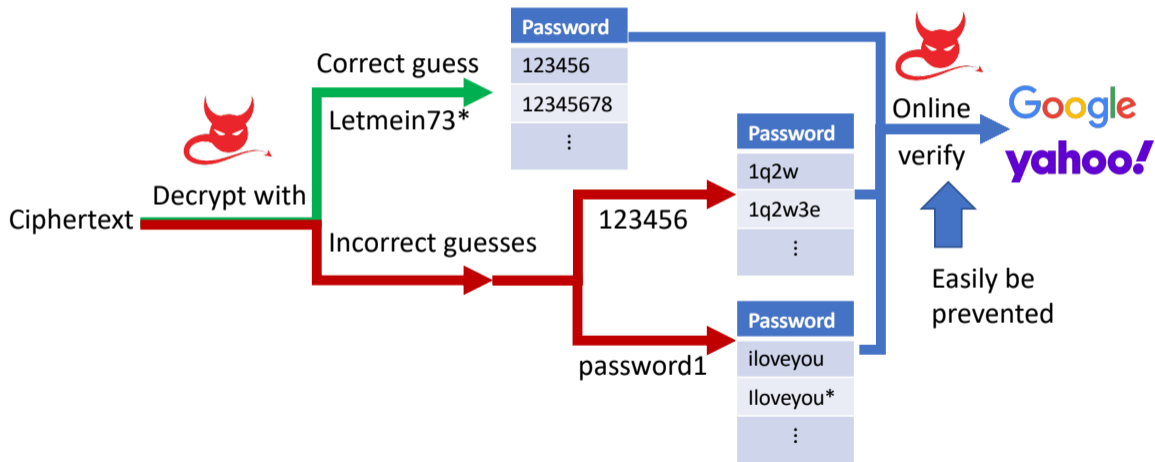
# Password vaults



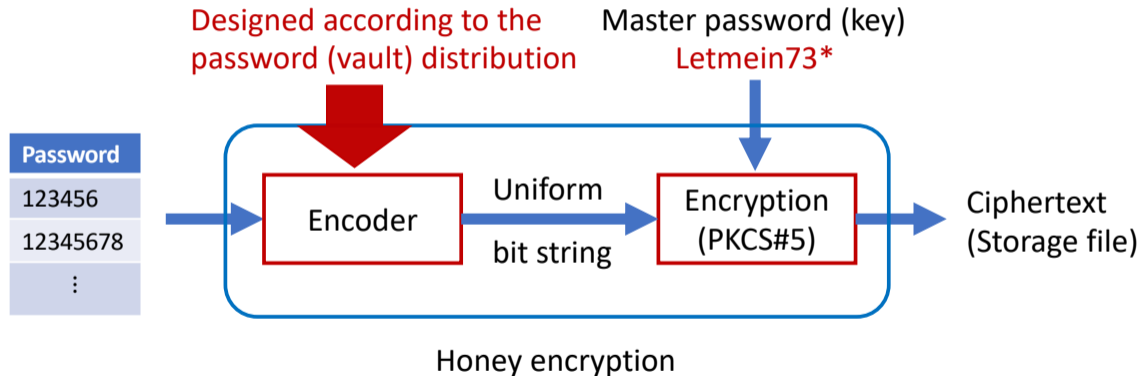
# Traditional password vaults suffer from offline guessing



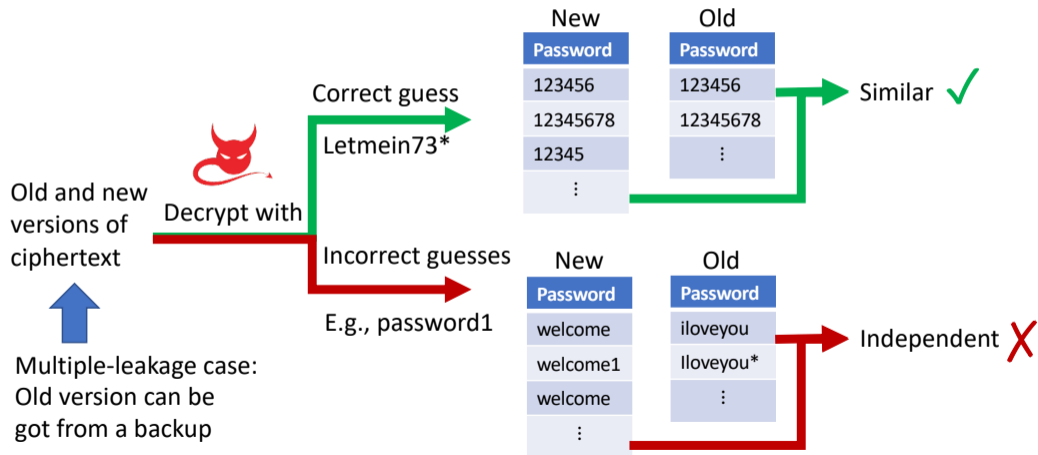
# Honey password vaults



# Honey password vaults: Design



# An open problem for honey password vaults: How to achieve update security if the user adds or changes a website password



# Our contribution

- ① New designs:
  - a A generic construction and an incremental update mechanism, achieving update security.
  - b An instantiation of the construction, generating more plausible-looking decoys.
- ② Security evaluation:
  - a Formally investigate the optimal strategy for distinguishing decoys and further propose practical attacks.
  - b Evaluate the current and our designs with the attacks.



# A generic construction for password vaults

## Probability model

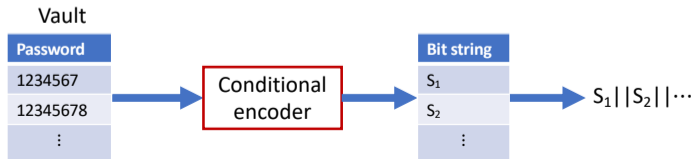
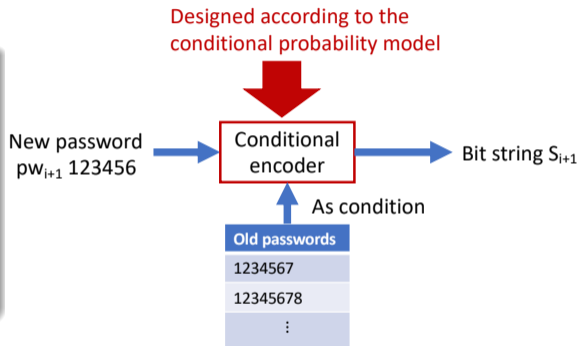
$$\Pr_{\text{real}}(V) = \prod_{i=0}^{n-1} \Pr_{\text{real}}(pw_{i+1} \mid pw_1, pw_2, \dots, pw_i). \quad (1)$$

- 1 Basic idea: A user generates the passwords one by one.  
 $\Pr_{\text{real}}(pw_{i+1} \mid pw_1, pw_2, \dots, pw_i)$  is the conditional probability that the user generates  $pw_{i+1}$  under given old passwords  $(pw_1, pw_2, \dots, pw_i)$ .
- 2 Our design: We use a conditional probability model  $\Pr_{\text{MSPM}}(\cdot|\cdot)$  (multi-similar-password model) to estimate  $\Pr_{\text{real}}(\cdot|\cdot)$ .

# A generic construction for password vaults

## Conditional encoder

- 1 Encode a new password  $pw_{i+1}$  given old passwords  $(pw_1, pw_2, \dots, pw_i)$ .
- 2 Designed according to the conditional probability model  $\Pr_{\text{MSPM}}(\cdot|\cdot)$  by Cheng et al.'s transformation [1].
- 3 Encode a vault password by password.



# An incremental update mechanism for password vaults

## Encoder+Encryption

- 1 Encoder: The conditional encoder.
- 2 Encryption: A prefix-keeping scheme, e.g., CTR-mode AES with PBKDF.

Old vault	Website	Username	Password index	Password	Bit string
	Google	Aaron	1	1234567	$S_1$
	yahoo!	Aaron1	2	12345678	$S_2$
	⋮	⋮	⋮	⋮	⋮

Add a password	Website	Username	Password index	Password	Bit string
	Google	Aaron	1	1234567	$S_1$
	yahoo!	Aaron1	2	12345678	$S_2$
	⋮	⋮	⋮	⋮	⋮
	facebook	Aaron	$i+1$	12345	$S_{i+1}$

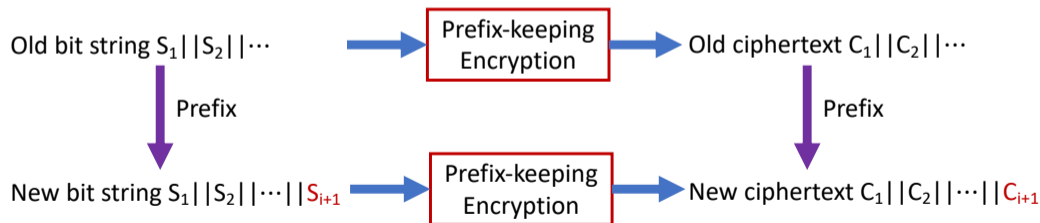
# An incremental update mechanism for password vaults

	Website	Username	Password index	Password	Bit string
Old vault	Google	Aaron	1	1234567	$S_1$
	yahoo!	Aaron1	2	12345678	$S_2$
	⋮	⋮	⋮	⋮	⋮
Delete a password	Google	Aaron	1	1234567	$S_1$
	yahoo!	Aaron1	Deleted	12345678	$S_2$
	⋮	⋮	⋮	⋮	⋮
Change a password	Google	Aaron	1	1234567	$S_1$
	yahoo!	Aaron1	$i+1$	12345678	$S_2$
	⋮	⋮	⋮	⋮	⋮
				12345	$S_{i+1}$

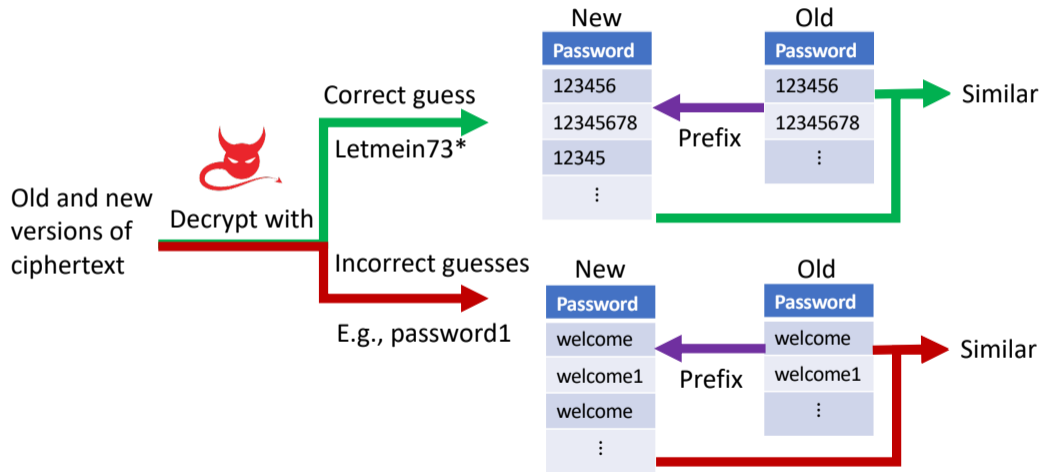
# An incremental update mechanism for password vaults

## Security

- 1 Old ciphertext is a prefix of the new one.
- 2 The attacker with the two versions of the ciphertext (multi-leakage case) degenerates to an attacker only with the current version (single-leakage case).



# Our incremental update mechanism achieves update security



# A new conditional probability model for password vaults

- ① Basic idea: A user generates a new password by 1) reusing an old one or 2) not (i.e., creating a brand new one).
- ② Construction:

$$\Pr_{\text{MSPM}}(pw_{i+1} \mid pw_1, pw_2, \dots, pw_i) \quad (2)$$

$$= \frac{1 - f(i)}{i} \sum_{i'=1}^i \Pr_{\text{SSPM}}(pw_{i+1} \mid pw_{i'}) + f(i) \Pr_{\text{SPM}}(pw_{i+1}).$$

- ③ Instantiation:

- a Single-similar-password model  $\Pr_{\text{SSPM}}$ : our simple design
- b Single-password model  $\Pr_{\text{SPM}}$ : Markov [2]
- c Unreused probability  $f$ : Nonlinear regression  $f(i) = 1/(\sum_{k=0}^3 a_k i^k)$

# How to evaluate the indistinguishability of decoys

## The optimal strategy of distinguishing real and decoy vaults

- 1 Ranking by the conditional probability of being real under the given ciphertext.  
The conditional probability for a vault  $V_i$  is proportional to the real-to-decoy probability ratio

$$\frac{\Pr_{\text{real}}(V_i)}{\Pr_{\text{decoy}}(V_i)}.$$

## Practical attacks

- 1 Cannot precisely calculate  $\Pr_{\text{real}}$ .
- 2 Estimate the ratio on the single password distribution and password-reuse features.



# Evaluating the existing and our honey vault schemes

## Experimental results

- 1 Our attack is more effective than the state-of-the-art attack (KL divergence attack [3]) against the existing schemes.
- 2 Our design brings 2.8x-7.5x online cost to attackers.

**Table 1:** The distinguishing accuracy of attacks against the honey vault schemes

Scheme	KL divergence attack [3]	Our attack
Chatterjee et al.'s [4]	86%	94%
Golla et al.'s [3] (static, $10^0$ )	52%	86%
Our design	58%	58%
Perfect design	50%	50%

## Future work

- ① The incremental update mechanism may apply to other applications using honey encryption.
- ② The probability model design for passwords and password vaults may be used for password guessing.

Thank you

# References I

- [1] Haibo Cheng et al. “Probability Model Transforming Encoders Against Encoding Attacks”. In: *USENIX Security 2019*, pp. 1573–1590.
- [2] Jerry Ma et al. “A Study of Probabilistic Password Models”. In: *IEEE S&P 2014*, pp. 538–552.
- [3] Maximilian Golla, Benedict Beuscher, and Markus Dürmuth. “On the security of cracking-resistant password vaults”. In: *ACM CCS 2016*, pp. 1230–1241.
- [4] Rahul Chatterjee et al. “Cracking-resistant password vaults using natural language encoders”. In: *IEEE S&P 2015*, pp. 481–498.